

ApproxJoin

Approximate Distributed Joins

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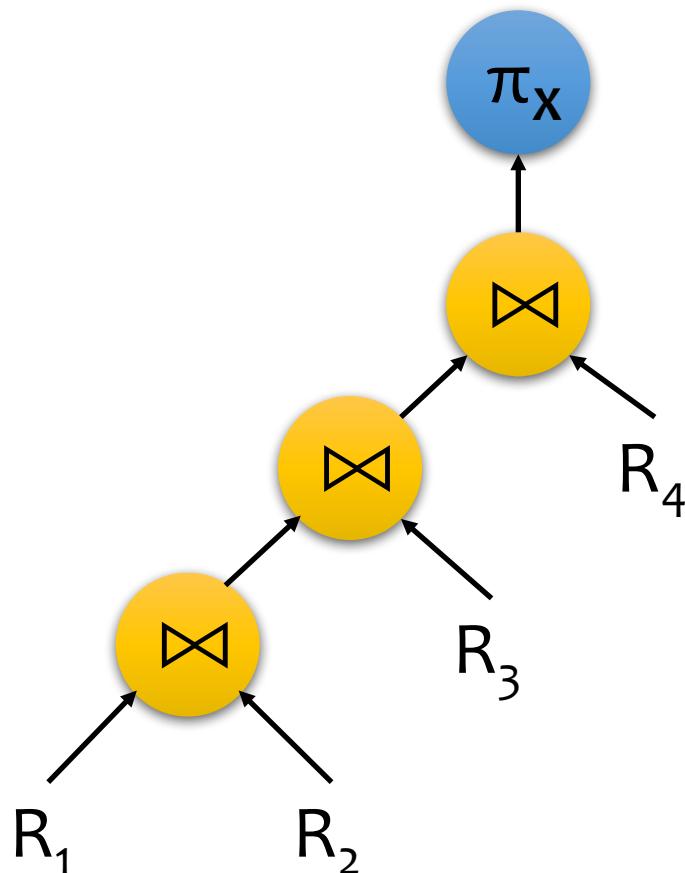
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Motivation



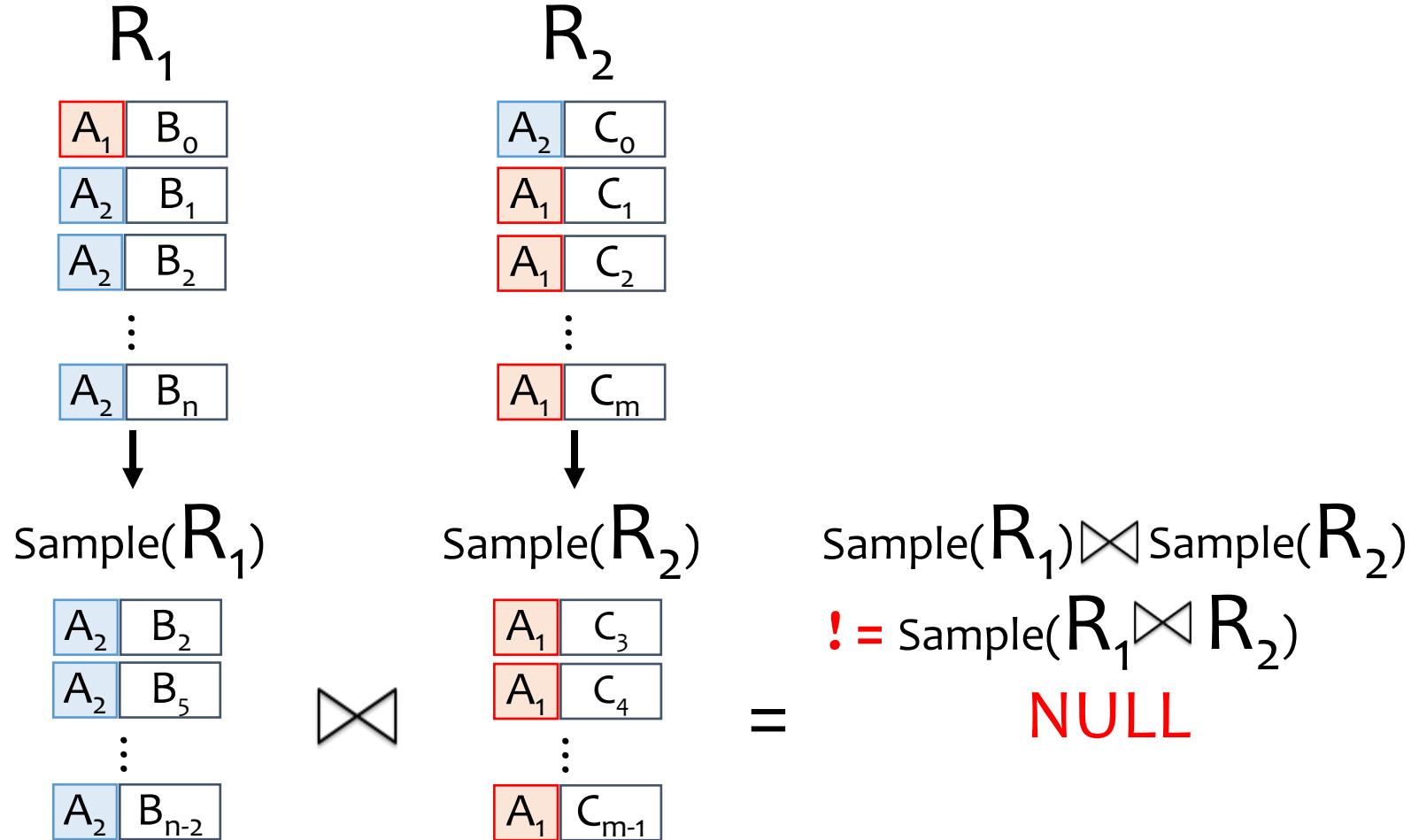
Join is a **critical** operation in big data analytics systems, but it is very **expensive**

Reduce the **overhead** of join operations using a **sampling-based** approach

Motivation

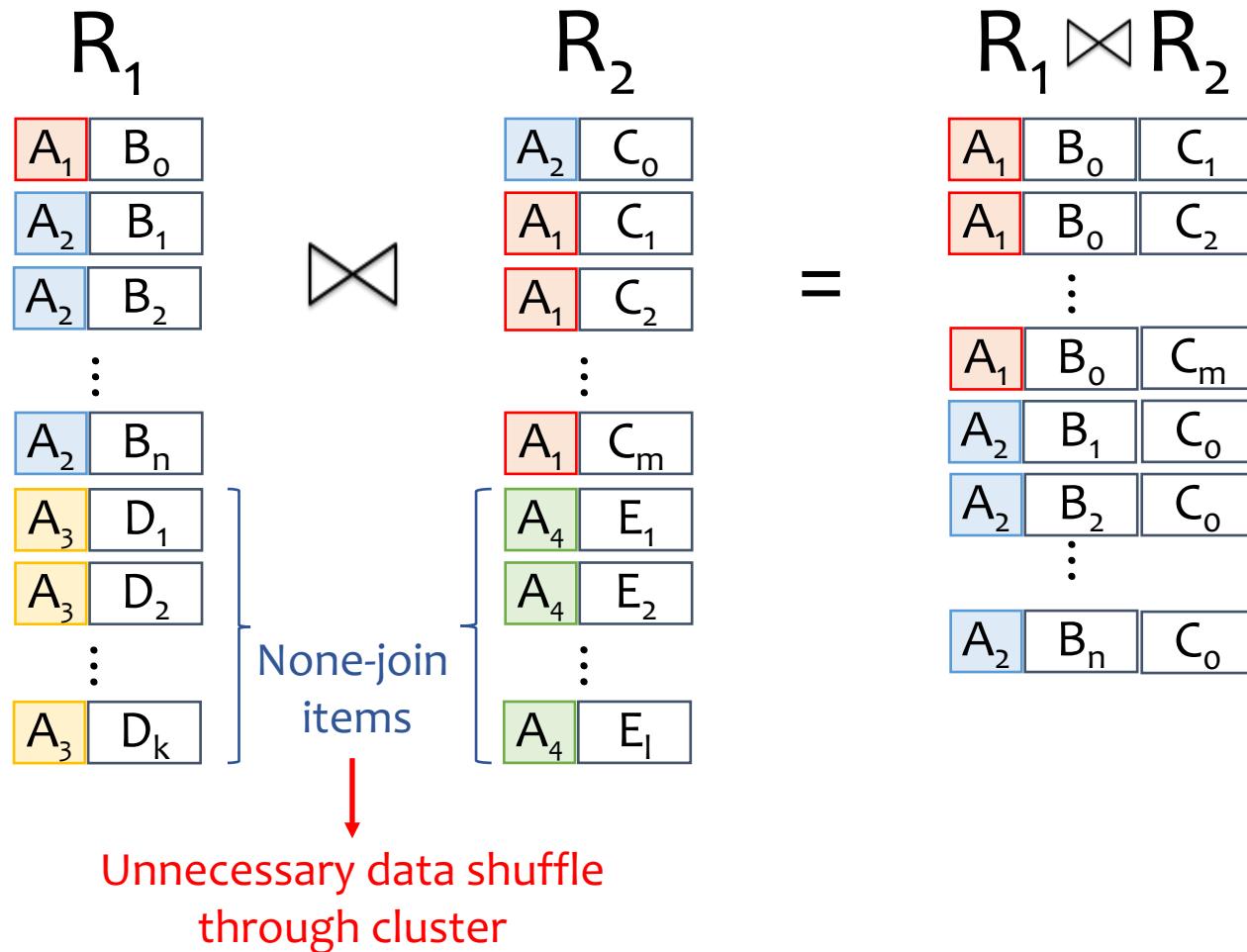
$$\begin{array}{c} R_1 \\ \begin{array}{|c|c|} \hline A_1 & B_0 \\ \hline A_2 & B_1 \\ \hline A_2 & B_2 \\ \vdots \\ A_2 & B_n \\ \hline \end{array} \end{array} \times \begin{array}{c} R_2 \\ \begin{array}{|c|c|} \hline A_2 & C_0 \\ \hline A_1 & C_1 \\ \hline A_1 & C_2 \\ \vdots \\ A_1 & C_m \\ \hline \end{array} \end{array} = \begin{array}{|c|c|c|} \hline A_1 & B_0 & C_1 \\ \hline A_1 & B_0 & C_2 \\ \hline \vdots & & \\ \hline A_1 & B_0 & C_m \\ \hline A_2 & B_1 & C_0 \\ \hline A_2 & B_2 & C_0 \\ \vdots \\ A_2 & B_n & C_0 \\ \hline \end{array}$$

Motivation



Sampling over joins is a challenging task regarding the output quality

Motivation



State-of-the-art Systems

AQUA (SIGMOD'99)
Sampling over joins (SIGMOD'99)

Requiring priori knowledge of inputs
(statistical info, indices)

RippleJoin (SIGMOD'99),
WanderJoin (SIGMOD'16)

Using online aggregation approach
for joins

SparkSQL (SIGMOD'15),
SnappyData (SIGMOD'16)

Using pre-existing samples to serve
queries

State-of-the-art Systems

AQUA (SIGMOD'99)

Sampling over joins (SIGMOD'99)

Requiring prior knowledge of inputs
(statistical info, indices)

Designed for single
node system

RippleJoin (SIGMOD'99)

WanderJoin (SIGMOD'16)

Using online aggregation approach
for joins

SparkSQL (SIGMOD'13)

SnappyData (SIGMOD'16)

Do not support
sampling over joins

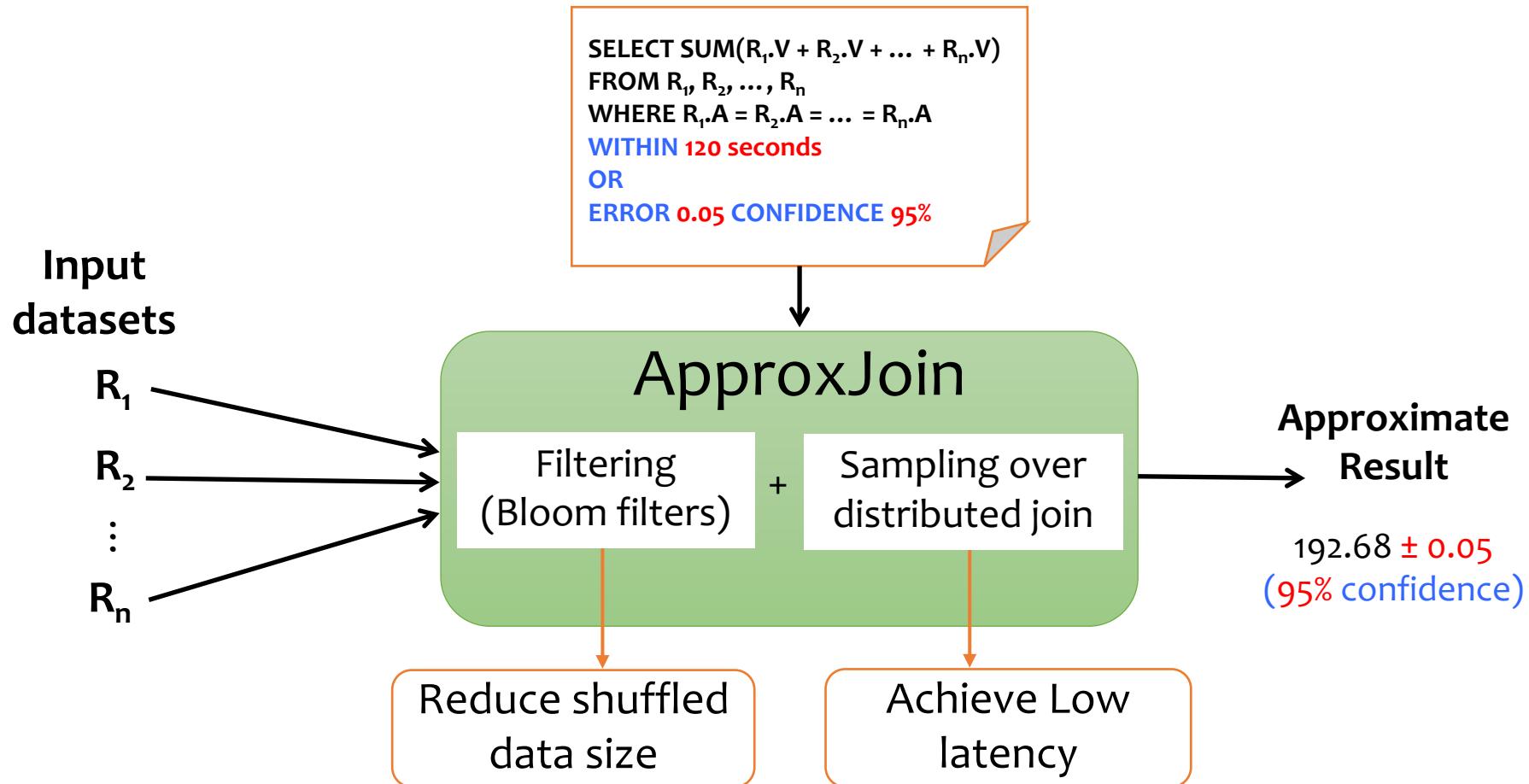
Sampling to serve
queries

Outline

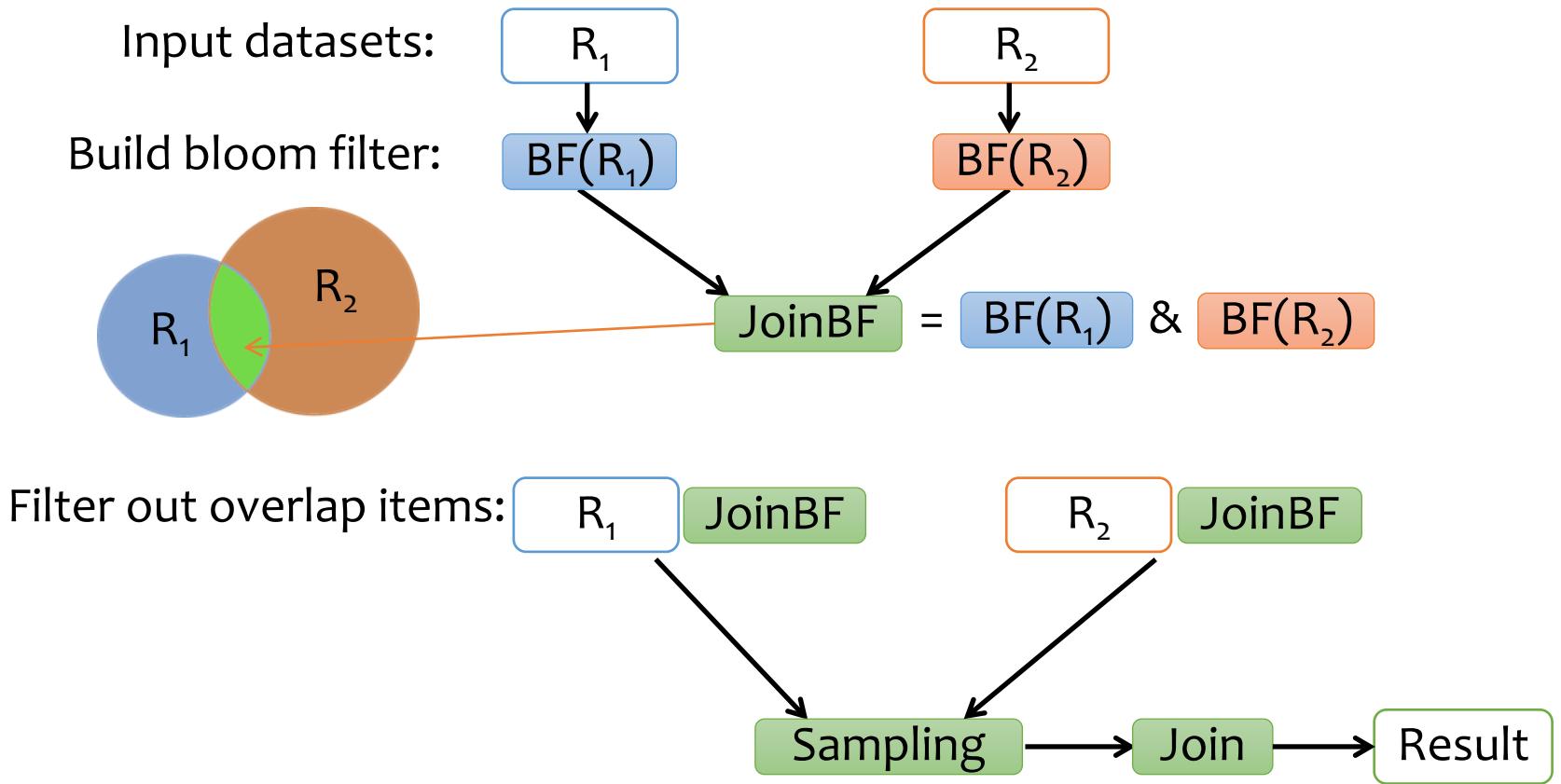
- Motivation

- Design
- Evaluation

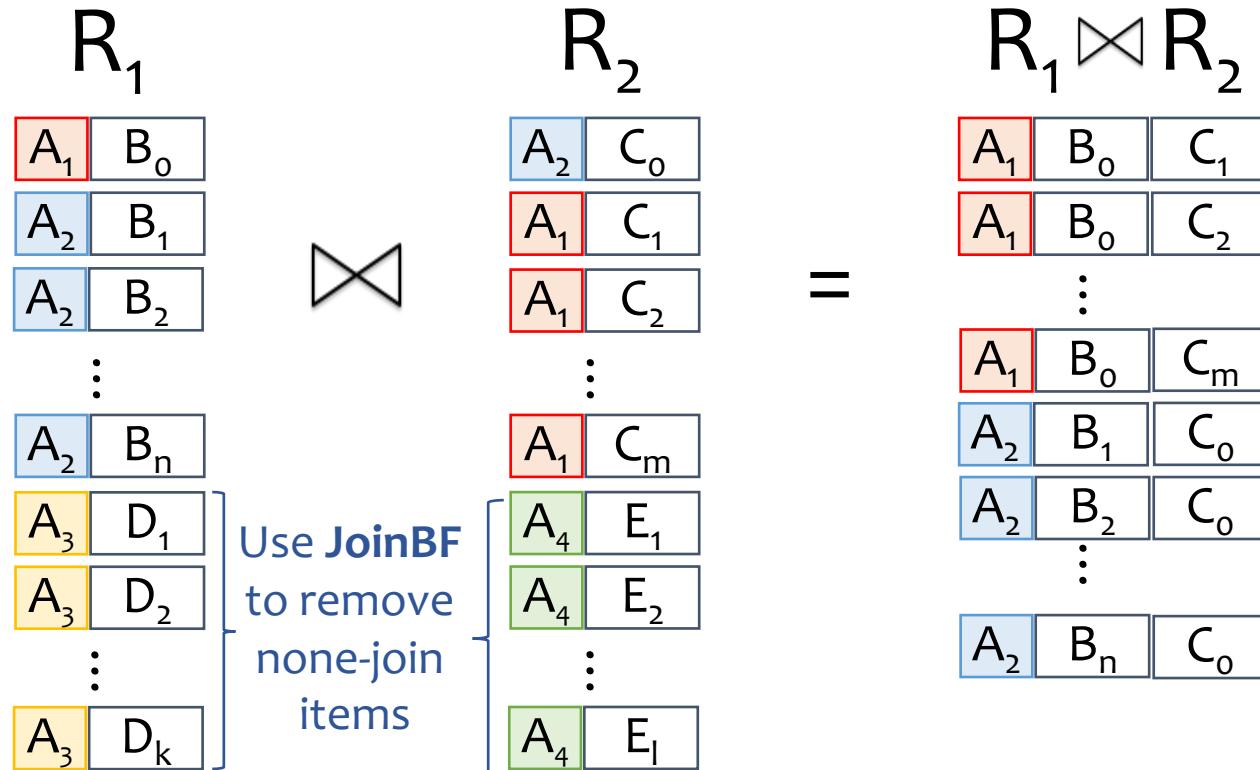
ApproxJoin: System Overview



ApproxJoin: Core Idea

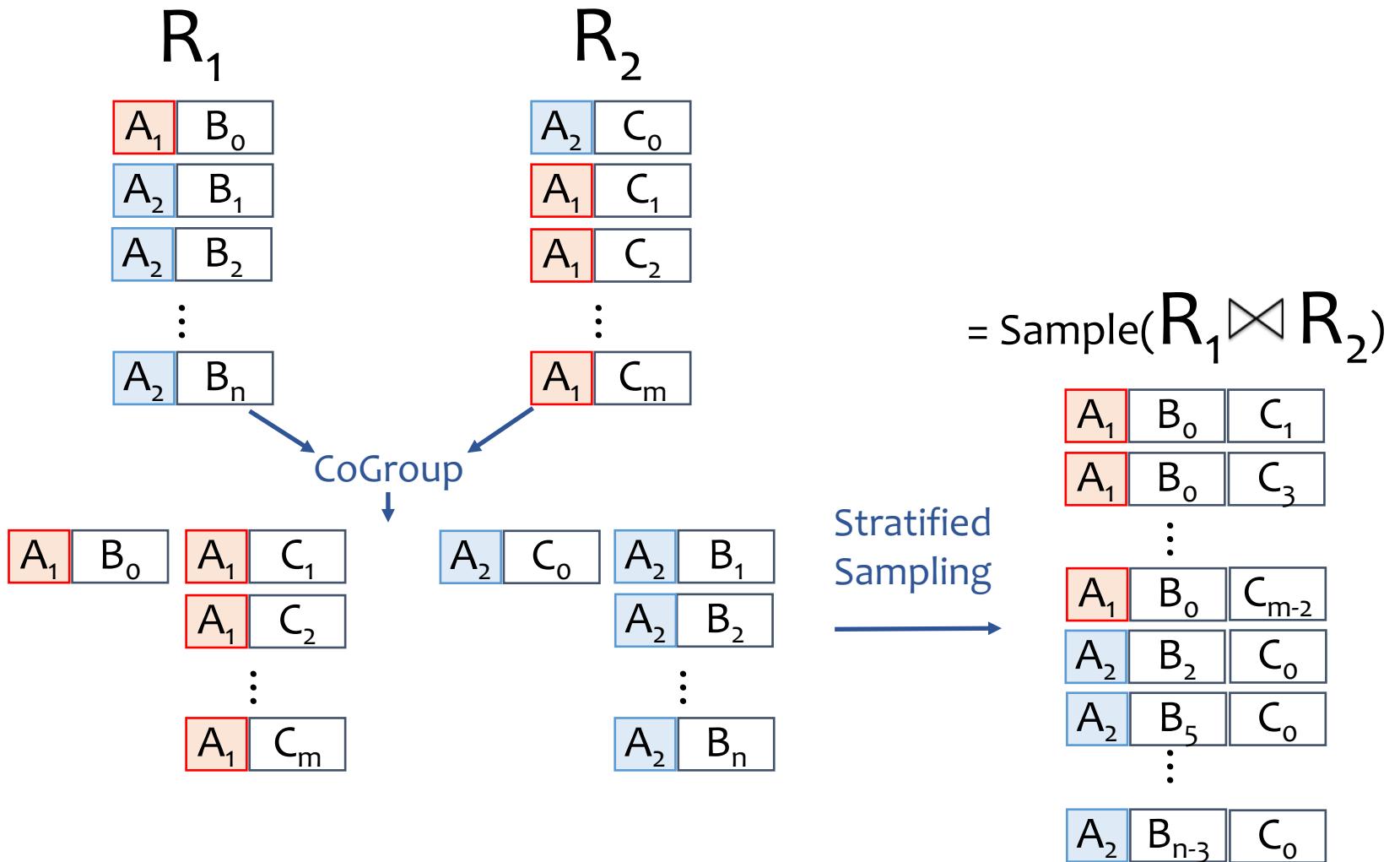


ApproxJoin: Filtering

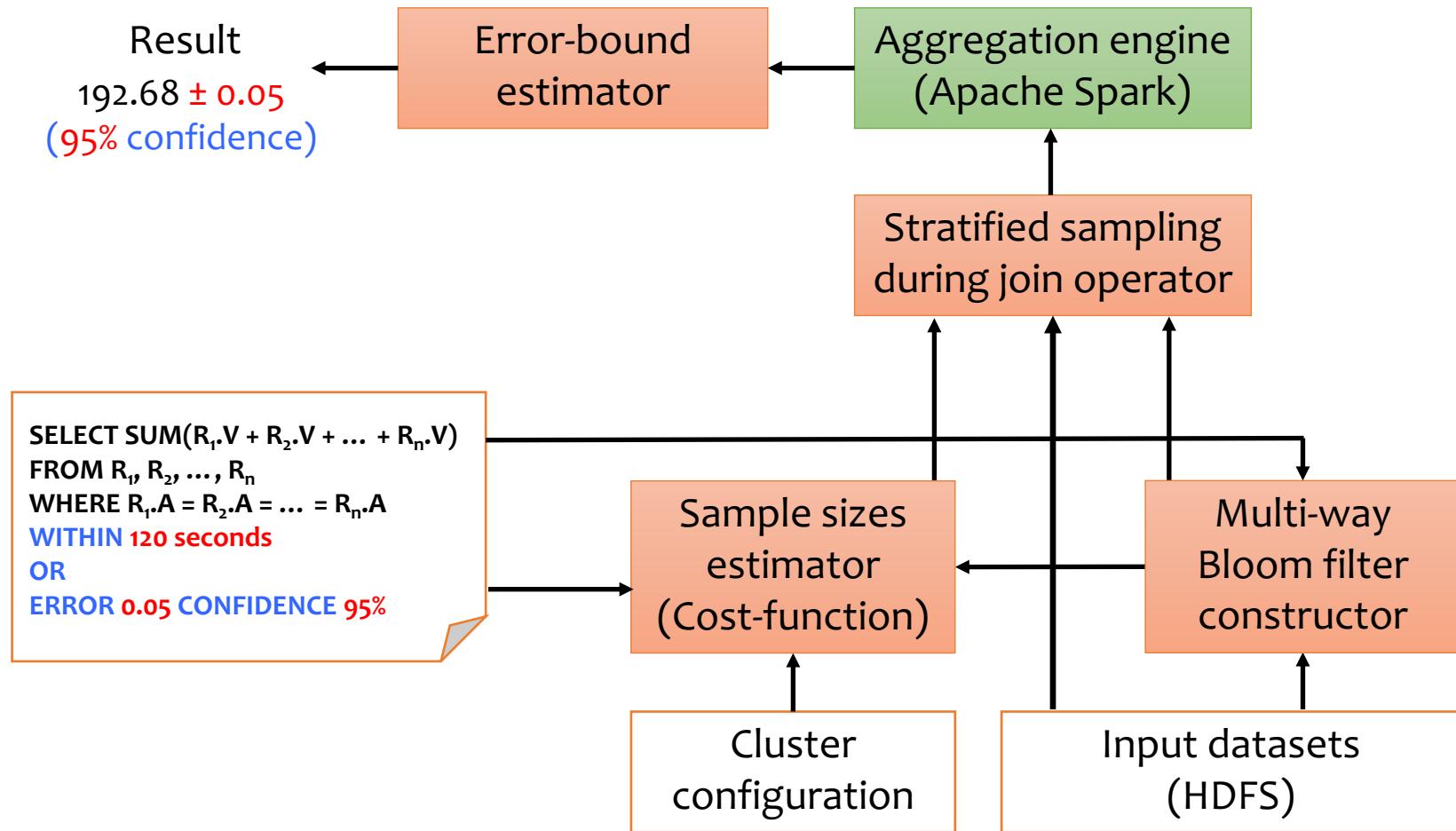


$$BF(R_1) = \{A_1, A_2, A_3\} \quad BF(R_2) = \{A_1, A_2, A_4\} \longrightarrow \text{JoinBF} = \{A_1, A_2\}$$

ApproxJoin: Sampling



ApproxJoin: Implementation



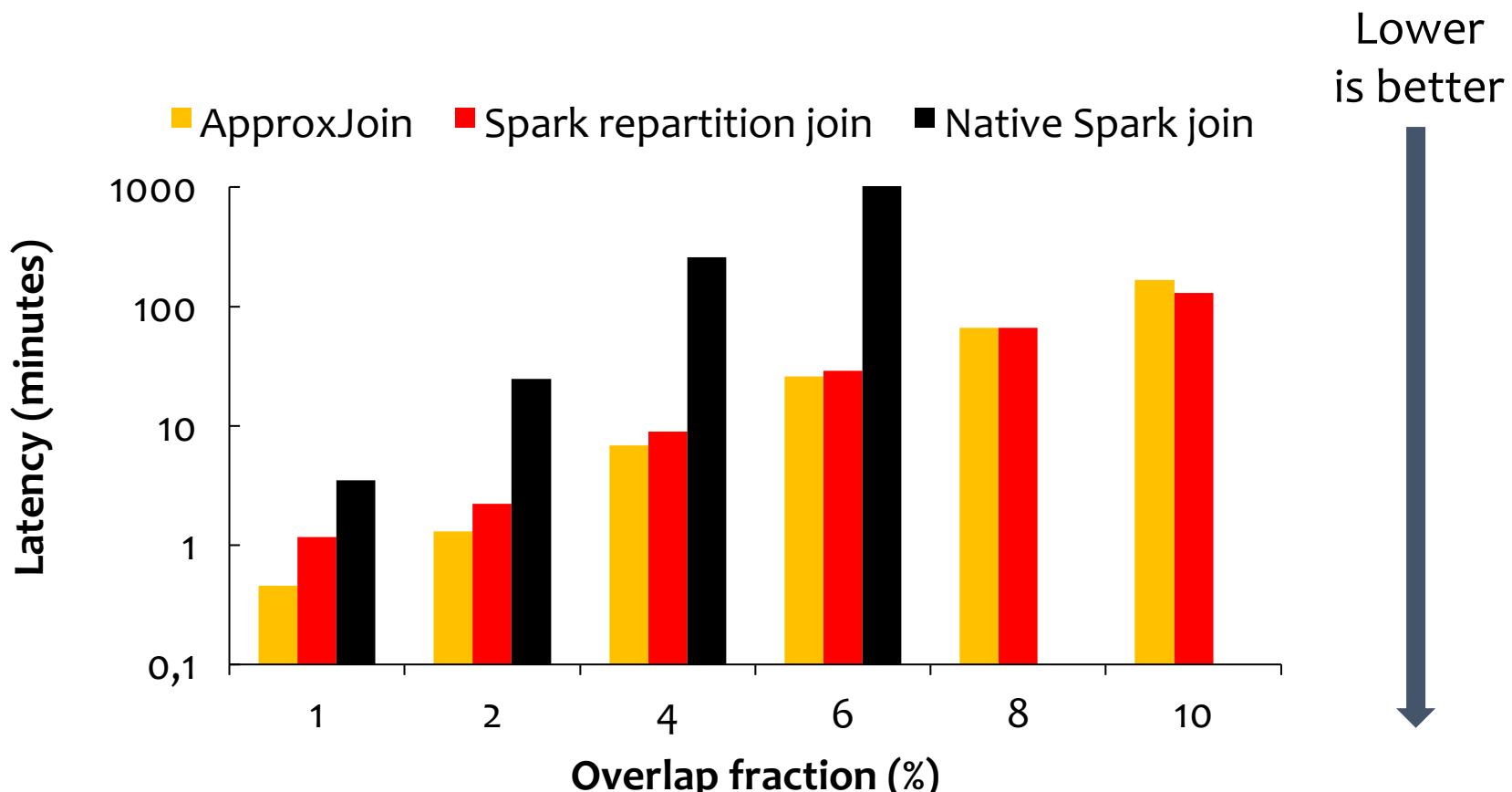
Outline

- Motivation
- Design
- Evaluation

Experimental Setup

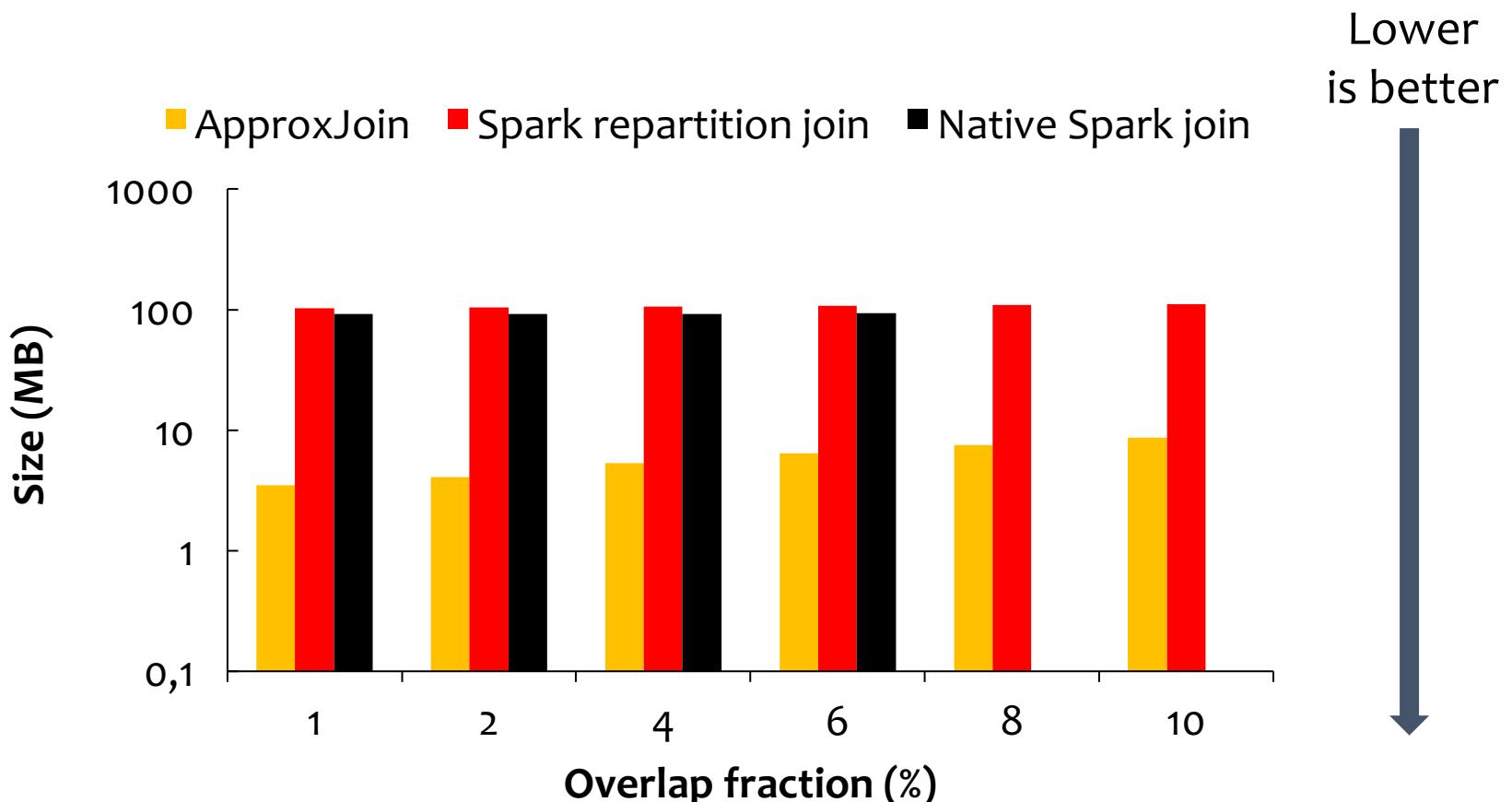
- Evaluation questions
 - Latency vs overlap fraction
 - Shuffled data size vs overlap fraction
 - Latency vs sampling fraction
 - Testbed
 - Cluster: 10 nodes
 - Datasets:
 - Synthesis: Poisson distribution datasets, TPC-H
 - CAIDA Network traffic traces; Netflix Prize
- 
- See the paper
for more
results!

Latency



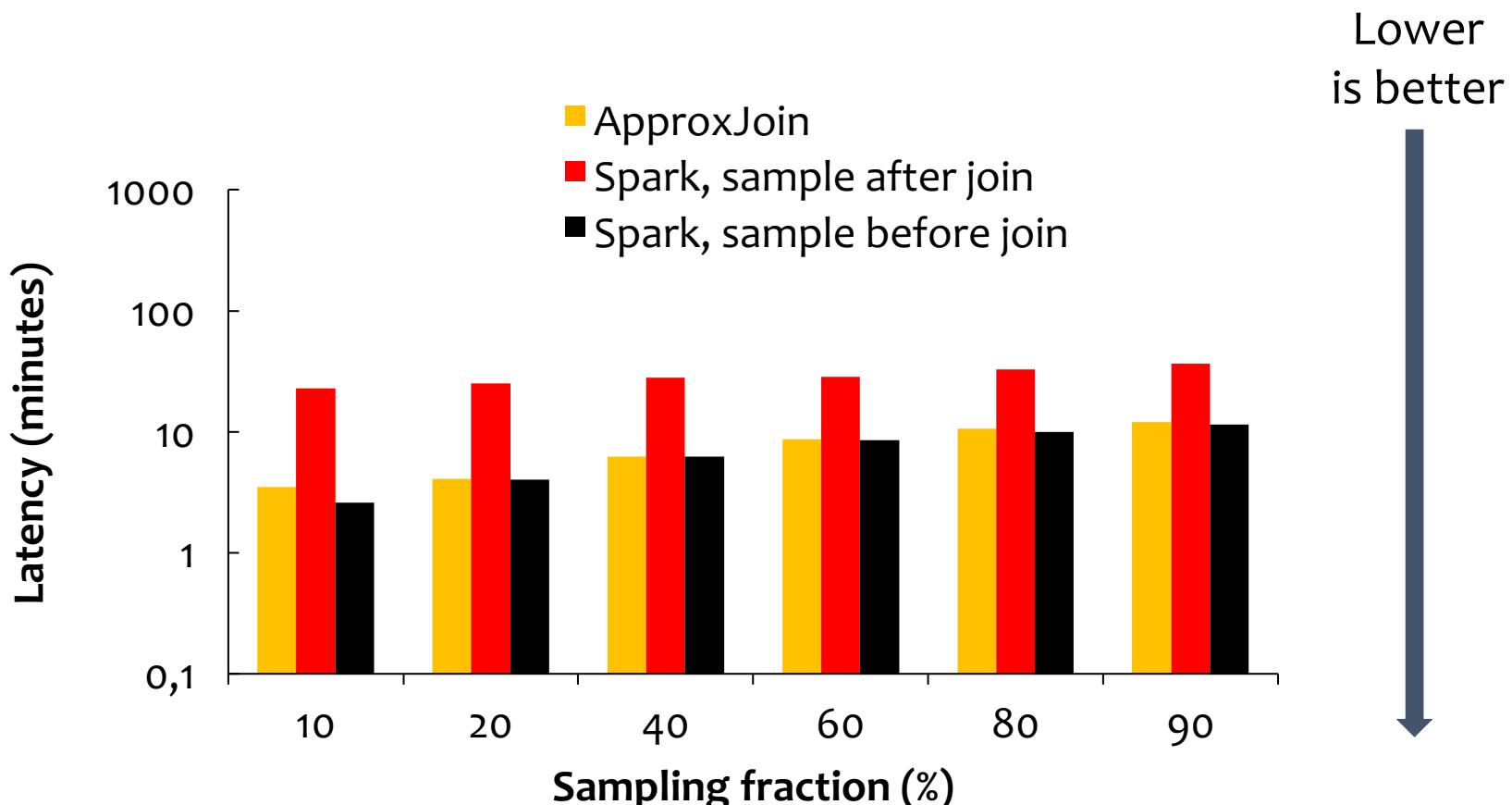
~2.6X and ~8X faster than Spark repartition join and native Spark join with overlap fraction of 1%

Shuffled Data Size



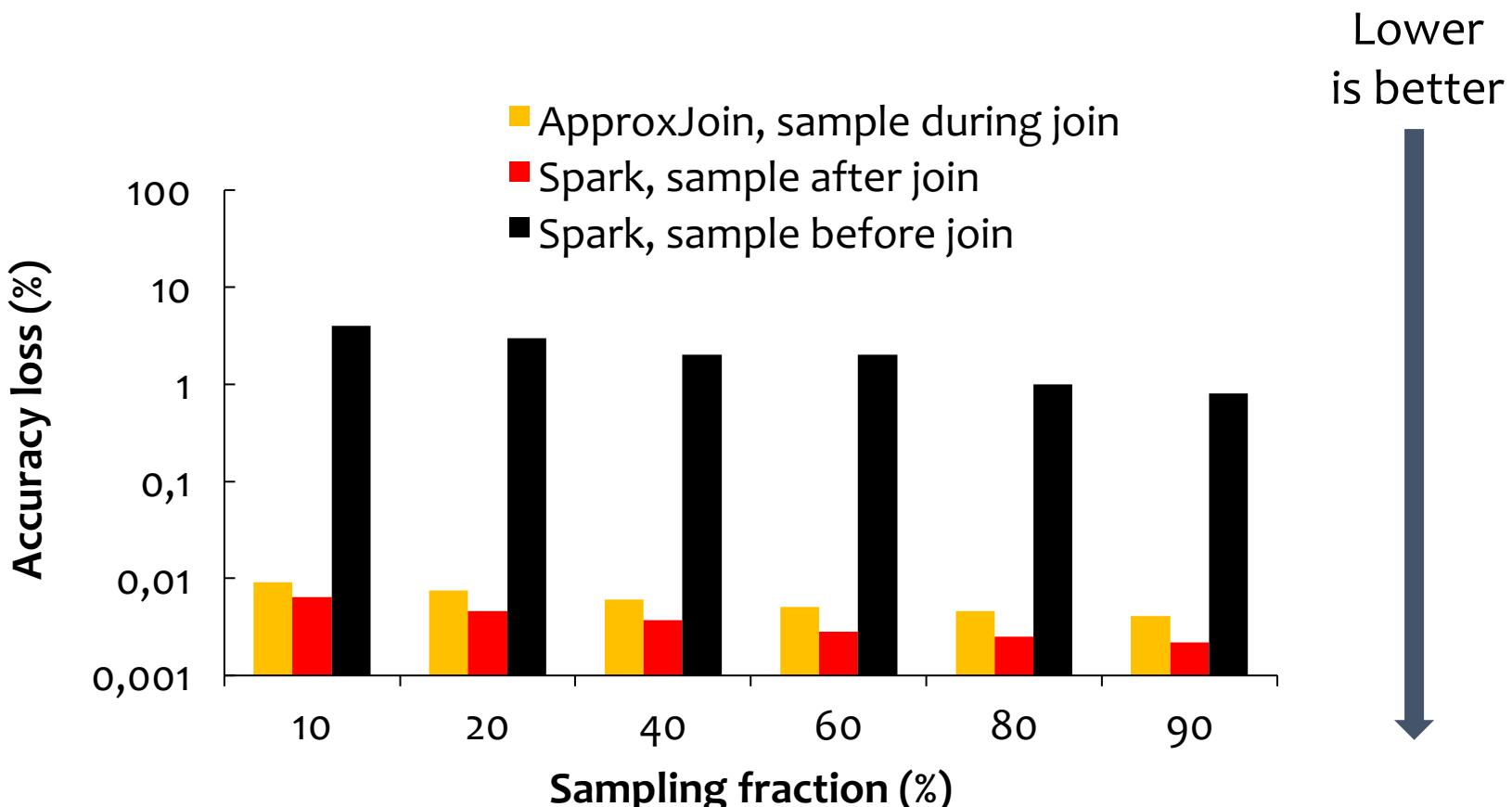
~29X and ~26X lower shuffled data size compared to Spark repartition join and native Spark join with overlap fraction of 1 %

Latency



(3X – 7X) faster than Spark with sampling *after* join
(1.01X – 1.3X) slower than Spark with sampling *before* join

Accuracy



Comparable accuracy to Spark with sampling *after* join
~42X more accurate than Spark with sampling *before* join

Outline

- Motivation
- Our work
- Conclusion

Conclusion

ApproxJoin: Approximate Distributed Joins

Transparent

Supports applications w/ minor code changes

Practical

Adaptive execution based on query budget

Efficient

Employs sketch & sampling techniques

Thank you!