

# Nested balanced ternary designs and Bhaskar Rao designs

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## Abstract

In this paper, we consider balanced ternary designs, BTDs, in which every block contains one element singly and the rest doubly. We call these packed BTDs, and we investigate three aspects of these designs: existence, nestings and signings. Construction methods generate classes of packed BTDs that are nested with balanced (BIBD) or partially balanced (PBIBD) incomplete block designs. Some of these classes are signed to produce  $c$ -Bhaskar Rao BTDs, most often with  $c = 0$ . Packed BTDs with block size three and five are studied in detail. The spectrum of possible indices for packed BTDs is determined. In particular, we prove every triple system with index  $3t$  is nested within a packed Bhaskar Rao balanced ternary design with  $K = 5$  and index  $8t$ . We give several new families of PBIBDs for block size 3, and show that each is nested within a BTD with block size 5 and index 4 or 6. We show that the necessary conditions are sufficient for the existence of this type of BTD, whether or not it has a design nested within.

## 1 Introduction

A balanced ternary design, or BTD, with parameters  $(V, B, R, K, \Lambda)$  is a collection of  $B$  blocks on  $V$  elements such that (1) each element occurs  $R$  times in the design, (2) each pair of distinct elements occurs  $\Lambda$  times in the design, and (3) each block contains  $K$  elements, where an element may occur 0, 1, or 2 times in a block (i.e., multiset).

BTDs are regular in the sense that each element occurs singly in  $\rho_1$  blocks and doubly in  $\rho_2$  blocks [6]. Because of this regularity, BTD parameters are most often given as  $(V; B; \rho_1, \rho_2, R; K; \Lambda)$ .

Counting arguments can be used to establish the necessity of the following well-known relationships among the parameters of any BTD:

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\* The authors wish to express their thanks to The Citadel Foundation for its support.

$$\begin{aligned} VR &= BK, \\ \Lambda(V - 1) &= \rho_1(K - 1) + 2\rho_2(K - 2), \text{ and} \\ R &= \rho_1 + 2\rho_2. \end{aligned}$$

An example of a BTD with parameters  $(4; 12; 3, 6, 15; 5; 16)$  appears in Figure 1. In the figure, each column represents a block of the design.

1	1	1	2	2	2	1	1	1	1	1	1
1	1	1	2	2	3	1	1	1	2	2	3
2	2	3	3	3	3	2	3	2	2	2	3
2	3	4	3	4	4	2	3	4	3	4	4
3	3	4	4	4	4	4	4	4	3	4	4

Figure 1: A BTD(4;12;3,6,15;5;16).

An alternative way of representing a BTD( $V; B; \rho_1, \rho_2, R; K; \Lambda$ ) is with a  $V$ -by- $B$  matrix,  $M$ , whose  $(i, j)^{th}$  entry denotes the number of times element  $i$  appears in block  $j$ .  $M$  is referred to as the incidence matrix of the BTD. The incidence matrix for the design in Figure 1 is given in Figure 2. It is straightforward to see that the inner product of any two distinct rows of an incidence matrix of a BTD( $V; B; \rho_1, \rho_2, R; K; \Lambda$ ) is  $\Lambda$ , the sum of the entries in any row  $R$ , and the sum of the entries in any column  $K$ .

	$b_1$	$b_2$	$b_3$	$b_4$	$b_5$	$b_6$	$b_7$	$b_8$	$b_9$	$b_{10}$	$b_{11}$	$b_{12}$
1	2	2	2	0	0	0	2	2	2	1	1	1
2	2	1	0	2	2	1	2	0	1	2	2	0
3	1	2	1	2	1	2	0	2	0	2	0	2
4	0	0	2	1	2	2	1	1	2	0	2	2

Figure 2: The incidence matrix for the BTD shown in Figure 1.

Sometimes it is possible to “sign” some of the entries in the rows of the incidence matrix of a BTD and still maintain constant row inner products. For example, changing all the ones to negative ones in the incidence matrix shown in Figure 2 produces a matrix where the inner product of any two rows is zero. Not every BTD incidence matrix can be signed. However, when one can be, the signed incidence matrix is called a Bhaskar Rao balanced ternary design or BRBTD.

Formally, a  $c$ -BRBTD( $V; B; \rho_1, \rho_2, R; K; \Lambda$ ) is a  $V$ -by- $B$  matrix  $\widehat{M}$  with entries from the set  $\{0, \pm 1, \pm 2\}$  such that (1) the inner product of any two distinct rows of  $\widehat{M}$  is  $c$ , and (2) the matrix  $M$  obtained from  $\widehat{M}$  by changing each negative entry of  $\widehat{M}$  to its corresponding absolute value yields the incidence matrix of a BTD( $V; B; \rho_1, \rho_2, R; K; \Lambda$ ).

A second way to modify the incidence matrix of a BTD is to change all the twos to ones. For example, changing the twos to ones in the matrix of Figure 2 produces the incidence matrix of a balanced incomplete block design, or BIBD, with parameters  $(4, 12, 9, 3, 6)$ . Not every BTD incidence matrix can be “reduced” to a

BIBD. However, when one can be, the BTB is nested as we will show in Corollary 2 given in Section 2.3.

A BIBD is a BTB where no element appears doubly in a block. Because  $\rho_2$  is zero in a BIBD, the parameters can be given as  $(v, b, r, k, \lambda)$  and the necessary parametric relationships as  $vr = bk$  and  $\lambda(v - 1) = r(k - 1)$ . These necessary conditions imply that the parameters are always fully specified by  $v, k$ , and  $\lambda$ . Thus, throughout the remainder of the paper when listing parameters of a BIBD we only list the triple  $(v, k, \lambda)$ .

The purpose of this paper is to describe a class of BTBs which can be used to construct BRBTBs and which are often nested. In Section 2, we define the class and show some general results. In Sections 3 and 4, respectively, we analyze designs with block size three and five. For each case we delineate the spectrum, construct related Bhaskar Rao designs, and identify nested designs.

For the interested reader, more details about balanced ternary designs appear in [2, 3, 9, 15, 17], about Bhaskar Rao designs in [4, 12, 14, 17, 18], and about nested designs in [10, 11, 13, 16].

## 2 Packed and Fully Packed BTBs

In this section, we introduce the concept of packed and fully packed BTBs and BRBTBs, derive necessary parametric relationships for them, and construct some infinite classes of them. We close the section with a discussion of nested designs and design reductions.

A block of odd size  $K = 2N + 1$  in a BTB is said to be *packed* if it contains one element that appears singly and  $N$  elements that appear doubly.

A BTB with odd block size  $K = 2N + 1$  is said to be packed if every block in the design is packed. For example, the BTB(3; 3; 1, 1, 3; 3; 2) with blocks  $\{1, 1, 2\}$ ,  $\{2, 2, 3\}$ , and  $\{3, 3, 1\}$  is a packed BTB.

A packed BTB with odd block size  $K = 2N + 1$  is said to be fully packed if it is simple (i.e., has no repeated blocks) and contains every possible packed block. Adding the blocks  $\{1, 2, 2\}$ ,  $\{2, 3, 3\}$ , and  $\{3, 1, 1\}$  to the above packed BTB gives a fully packed BTB(3; 6; 2, 2, 6; 3; 4).

A  $c$ -BRBTB is called packed [fully packed] if the underlying BTB is packed [fully packed]. The 0-BRBTB(3; 3; 1, 2, 5; 5; 8) shown in Figure 3 is fully packed.

	$b_1$	$b_2$	$b_3$
1	2	2	-1
2	2	-1	2
3	-1	2	2

Figure 3: A fully packed 0-BRBTB(3;3;1,2,5;5,8).

BTBs and BRBTBs can be packed without being fully packed. The 0-BRBTB(5; 10; 2, 4, 10; 5; 8) of Figure 4 is one such BRBTB.

	$b_1$	$b_2$	$b_3$	$b_4$	$b_5$	$b_6$	$b_7$	$b_8$	$b_8$	$b_{10}$
1	2	2	0	0	-1	0	0	2	2	-1
2	2	0	2	-1	0	0	2	0	-1	2
3	0	2	-1	2	0	2	0	-1	0	2
4	0	-1	2	0	2	2	-1	0	2	0
5	-1	0	0	2	2	-1	2	2	0	0

Figure 4: A packed 0-BRBTD(5;10;2,4,10;5;8) that is not fully packed.

The BTD associated with the 0-BRBTD(5; 10; 2, 4, 10; 5; 8) given in Figure 4 is the union of two BTD(5; 5; 1, 2, 5; 5; 4)s, namely  $\{b_1, b_2, b_3, b_4, b_5\}$  and  $\{b_6, b_7, b_8, b_9, b_{10}\}$ . However, the signing of the “whole” BTD is not a signing for either of the “half” BTDs.

## 2.1 Parametric relationships for packed BTDs

We begin by establishing necessary conditions for the parameters of packed BTDs. These conditions are then used to set bounds on the spectrum of fully packed and simple packed BTDs.

**Lemma 1** *If a BTD( $V; B; \rho_1, \rho_2, R; 2N + 1; \Lambda$ ) is packed, then  $\Lambda$  is even.*

PROOF: If  $D$  is a packed BTD, exactly one element per block appears singly. Thus whenever distinct  $v_1$  and  $v_2$  appear in the same block, one or both occurs with multiplicity two. The result follows.  $\square$

**Lemma 2** *If a BTD( $V; B; \rho_1, \rho_2, R; 2N + 1; \Lambda$ ) is packed, then*

- (a)  $B = \rho_1 V$ ,
- (b)  $\rho_2 = N\rho_1$ ,
- (c)  $R = K\rho_1$ , and
- (d)  $\Lambda(V - 1) = 4\rho_1 N^2$ .

PROOF: Assume  $D = \text{BTD}(V; B; \rho_1, \rho_2, R; 2N + 1; \Lambda)$  is a packed BTD. Since every block of  $D$  contains exactly one singleton and since every element is a singleton exactly  $\rho_1$  times, there must be  $\rho_1 V$  blocks in the design. Similarly counting doubletons we get  $NB = V\rho_2$ , which together with (a) implies  $\rho_2 = N\rho_1$ . The remaining two conditions follow from (a), (b), the necessary parametric conditions for BTDs stated in Section 1, and the fact that  $K = 2N + 1$ .  $\square$

**Theorem 1** *If a BTD( $V; B; \rho_1, \rho_2, R; 2N + 1; \Lambda$ ) is packed, then it is fully specified by the three parameters  $V, K$ , and  $\Lambda$ .*

PROOF: If a BTD( $V; B; \rho_1, \rho_2, R; 2N + 1; \Lambda$ ) is packed, the equation of part (d) of Lemma 2 can be used to define  $\rho_1$  in terms of  $V, K$ , and  $\Lambda$ . This formulation can then be substituted into (a) through (c) of that lemma, and the result follows.  $\square$

Because of the above result, throughout the remainder of the paper, we use the abbreviated notation  $BTD(V, K, \Lambda)$  or  $BTD(V, 2N + 1, \Lambda)$  when referring to a packed  $BTD(V; B; \rho_1, \rho_2, R; K = 2N + 1; \Lambda)$ .

**Theorem 2** *A fully packed  $BTD(V, 2N + 1, \Lambda)$  exists if and only if  $V \geq N + 1$  and  $\Lambda = 4\binom{V-1}{N}N^2/(V - 1)$ . When a fully packed  $BTD$  does exist for a triple  $(V, K, \Lambda)$ , it is unique.*

PROOF: *Since a fully packed design contains every possible packed block exactly once, these designs are straightforward to construct and clearly unique for a given set of parameters. On the other hand, if  $D = BTD(V, 2N + 1, \Lambda)$  is fully packed with element set  $E$ , then each element  $v$  must appear singly in a block with all possible  $N$ -tuples of  $E - \{v\}$ . Thus,  $\rho_1 = \binom{V-1}{N}$ . This, used in conjunction with Lemma 2 (d), implies  $\Lambda = 4\binom{V-1}{N}N^2/(V - 1)$ .  $\square$*

**Corollary 1** *If  $D = BTD(V, 2N + 1, \Lambda)$  is both simple and packed, then  $\Lambda \leq 4\binom{V-1}{N}N^2/(V - 1)$ .*

## 2.2 Constructing packed BTDS

Below we construct two infinite classes of packed designs. The method used in these constructions builds packed BTDS from already existing BIBDs. The spectrum of the BTDS produced by the second method is not as far reaching as that from the first method; however, the parameters of the BTDS are generally smaller. The designs from both constructions can be signed to produce BRBTDS.

Not all construction methods for building packed designs are dependent on having a starter design. For example see Theorem 6.

**Theorem 3** *Let  $V \geq 4$ . If there exists a  $BIBD(V, 3, \lambda)$ , then there exists a packed  $BTD(V, 7, 6\lambda(V - 3))$ . If the  $BIBD$  is not complete, then this  $BTD$  will not be fully packed. The incidence matrix of the constructed  $BTD$  can be signed to create a  $c$ -BRBTDS with  $c = 2\lambda(V - 3) = \Lambda/3$ .*

PROOF: Assume that  $V \geq 4$ , and that  $D_1$  is a  $BIBD$  with parameters  $(V, b, r, 3, \lambda)$  and elements  $\{1, 2, \dots, V\}$ . For each block  $\{v_1, v_2, v_3\}$  in  $D_1$  and each element  $v_4$  in  $\{1, 2, \dots, V\} \setminus \{v_1, v_2, v_3\}$ , define a new block  $\{v_1, v_1, v_2, v_2, v_3, v_3, v_4\}$ . Call the collection of all such new blocks  $D_2$ . By construction,  $D_2$  contains  $b(V - 3)$  size seven blocks over  $V$  elements. Since each element appears in  $r$  blocks of  $D_1$ , it will appear doubly in  $r(V - 3)$  blocks of  $D_2$  and singly in  $(b - r)$  blocks. Also, since each pair of distinct elements appear together in  $\lambda$  blocks of  $D_1$ , and each without the other in  $r - \lambda$  blocks, each distinct pair will appear both doubly in  $\lambda(V - 3)$  blocks of  $D_2$  and one singly and the other doubly in  $2(r - \lambda)$  blocks. Thus,  $D_2$  is a  $BTD$  where each pair of distinct elements appears in  $4\lambda(V - 3) + 2 \cdot 2(r - \lambda)$  blocks. The necessary conditions on  $BIBD$  parameters imply that  $r = \lambda(V - 1)/2$ . This can be used to rewrite the index (pair repetition number) given above as  $6\lambda(V - 3)$ . If  $D_1$  is not complete,  $D_2$  will not be fully packed.

Assume  $M$  is the incidence matrix for  $D_2$ . Change each one in  $M$  to a negative one. Recall that since each pair of distinct elements  $v_1$  and  $v_2$  of  $D_1$  appear together in  $\lambda$  blocks of  $D_1$  and each without the other in  $r - \lambda$  blocks, each pair will appear both doubly in  $\lambda(V - 3)$  blocks of  $D_2$  and one singly and the other doubly in  $2(r - \lambda)$  blocks. Thus, in the signed matrix the inner product of row  $v_1$  and row  $v_2$  will be  $4\lambda(V - 3) - 2 \cdot 2(r - \lambda) = 2\lambda(V - 3)$ .  $\square$

Before leaving Theorem 3 we mention that it is well known that the necessary conditions for a BIBD( $v, 3, \lambda$ ) are sufficient [10].

A BIBD is said to be resolvable if the design blocks can be partitioned into classes, called parallel classes, such that each element of the design appears exactly once in the blocks of a class.

A BIBD is said to be near-resolvable if the design blocks can be partitioned into classes (near-parallel classes) such that each class is missing exactly one element and every element of the design is absent from exactly one class.

If a BIBD( $v, K, \lambda$ ) is near-resolvable, then each element must appear once in  $v - 1$  classes so  $r = v - 1$ , which forces  $\lambda = K - 1$ . Since each class is missing a single element and since each block has size  $K$ ,  $v \equiv 1 \pmod{K}$ .

The necessary conditions for the existence of a near resolvable design are known to be sufficient for  $K \leq 7$ . For  $K \leq 5$ , see Table I.6.26 in [1]. For  $K = 6$  see [2]. In [3], Abel et al. reduced the exception list for  $K = 7$  to five open cases. These five cases have since been solved by Abel [4].

**Theorem 4** *If there exists a near-resolvable BIBD( $V, K, K - 1$ ), then there exists a packed BTD( $V, 2K + 1, 4K$ ). The incidence matrix of the constructed BTD can be signed to create a  $4(K - 2)$ -BRBTD.*

PROOF: Assume  $D_1$  is a near-resolvable BIBD( $V, K, K - 1$ ), where  $C_{v_i}$  is the near-parallel class missing element  $v_i$ . For each element  $x$  and each block  $\{v_1, \dots, v_K\}$  in  $C_x$ , define a new block  $\{v_1, v_1, \dots, v_K, v_K, x\}$ . Call the collection of all such new blocks  $D_2$ . By construction,  $D_2$  contains  $b$  size  $2K + 1$  blocks over  $V$  elements. Since each element appears in  $V - 1$  blocks of  $D_1$  and is missing from exactly one class, each element will appear doubly in  $V - 1$  blocks of  $D_2$  and singly in  $b/V = (V - 1)/K$  blocks. Also, since each pair of distinct elements appear together in  $K - 1$  blocks of  $D_1$ , and each is missing from exactly one near-resolvable class, each distinct pair will appear both doubly in  $K - 1$  blocks of  $D_2$  and one singly and the other doubly in two blocks. Thus,  $D_2$  is a BTD where each pair of distinct elements appears  $4(K - 1) + 4 = 4K$  times.

Assume  $M$  is the incidence matrix for  $D_2$ . Change each one in  $M$  to a negative one. Recall that since each pair of distinct elements will appear together in  $K - 1$  blocks of  $D_1$ , and each is missing from exactly one near-resolvable class, each distinct pair will appear both doubly in  $K - 1$  blocks of  $D_2$  and one singly and the other doubly in two blocks. Thus, in the signed matrix the inner product of row  $v_1$  and row  $v_2$  will be  $4(K - 1) - 4 = 4(K - 2)$ .  $\square$

### 2.3 Nested packed BTDs and BTD reductions

We next turn our attention to nested designs and reductions, and their relevance to packed BTDs.

If each block  $b_i$  of a design  $D_1$  can be partitioned into sub (multi) sets  $c_{ij}$ ,  $j = 1, 2, \dots, m$  such that  $\{c_{ij}\}_{i=1}^{i=B}$  is a design (with the assigned decomposition of blocks) for each  $j$ , then  $D_1$  is said to be a nested design ([14], [15], [17]).

**Theorem 5** *If  $D = BTD(V, K, \Lambda)$  is a fully packed BTD, then  $D$  is nested.*

PROOF: Let  $K = 2N + 1$ . For each block  $b_i = \{v_1, v_1, v_2, v_2, \dots, v_N, v_N, v_{N+1}\}$  in  $D$ , let  $c_{i1} = \{v_1, v_2, \dots, v_{N+1}\}$  and  $c_{i2} = \{v_1, v_2, \dots, v_N\}$ . Both  $D_1 = \{c_{i1}\}_{i=1}^{i=B}$  and  $D_2 = \{c_{i2}\}_{i=1}^{i=B}$  are designs. Here  $D_1$  is a multiple of the complete (containing every possible  $N + 1$ -tuple) BIBD( $V, N + 1, \lambda_1$ ), and  $D_2$  is also a multiple of a complete BIBD( $V, N, \lambda_2$ ) design where the blocks have size  $N$ .  $\square$

This definition of a nested *design* does not specify the design type of the sub-designs. In fact, it does not even demand that the sub-designs all be of the same type. For example, when  $K = 3$ , the design  $D_1$  in the proof of Theorem 5 is a BIBD( $V, 2, 2$ ), while  $D_2$  is a one-design, i.e., a set of elements.

Fully packed BTDs are not the only packed BTDs that are nested. The subdivision of blocks used in Theorem 5 can be applied to the BTDs of Theorems 3 and 4 to show that they too are nested.

Nested BTDs have occurred in the literature previously. Theorem 6 is one such example.

**Theorem 6** [21], [17]. *If  $x$  be a primitive root of  $GF(V)$ , where  $V = 2t+1$  is a prime or prime power and  $t$  is odd, then the initial block  $\{0, x^0, x^0, x^2, x^2, \dots, x^{2t-2}, x^{2t-2}\}$  gives, when developed additively over  $GF(V)$ , a nested packed BTD which reduces to two BIBDs with initial sub-blocks  $\{0, x^0, x^2, \dots, x^{2t-2}\}$  and  $\{x^0, x^2, \dots, x^{2t-2}\}$ .*

Somewhat differently, we say  $D_2$  is a *reduction* of a packed BTD  $D_1$ , if  $D_2$  is a design and  $\{v_1, v_2, \dots, v_{N+1}\}$  is a block of  $D_2$  if and only if  $\{v_1, v_1, v_2, v_2, \dots, v_N, v_N, v_{N+1}\}$  is a block of  $D_1$ . For example, the BIBD(4, 3, 6) with blocks  $\{1, 2, 3\}$ ,  $\{1, 2, 3\}$ ,  $\{1, 3, 4\}$ ,  $\{2, 3, 4\}$ ,  $\{2, 3, 4\}$ ,  $\{2, 3, 4\}$ ,  $\{2, 3, 4\}$ ,  $\{1, 2, 4\}$ ,  $\{1, 3, 4\}$ ,  $\{1, 2, 4\}$ ,  $\{1, 2, 3\}$ ,  $\{1, 2, 4\}$  and  $\{1, 3, 4\}$  is a reduction of the packed BTD(4; 12; 3, 6, 15; 5; 16) shown in Figure 1 of Section 1.

If  $D_1$  is a nested design with reduction  $D_2$ , it is possible that  $D_2$  is one of the subdesigns of  $D_1$  that makes up the nesting. We illustrate this using Theorem 6. Consider the field of seven elements  $\{0, 1, \dots, 6\}$  with operations mod 7. The element 3 is a generator of the cyclic group of non-zero elements ( $3^1 = 3, 3^2 = 2, 3^3 = 6, 3^4 = 4, 3^5 = 5, 3^6 = 1$ ). Setting  $x = 3$ , we get  $\{x^0, x^2, x^4\} = \{1, 2, 4\}$ . Thus the starter block  $\{0, 1, 1, 2, 2, 4, 4\}$  additively generates a packed BTD(7, 7, 6). By splitting the starter block into  $\{0, 1, 2, 4\}$  and  $\{1, 2, 4\}$  one obtains starter blocks for a BIBD(7, 4, 2) and a BIBD(7, 3, 1) respectively, thus showing that the BTD is nested. Clearly the BIBD(7, 4, 2) is a reduction of the BTD(7, 7, 6).

Below we give necessary and sufficient conditions for a packed BTD to reduce to a BIBD.

For every distinct pair of elements  $v_1$  and  $v_2$  in a packed BTD, define  $d_2(v_1, v_2)$  to be the number of blocks which contain one of  $v_1$  or  $v_2$  singly and the other doubly, and define  $d_4(v_1, v_2)$  to be the number of blocks which contain each of  $v_1$  and  $v_2$  doubly.

**Theorem 7** *A packed BTD( $V, 2N + 1, \Lambda$ ) reduces to a BIBD( $V, N + 1, \lambda$ ), if and only if  $d_2$  and  $d_4$  are constants independent of the choice of the elements  $v_1$  and  $v_2$ .*

PROOF: Assume the packed BTD( $V, 2N + 1, \Lambda$ ) reduces to a BIBD( $V, N + 1, \lambda$ ). Then if  $v_1, v_2$  and  $v_3, v_4$ , are two pairs of distinct elements of the design, the equations  $\Lambda = 2d_2(v_1, v_2) + 4d_4(v_1, v_2) = 2d_2(v_3, v_4) + 4d_4(v_3, v_4)$  and  $\lambda = d_2(v_1, v_2) + d_4(v_1, v_2) = d_2(v_3, v_4) + d_4(v_3, v_4)$  yield  $(\Lambda - 2\lambda)/2 = d_4(v_1, v_2) = d_4(v_3, v_4)$  and  $(4\lambda - \Lambda)/2 = d_2(v_1, v_2) = d_2(v_3, v_4)$  and the result follows.

Conversely, let  $D_1$  be a packed BTD( $V, 2N + 1, \Lambda$ ) with  $d_2$  and  $d_4$  constants independent of any  $v_1$  and  $v_2$ . Further, let  $D_2$  be the collection of blocks obtained from  $D_1$  by changing each doubleton in a  $D_1$  block to a singleton. Since  $D_1$  is packed, each block in  $D_2$  will be of size  $N + 1$ . The number of blocks a pair of elements appears in does not change in going from  $D_1$  to  $D_2$  blocks, only the multiplicity changes, i.e. all multiplicities are reduced to one. Thus, each pair of elements will appear together in  $\lambda = d_2 + d_4$  blocks. Finally since  $\rho_1$  and  $\rho_2$  are constants for  $D_1$ , each element will appear in  $r = \rho_1 + \rho_2$  reduced blocks. We have shown  $D_2$  is a block system with a constant replication number for each element and a constant index for each pair of elements. Thus, it follows that  $D_2$  is a BIBD( $V, N + 1, \lambda$ ).  $\square$

**Corollary 2** *If a packed BTD( $V, 2N + 1, \Lambda$ )  $D_1$  reduces to a BIBD( $V, N + 1, \lambda$ ),  $D_2$ , then  $D_1$  is nested. One nesting consists of  $D_2$  and a second BIBD.*

PROOF: Assume  $D_1$  is a packed BTD( $V, 2N + 1, \Lambda$ ) that reduces to  $D_2$  a BIBD( $V, N + 1, \lambda$ ). For each block  $b_i = \{v_1, v_1, v_2, v_2, \dots, v_N, v_N, v_{N+1}\}$  in  $D_1$ , let  $c_{i1} = \{v_1, v_2, \dots, v_{N+1}\}$  and  $c_{i2} = \{v_1, v_2, \dots, v_N\}$ . By assumption  $D_2 = \{c_{i1}\}_{i=1}^{i=B}$  is a BIBD( $V, N + 1, \lambda$ ). Thus Theorem 7 tells us that  $d_2$  and  $d_4$  are constants such that  $\Lambda = 2d_2 + 4d_4$  and  $\lambda = d_2 + d_4$ . Let  $D_3 = \{c_{i2}\}_{i=1}^{i=B}$ . Since each element appears in  $\rho_2$  blocks of  $D_1$  as a doubleton, each element will appear in  $\rho_2$  blocks of  $D_3$ . Also since each pair of distinct elements appear together in  $d_2 + d_4$   $D_1$  blocks and  $d_2$   $D_2$  blocks, it follows that they appear together in  $d_4$   $D_3$  blocks. We have shown  $D_3$  is an equireplicate block system, with block size  $N$ , and constant index for all pairs of elements. Thus,  $D_3$  will be a BIBD( $V, N, d_4$ ).  $\square$

### 3 Analysis of packed BTDs with block size 3

Each block in a packed BTD( $V, 3, \Lambda$ ) has the form  $\{v_1, v_1, v_2\}$  for some distinct pair of elements  $v_1$  and  $v_2$ . Because of this correspondence between blocks and pairs of distinct elements, it is straightforward to describe completely the fully packed

$BTD(V, 3, \Lambda)$ s and their spectrum, as well as the spectrum of the simple packed  $BTD(V, 3, \Lambda)$ s that are not fully packed. The designs are interesting in that the entire spectrum of such designs is nested and can be signed.

**Theorem 8** *If a packed  $BTD(V, 3, \Lambda)$  is simple, then either  $\Lambda = 4$  and the design is fully packed, or  $\Lambda = 2$  and the design is not fully packed.*

PROOF: No block in a simple BTM can be repeated. Thus, for simple BTMs with block size three and distinct elements  $v_1$  and  $v_2$ , either the BTM contains both of the blocks  $\{v_1, v_1, v_2\}$  and  $\{v_1, v_2, v_2\}$  or exactly one of them. If both blocks are included in the BTM, it follows that  $\Lambda = 4$  and the BTM is fully packed. If only one is included, it follows that  $\Lambda = 2$  and the BTM is packed but not fully packed.  $\square$

**Theorem 9** *For every  $V \geq 2$ , there exists a unique fully packed  $BTD(V, 3, 4)$  whose incidence matrix can be signed to create a 0-BRBTM.*

PROOF: The existence and uniqueness of the designs follows from Theorem 2. Now suppose  $M$  is the incidence matrix of a fully packed BTM with element set  $\{1, 2, \dots, V\}$ . Change a one in  $M$  to a negative one if and only if the one corresponds to  $v_2$  in the block  $\{v_1, v_1, v_2\}$  and  $v_2 < v_1$ . Since the design also contains the block  $\{v_2, v_2, v_1\}$ , the result follows.  $\square$

The blocks of a simple  $BTD(7, 3, 2)$  that is packed but not fully packed are shown in Figure 5. The design was created using a cyclic construction method. The blocks shown in column 1 of the example are called the starter blocks of the design. Each row of blocks is then cyclically generated from the starter block of the row by repeatedly adding one (mod 7) to the block elements.

001	112	223	334	445	556	660
002	113	224	335	446	550	661
003	114	225	336	440	551	662

Figure 5: A simple packed  $BTD(7,3,2)$  that is not fully packed.

**Theorem 10** *For every odd  $V = 2t + 1 \geq 3$ , there exists a simple packed  $BTD(2t + 1, 3, 2)$  which is not fully packed.*

PROOF: We use the cyclic construction method illustrated above to construct a  $BTD(2t + 1, 3, 2)$ . If  $V = 2t + 1$ , then the starter blocks are  $\{0, 0, i\}$ ,  $i \in \{1, \dots, t\}$ .  $\square$

**Theorem 11** *If  $D = BTD(2t, 3, \Lambda)$  is a simple packed BTM, then  $D$  is fully packed and  $\Lambda = 4$ .*

PROOF: Assume the design  $D = BTD(2t, 3, \Lambda)$  is packed but not fully packed. Since  $D$  is not fully packed, Theorem 8 tells us that  $\Lambda = 2$ . Further,  $B = \binom{V}{2}$  since every pair of distinct elements appear together in exactly one block. Thus,  $VR = BK$ , which implies  $VR = 3B = 3V(V - 1)/2$  which yields  $2R = 3(V - 1)$ . The left hand side of this last equation is even while the right hand side is odd. The contradiction implies the BTM must have been fully packed.  $\square$

Reducing all twos to ones in the incidence matrix of any simple packed  $\text{BTD}(V, 3, 2)$  that is not fully packed produces a matrix that represents the set of all pairs of distinct elements in the design. Similarly, reducing all twos to ones in the incidence matrix of a fully packed  $\text{BTD}(V, 3, 4)$  represents two copies of this same set. These facts can be formally stated as:

**Theorem 12** *If a simple  $\text{BTD}(V, 3, 2)$  is packed but not fully packed, then the  $\text{BTD}$  is nested. The  $\text{BTD}$  reduces to a one-factor (block size 2) and the remainders, i.e. the block elements not appearing in the reduction, form a list of the elements of the design (a one-design). If the  $\text{BTD}$  is fully packed, it reduces to two copies of the one-factor. In particular, for every admissible  $\Lambda$ , there exists a packed  $\text{BTD}(V, 3, \Lambda)$  that is nested.*

#### 4 Analysis of packed $\text{BTD}$ s with block size 5

In this section, we analyze packed  $\text{BTD}(V, 5, \Lambda)$ s. We begin our analysis by signing the fully packed  $\text{BTD}(V, 5, \Lambda)$ s. The remainder of the section focuses on the spectrum, signings and nestings of packed  $\text{BTD}(V, 5, \Lambda)$ s that are not fully packed.

**Theorem 13** *For every  $V \geq 3$  and  $\Lambda = 8(V - 2)$ , there exists a unique fully packed  $\text{BTD}(V, 5, \Lambda)$  whose incidence matrix can be signed to create a 0- $\text{BRBTD}$ .*

**PROOF:** The existence and uniqueness of the designs follows from Theorem 2. Let  $M$  be the incidence matrix of a fully packed  $\text{BTD}(V, 5, \Lambda)$  and let  $\widehat{M}$  be the matrix obtained by changing every one in  $M$  to a negative one. If  $v_1$  and  $v_2$  are any two distinct elements of the design, they appear together in  $V - 2$  blocks of the form  $\{v_1, v_1, v_2, v_2, v_i\}$ ,  $V - 2$  blocks of the form  $\{v_1, v_1, v_2, v_i, v_i\}$ , and  $V - 2$  blocks of the form  $\{v_1, v_2, v_2, v_i, v_i\}$ . Thus, the inner product of row  $v_1$  and  $v_2$  in  $\widehat{M}$  will be  $4(V - 2) - 2(V - 2) - 2(V - 2)$  or zero.  $\square$

We are now ready to prove the main result of the section. This requires a structure lemma of H. Agrawal's (stated as Lemma 3) on binary equi-replicate designs. A binary equi-replicate design is a collection of  $b$  size  $k$  blocks (i.e. sets) over a  $v$ -set of elements such that each element appears in  $r$  blocks.

**Lemma 3** [5] *The elements of every binary equi-replicate design with  $bk = vr$  and  $b = mv$ , can be arranged in a  $k$ -by- $b$  array such that each column represents a block of the design and each row contains  $m$  copies of every element.*

**Theorem 14** *Every  $\text{BIBD}(V, 3, 3s)$  is the reduction of a packed  $\text{BTD}(V, 5, 8s)$ . Moreover, the incidence matrix of the  $\text{BTD}$  can be signed to create a 0- $\text{BRBTD}$ .*

**PROOF:** Assume  $D_1$  is a  $\text{BIBD}(V, B, R, 3, 3s)$ , (recall this implies by the necessary conditions of  $\text{BIBD}$ s that  $R = 3s(V - 1)/2$  and  $B = sV(V - 1)/2$ ). For each block  $b_{1i} = \{v_1, v_2, v_3\}$  in  $D_1$ , set  $b_{2i} = \{v_1v_2, v_1v_3, v_2v_3\}$ , the set of unordered pairs of  $b_{1i}$ . Call the collection of all such blocks  $D_2$ . Now  $D_2$  contains  $B$  size three blocks

over  $V_2 = \binom{V}{2}$  elements, namely the unordered pairs of  $D_1$ , each of which appears  $3s$  times in  $D_2$ . Since the replication number of  $D_2$  is the index  $\lambda = 3s$  of  $D_1$ , it follows that  $D_2$  is an equi-replicate design with  $m = s$ . Thus, we can use Agrawal's lemma to get a 3-by- $b$  array,  $M$ , in which every pair of  $D_1$  elements appears  $s$  times in each row, and each column represents a block of  $D_2$ .

For each column  $i$  of  $M$ , there exist three distinct elements  $v_1, v_2$  and  $v_3$  of  $D_1$  such that  $v_1v_2$  is the element in position  $(1, i)$  of  $M$ ,  $v_1v_3$  is in position  $(2, i)$ , and  $v_2v_3$  is in position  $(3, i)$ . Define  $b_{3i} = \{v_1, v_2, v_3\}$ . Call the collection of all such blocks  $D_3$ .  $D_3$  contains  $B$  size five blocks over the  $V$  elements of  $D_1$ . When pair  $v_1v_2$  appears in column  $i$  and in either row one or two of  $M$ ,  $b_{3i}$  will contain one of them singly and the other doubly. When the pair appears in row three,  $b_{3i}$  will contain both doubly. Thus, pair  $v_1v_2$  appears together  $\Lambda = 2 \cdot 2s + 4s = 8s$  times in  $D_3$ . The only time  $v_1$  appears doubly in  $b_{3j}$  is if there exists  $x$  such that  $v_1x$  is the element in position  $(3, i)$  of  $M$ . Since  $v_1$  is in  $V - 1$  pairs of elements and since each pair appears  $s$  times per row of  $M$ , we can conclude  $\rho_2$  is constant and has value  $s(V - 1)$ . Let  $\rho_1(x)$  represent the number of times  $x$  appears singly in a block of  $D_3$ . If  $x$  appears singly in block  $b_{3i}$ , then  $x$  appears in 4 pairs in  $b_{3i}$ . If  $x$  appears doubly in block  $b_{3i}$ , then  $x$  appears in 6 pairs in  $b_{3i}$ . Hence, each element  $x$  of  $D_3$  appears in  $4\rho_1(x) + 6\rho_2$  element pairs of  $D_3$ . But from above we know each element of  $D_3$  appears in  $8s(V - 1)$  pairs of  $D_3$ . Thus,  $4\rho_1(x) + 6\rho_2 = 8s(V - 1)$ , which implies  $\rho_1$  is constant,  $s(V - 1)/2$ , regardless of  $x$ . It now follows that  $D_3$  is a packed BTB( $V, 5, 8s$ ).

In the incidence matrix  $M_3$  of  $D_3$ , change each one to a negative one. From the construction of the blocks, one can see that the inner product of any two rows of  $\widehat{M}_3$  is zero. Thus,  $\widehat{M}_3$  is a 0-BRBTB. If the 0-BRBTB is reduced, with the minus sign being retained, the result is a  $c$ -Bhaskar Rao BIBD( $v, 3, 3s$ ), i.e., a  $c$ -BRD( $v, 3, 3s$ ) with  $c = -s$ . □

Theorem 14 uses a BIBD to build an equi-replicate design. Then adhering to the method prescribed by Agrawal in [5], an array on the equi-replicate design elements is generated. Lastly, the original BIBD blocks and information from the Agrawal array are used to construct packed BTB blocks. We illustrate Theorem 14 by building a packed BTB(5, 5, 8) from a BIBD(5, 3, 3) on elements  $\{0, 1, 2, 3, 4\}$ . The ten blocks of the BIBD(5, 3, 3) are given in row 2 of Figure 6. Note in order to save space, block  $\{a, b, c\}$  has been written as  $abc$ . The three 2-sets produced by each BIBD block are used to form the blocks of an equi-replicate design, and an Agrawal array is generated on the equi-replicate design elements. One such Agrawal array is shown in rows 4–6 of Figure 6. It is important to note that Agrawal's construction guarantees an array such that each equi-replicate element appears exactly once in each row of the array and such that each column represents a block of the equi-replicate design. Because of this guarantee, we can use any row of the Agrawal array to tell us which two elements of the BIBD blocks should be doubletons in the BTB blocks; we use row six. The blocks of the BTB are shown in row 8 of Figure 6.

It is straightforward to build the incidence matrix of the BTB shown in row 8 of Figure 6. Changing each 1 in the matrix to a  $-1$  produces a 0-BRBTB.

The blocks of the BIBD									
013	034	014	023	012	134	123	124	024	234
The Agrawal array for the equi-replicate design									
{0,1}	{0,3}	{1,4}	{0,2}	{1,2}	{3,4}	{1,3}	{2,4}	{0,4}	{2,3}
{0,3}	{0,4}	{0,1}	{2,3}	{0,2}	{1,3}	{1,2}	{1,4}	{2,4}	{3,4}
{1,3}	{3,4}	{0,4}	{0,3}	{0,1}	{1,4}	{2,3}	{1,2}	{0,2}	{2,4}
The blocks of the BTD									
11330	33440	00441	00332	00112	11443	22331	11224	00224	22443

Figure 6: An illustration of the sequence of structures Lemma 6 uses to construct packed BTDs.

In the above theorem, the construction implies that the original BIBD is a reduction of the constructed BTD; the construction also implies that the BTD will be simple whenever the BIBD is. This allows us to describe the  $\Lambda$ -spectrum of such designs.

**Lemma 4** *If a packed BTD( $V, 5, \Lambda$ ) reduces to a BIBD( $V, 3, \lambda$ ), then  $\Lambda = 8d_4$  and  $\lambda = 3d_4$ , where  $d_4$  is as in Theorem 7.*

PROOF: From Lemma 2 we know  $\Lambda = 4\rho_1 \cdot 4/(V-1)$  and  $\rho_2 = 2\rho_1$ . From the necessary conditions for BIBDs we know  $\lambda = 2r/(V-1)$ . Thus,  $6\rho_1/(V-1) = d_2 + d_4$  and  $16\rho_1/(V-1) = 2d_2 + 4d_4$ . Solving these as a system of equations yields  $d_2 = 2d_4$  and the result follows.  $\square$

**Lemma 5** *If a BTD( $V, 5, \Lambda$ ) is packed, then  $\Lambda \geq 4$ . Further, if the BTD reduces to a BIBD( $V, 3, \lambda$ ), then  $\Lambda \equiv 0 \pmod{8}$ ,  $\lambda \equiv 0 \pmod{3}$ , and  $\lambda = 3\Lambda/8$ .*

PROOF: Every block of a packed BTD( $V, 5, \Lambda$ ) has form  $\{v_1, v_1, v_2, v_2, v_3\}$  for distinct  $v_1, v_2$  and  $v_3$ . Since pair  $v_1, v_2$  appears four times in the block, it follows that  $\Lambda \geq 4$ . Lemma 4 tells us  $\lambda = 3d_4$  and  $\Lambda = 8d_4$ . Thus,  $\Lambda \equiv 0 \pmod{8}$ ,  $\lambda \equiv 0 \pmod{3}$  and  $\lambda = 3\Lambda/8$ .  $\square$

**Theorem 15** (a) *If  $V$  is odd, then the  $\Lambda$ -spectrum of the simple packed BTD( $V, 5, \Lambda$ )s that reduce to a simple BIBD( $V, 3, \lambda$ ) is the set  $\{\Lambda = 8s : 8 \leq 8s \leq 8(V-2)/3\}$ .*

(b) *If  $V$  is even, then the  $\Lambda$ -spectrum of the simple packed BTD( $V, 5, \Lambda$ )s that reduce to a simple BIBD( $V, 3, \lambda$ ) is the set  $\{\Lambda = 16s : 16 \leq 16s \leq 8(V-2)/3\}$ .*

PROOF: It is well-known that simple BIBD( $V, 3, \lambda$ )s exist if and only if (1)  $\lambda$  and  $V$  satisfy the necessary conditions for a BIBD and (2)  $\lambda \leq V-2$  (see p.85 of [11]). When  $V$  is odd this implies simple BIBD( $V, 3, 3s$ )s exist for  $3s \leq V-2$ . When  $V$  is even this implies simple BIBD( $V, 3, 3s$ )s exist for  $s$  even and  $3s \leq V-2$ . This together with Theorem 14, and Lemma 5 yields the result.  $\square$

**Corollary 3** *If  $V$  is even, then the  $\Lambda$ -spectrum of the simple packed  $BTD(V, 5, \Lambda)$ s is the same as the spectrum of the simple packed  $BTD(V, 5, \Lambda)$ s that reduces to a simple  $BIBD(V, 3, \lambda)$ , namely the set  $\{\Lambda = 16s : 16 \leq 16s \leq 8(V - 2)/3\}$ .*

PROOF: *If  $V$  is even, then the index  $\Lambda$  must be a multiple of 16 since Lemma 2 (d) tells us that  $\Lambda(V - 1) = 16\rho_1$  for any packed  $BTD(V, 5, \Lambda)$ . But the  $\Lambda$ -spectrum of the simple nested packed  $BTD$ s already covers all such values.  $\square$*

**Corollary 4** *The necessary conditions are sufficient for the existence of a simple packed 0-BRBTDD( $V, 5, \Lambda$ ) which reduces to a  $BIBD(V, 3, \lambda)$ .*

Although Theorem 14 constructed simple packed  $BTD(V, 5, \Lambda)$ s that reduced to simple  $BIBD(V, 3, \lambda)$ s, there is no guarantee that every simple packed  $BTD(V, 5, \Lambda)$  does so; in fact there is no guarantee that they even reduce to  $BIBD$ s. Below we show two constructions that illustrate these facts. The first construction is adapted from page 135 of [6].

**Theorem 16** *If a  $BIBD(V, 3, 1)$  exists, then a simple packed  $BTD(V, 5, 8)$  exists that reduces to a  $BIBD(V, 3, 3)$  that is not simple. Further, the  $BTD(V, 5, 8)$  can be signed to produce a 0-BRBTDD( $V, 5, 8$ ).*

PROOF: Let  $D_1$  be a simple  $BIBD(V, 3, 1)$  with incidence matrix  $M_1$ , and let  $\widehat{M}_2$  be the 0-BRBTDD(3, 5, 8) shown in Figure 3. (Note  $M_2$  is a simple  $BTD$  that reduces to a  $BIBD$  that is not simple.) We build a new matrix  $\widehat{M}_3$  by replacing each element of  $M_1$  with a triple of elements. In particular, replace the three ones in each column of  $M_1$  with the three different rows of  $\widehat{M}_2$  and replace each zero in  $M$  with a triple of three zeros. Since  $M_1$  is the incidence matrix of a simple  $BIBD(V, 3, 1)$  and  $\widehat{M}_2$  a simple 0-BRBTDD(3, 5, 8),  $\widehat{M}_3$  will be a simple 0-BRBTDD( $V, 5, 8$ ). However, each of the columns of  $M_3$  generated from a single column of  $M_1$  will reduce to the same block.  $\square$

Note a  $BIBD(V, 3, 1)$  exists if and only if  $V \equiv 1, 3 \pmod{6}$  [11].

Let  $D_1$  be the packed  $BTD(5, 5, 4)$  with blocks 11225, 44551, 11334, 33552, and 22443. The blocks of  $D_1$  reduce to the blocks 125, 451, 134, 352, and 243. Since pair 12 appears once in the reduced blocks while pair 15 appears twice,  $D_1$  is an example of a packed  $BTD(5, 5, 4)$  that does not reduce to a  $BIBD$ . We use  $D_1$  to build a class of  $BTD(V, 5, 4)$ s that do not reduce to  $BIBD$ s.

**Theorem 17** *If a  $BIBD(V, 5, \lambda)$  exists, then a packed  $BTD(V, 5, 4\lambda)$  which does not reduce to a  $BIBD$  exists. If  $v \equiv 1, 5 \pmod{20}$ , then  $\lambda$  can be taken to be one.*

PROOF: Let  $D_1$  be the  $BTD(5, 5, 4)$  described above and let  $D_2$  be a  $BIBD(V, 5, \lambda)$  that contains distinct elements  $v_1, v_2$ , and  $v_3$ . For every block  $b$  in  $D_2$  generate a copy of  $D_1$  using the elements of  $b$ . In the case where one or more of  $v_1, v_2$ , or  $v_3$  is in  $b$ , identify  $v_1$  with 1,  $v_2$  with 2, and  $v_3$  with 3. The collection of generated blocks will be a packed  $BTD(V, 5, 4\lambda)$ . When the blocks are reduced,  $v_2$  and  $v_3$  will appear together in the reduced blocks twice as many times as  $v_1$  and  $v_2$ .  $\square$

We continue our analysis of the five case by examining packed BTDs that do not reduce to BIBDs. Here a different type of design, a partially balanced incomplete block design, arises. Let  $X$  be a  $v$ -set with a symmetric  $m$ -association scheme defined on it, then a partially balanced incomplete block design with  $m$  associate classes (PBIBD( $M$ )) having parameters  $(v; b; r; \lambda_1, \lambda_1, \dots, \lambda_m)$  is a design based on the  $v$ -set  $X$  with  $b$  blocks each of size  $k$  such that (1) every element occurs in  $r$  blocks, and (2) each pair of  $i$ -th associates occur together in  $\lambda_i$  blocks; see [23].

If  $v$  is an odd integer and  $D$  is an additively generated collection of  $b$  blocks of size  $k$  over  $Z_v$  with  $s$  starter blocks, then  $D$  is a PBIBD( $(V-1)/2$ ) with  $r = 3s$ . The  $(v-1)/2$  associate classes are defined by the non-zero differences less than  $(v-1)/2$  (i.e.,  $x$  and  $y$  are  $i$ -th associates iff  $x - y = \pm i \pmod{V}$ ,  $1 \leq i \leq (V-1)/2$ ), and each  $\lambda_i$  is defined by the number of times  $i$  appears as a difference in the starter blocks. For example, in the additively generated design over  $Z_9$ , with starter blocks  $\{0, 1, 4\}$  and  $\{0, 2, 3\}$  there are 4 associate classes and  $\lambda_1 = \lambda_3 = 2$  and  $\lambda_2 = \lambda_4 = 1$ .

As seen above the indices corresponding to the various associate classes in a PBIBD( $m$ ) do not have to be unique. We are interested in PBIBDs where there are exactly two unique  $\lambda_i$ s. The above example is one such situation. In the example each element  $x$  appears twice with  $x+1$ ,  $x+3$ ,  $x+6$ , and  $x+8$  (the addition is mod 9) and once with  $x+2$ ,  $x+4$ ,  $x+5$ ,  $x+7$ .

If we extend the starter blocks from this example to  $\{0, 0, 1, 4, 4\}$  and  $\{0, 0, 3, 2, 2\}$ , it is straightforward to check that they additively generate a packed BTD(9, 5, 4) which reduces to the example PBIBD.

The following two theorems use the idea illustrated above. That is, in the proof of each theorem a PBIBD with two unique  $\lambda_i$ s is generated and then extended to a BTD that reduces back to the PBIBD. In the two proofs we also make use of the fact that a design additively generated from starter blocks, where each element appears zero, one or two times, will be a BTD iff each non-zero difference appears the same number of times.

**Theorem 18** *For every  $V \equiv 1 \pmod{4}$ , there exists a cyclic nested packed BTD( $V, 5, 4$ ) which reduces to a cyclic PBIBD.*

PROOF: Let  $V = 4t + 1$ . Additively generate the blocks of a PBIBD over  $Z_V$  using the  $t$  starter blocks  $\{0, 1, 2t\}$ ,  $\{0, 2, 2t-1\}$ ,  $\dots$ ,  $\{0, t, t+1\}$ . Calculating and counting the differences of the elements in the starter blocks shows that each of  $1, 3, \dots, 2t-1$  and their negatives mod  $v$  will be a difference two times, while each of  $2, 4, \dots, 2t$  and their negatives mod  $v$  will be a difference one time. Now extend each starter block to five elements by repeating zero and the non-zero even element in the block. The construction of the extended starter blocks increases the number of times the differences  $1, 3, \dots, 2t-1$  appear by one, and the number of times the differences  $2, 4, \dots, 2t$  appear by two. Thus the design additively generated by the extended starter blocks is a BTD( $V, 5, 4$ ). Clearly it reduces to the original PBIBD.

If we delete the element 0 from each of the starter blocks of the PBIBD used above, and additively generate these new blocks over  $Z_V$ , we get an equi-replicate design with each block having size two and  $r = 2t$  ( $r = 2t$  since there are  $t$  starter blocks

with two elements per block). Further each element  $x$  appears once with element  $y$  if and only if  $y = x \pm 2$ ,  $i = 1, 2, \dots, t$ . Thus this design is a PBIBD and it follows that the BTB is nested.  $\square$

**Theorem 19** *For every  $V \equiv 1(\text{mod } 8)$ , there exists a cyclic nested packed BTB  $(V, 5, 6)$  which reduces to a cyclic PBIBD.*

PROOF: Let  $V = 8t + 1$ . Additively generate the blocks of a PBIBD over  $Z_V$  using the  $3t$  starter blocks:  $\{0, i, t + 4i - 2\}$ ,  $\{0, -i, t + 2i - 1\}$  and  $\{0, i, t + 4i\}$ ,  $i = 1, 2, \dots, t$ . Calculating and counting the differences of the elements in the starter blocks confirms that each of 1 through  $t$  is a difference three times, while each of  $t + 1$  through  $4t$  is a difference two times (so the two indices are 3 and 2). Now extend each starter block by repeating the two non-zero elements in the block. The construction of the extended starter blocks increases the number of times the differences  $1, 2, \dots, t$  appear by three, and the number of times the differences  $t + 1, t + 2, \dots, 4t$  appear by four. Thus the design additively generated by the extended starter blocks is a BTB  $(V, 5, 6)$ . Clearly it reduces to the original PBIBD.

If we delete the element 0 from each of the  $3t$  starter blocks of the PBIBD used above, and additively generate these new blocks over  $Z_V$ , we get an equi-replicate design with each block having size two and  $r = 6t$  ( $r = 6t$  since there are  $3t$  starter blocks with two elements per block). Further each element  $x$  fails to appear with  $y$  if  $y = x \pm 1$ , otherwise  $x$  appears exactly once with  $y$ . Thus this design is a PBIBD and it follows that the BTB is nested.  $\square$

Note the PBIBD families used in the proofs of the last two theorems do not appear in Clatworthy [9] since the PBIBD association scheme, in each of our examples, has  $(v - 1)/2$  associate classes; whereas the Clatworthy designs have two associate classes

We close with a general result about the existence of packed BTB.

**Lemma 6** *If a packed BTB  $(V, 5, \Lambda)$  exists, then  $\Lambda = 2s \geq 4$ . Further, (1)  $V = 2t$  implies  $\Lambda = 16s$ , (2)  $V = 4t + 3$  implies  $\Lambda = 8s$ , and (3)  $V = 8t + 5$  implies  $\Lambda = 4s$ .*

PROOF: Lemmas 1 and 5 imply  $\Lambda = 2s \geq 4$ . Since  $K = 5$ , Lemma 2 (d) tells us that  $\Lambda(V - 1) = 16\rho_1$ . Thus, when  $V$  is even,  $V - 1$  is odd which implies  $\Lambda = 16s$ . When  $V = 4t + 3$ ,  $V - 1$  is twice an odd number which implies  $\Lambda = 8s$ . When  $V = 8t + 5$ ,  $V - 1$  is four times an odd number which implies  $\Lambda = 4s$ .  $\square$

**Theorem 20** *The necessary conditions are sufficient for the existence of a packed (and/or nested) BTB  $(V, 5, \Lambda)$ .*

PROOF: We follow the cases outlined in Lemma 6. (1) Let  $V = 2t$ . Then since there is a BIBD  $(2t, 3, 6)$ , applying the construction of Lemma 14 produces a BTB  $(2t, 5, 16)$ . (2) Similarly if  $V = 4t + 3$ , since there is a BIBD  $(V, 3, 3)$  applying the construction of Lemma 14 produces a BTB  $(4t + 3, 5, 8)$ . (3) If  $V = 8t + 5$ , the construction in Theorem 18 produces a BTB  $(8t + 5, 5, 4)$ . Multiples of these indices may be obtained from multiple copies of these designs. If  $V = 8t + 1$ , every even value of  $\Lambda$  greater than two may be obtained from a linear combination of 4 and 6 using the constructions in Theorems 18 and 19.  $\square$

## Acknowledgment

The authors would like to thank the referee for the helpful suggestions.

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(Received 9 May 2006; revised 12 Dec 2007)