#### A Formal Study of Model Inversion Attacks

Xi Wu

xiwu@cs.wisc.edu

Joint work with Matt Fredrikson, Somesh Jha and Jeffrey F. Naughton

November 9, 2016

Model Inversion Attacks

- Model Inversion Attacks
  - A kind of privacy attacks which try to "back out" sensitive data.

- Model Inversion Attacks
  - A kind of privacy attacks which try to "back out" sensitive data.

· Main Results to Discuss

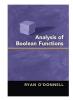
- Model Inversion Attacks
  - A kind of privacy attacks which try to "back out" sensitive data.

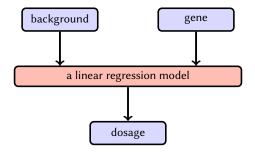
- Main Results to Discuss
  - The connection between model inversion and Boolean analysis.

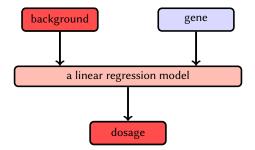
- Model Inversion Attacks
  - A kind of privacy attacks which try to "back out" sensitive data.

- Main Results to Discuss
  - The connection between model inversion and Boolean analysis.
  - · Found major applications in complexity theory.

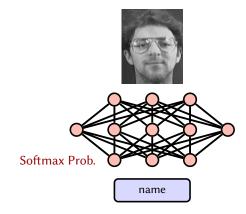




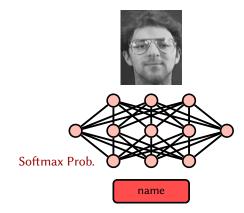




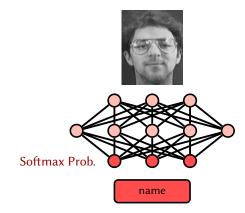
• Going from dosage and background to the genetic marker.



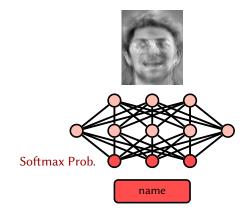
· Recover image from name:



· Recover image from name: Not good if one only knows the name..



· Recovery is sensible if softmax probabilities are known.



· Recovery is sensible if softmax probabilities are known.

Actually, many previous attempts. (not necessarily under "data privacy.")

Actually, many previous attempts. (not necessarily under "data privacy.")

- Inverting feedforward neural networks using linear and nonlinear programming. Lu et al., 1999
- *Image Reconstruction from Bag-of-Visual-Words* Kato and Harada, CVPR 2014.
- *Image reconstruction based on local feature descriptors* Maryam Daneshi, JQ Guo, 2011
- From Bits to Images: Inversion of Local Binary Descriptors d'Angelo et al. arXiv 2012

Actually, many previous attempts. (not necessarily under "data privacy.")

- Inverting feedforward neural networks using linear and nonlinear programming. Lu et al., 1999
- *Image Reconstruction from Bag-of-Visual-Words* Kato and Harada, CVPR 2014.
- *Image reconstruction based on local feature descriptors* Maryam Daneshi, JQ Guo, 2011
- From Bits to Images: Inversion of Local Binary Descriptors d'Angelo et al. arXiv 2012
- Essence: Sensible recovery from highly compressed information.



• Black-box model inversion attacks for Boolean models.

- Black-box model inversion attacks for Boolean models.
- · Formulate noiseless and noisy models and study their "invertibility."

- Black-box model inversion attacks for Boolean models.
- · Formulate noiseless and noisy models and study their "invertibility."
- Connect invertibility to notions in Boolean analysis.

- The general framework to study model inversion attacks.
  - E.g. framework to model white-box attacks.
  - Section 2.

- The general framework to study model inversion attacks.
  - E.g. framework to model white-box attacks.
  - Section 2.
- Special structure of machine learning models in white-box attacks.
  - · Sequential compositions in a model as "communication games."
  - Section 5.A.

- The general framework to study model inversion attacks.
  - E.g. framework to model white-box attacks.
  - Section 2.
- Special structure of machine learning models in white-box attacks.
  - · Sequential compositions in a model as "communication games."
  - Section 5.A.
- · Computational power of restricted communication games.
  - · Very limited communication channel can leak "everything."
  - Section 5.B.

- The general framework to study model inversion attacks.
  - E.g. framework to model white-box attacks.
  - Section 2.
- Special structure of machine learning models in white-box attacks.
  - · Sequential compositions in a model as "communication games."
  - Section 5.A.
- · Computational power of restricted communication games.
  - · Very limited communication channel can leak "everything."
  - Section 5.B.

Please refer to the paper.

## Boolean Analysis (1/2)

- Studies Boolean functions  $f : \{-1, 1\}^n \mapsto \{-1, 1\}$ .
  - $b \in \{0,1\} \mapsto (-1)^b$ .
  - Found many applications in theoretical computer science (circuit complexity, learning theory, cryptography, ...).

## Boolean Analysis (1/2)

- Studies Boolean functions  $f : \{-1, 1\}^n \mapsto \{-1, 1\}$ .
  - $b \in \{0,1\} \mapsto (-1)^b$ .
  - Found many applications in theoretical computer science (circuit complexity, learning theory, cryptography, ...).

#### Definition (Difference Operator)

 $\mathsf{D}_i$  is a linear operator applied to a Boolean function f such that  $(\mathsf{D}_i\,f)(x) = \frac{f(x^{i\to 1}) - f(x^{i\to -1})}{2}.$ 

Intuition: Discrete "derivative."

## Boolean Analysis (1/2)

- Studies Boolean functions  $f : \{-1, 1\}^n \mapsto \{-1, 1\}$ .
  - $b \in \{0,1\} \mapsto (-1)^b$ .
  - Found many applications in theoretical computer science (circuit complexity, learning theory, cryptography, ...).

#### Definition (Difference Operator)

 $\mathsf{D}_i$  is a linear operator applied to a Boolean function f such that  $(\mathsf{D}_i\,f)(x) = \frac{f(x^{i\to 1}) - f(x^{i\to -1})}{2}.$ 

Intuition: Discrete "derivative."

Definition (Influence)

$$\mathbf{Inf}_i[f] = \Pr_{x \sim \{\text{-}1,1\}^n}[f(x^{i \to 1}) \neq f(x^{i \to -1})]$$

**Intuition**: Fraction of input that  $x_i$  has influence.

Xi Wu

# Boolean Analysis (2/2)

• 
$$N_{\rho}(x)$$
.  $\widetilde{x} \sim N_{\rho}(x)$  if  
 $\widetilde{x}_j = \begin{cases} x_j & \text{w.p. } \frac{1+\rho}{2} \\ 1-x_j & \text{w.p. } \frac{1-\rho}{2} \end{cases}$ 

### Boolean Analysis (2/2)

• 
$$N_{\rho}(x)$$
.  $\tilde{x} \sim N_{\rho}(x)$  if  
 $\tilde{x}_j = \begin{cases} x_j & \text{w.p. } \frac{1+\rho}{2} \\ 1-x_j & \text{w.p. } \frac{1-\rho}{2} \end{cases}$ 

#### Definition (Noise Stability)

$$\begin{array}{l} \operatorname{Let} -1 \leq \rho \leq 1. \; \operatorname{Stab}_{\rho}[f] = \mathbb{E}_{\substack{x \sim \{\text{-}1,1\}^n \\ y \sim N_{\rho}(x)}} \left[ f(x) f(y) \right]. \end{array}$$

Intuition: Measure the change of f under noise.

#### Boolean Analysis (2/2)

• 
$$N_{\rho}(x)$$
.  $\widetilde{x} \sim N_{\rho}(x)$  if  

$$\widetilde{x}_{j} = \begin{cases} x_{j} & \text{w.p. } \frac{1+\rho}{2} \\ 1-x_{j} & \text{w.p. } \frac{1-\rho}{2} \end{cases}$$

#### Definition (Noise Stability)

$$\begin{array}{l} \operatorname{Let} -1 \leq \rho \leq 1. \; \operatorname{Stab}_{\rho}[f] = \mathbb{E}_{\substack{x \sim \{\text{-}1,1\}^n \\ y \sim N_{\rho}(x)}} \left[ f(x) f(y) \right]. \end{array}$$

Intuition: Measure the change of f under noise.

#### Definition (Stable Influence)

Let 
$$0 \leq \rho \leq 1$$
.  $\operatorname{Inf}_{i}^{(\rho)}[f] = \operatorname{Stab}_{\rho}[\operatorname{D}_{i} f] = \mathbb{E}_{x \sim \{-1,1\}^{n}} [\operatorname{D}_{i} f(x) \operatorname{D}_{i} f(y)].$   
 $y \sim N_{\rho}(x)$ 

**Intuition**: Measure the change of influence of  $x_i$  under noise. Note: when  $\rho = 1$ , this reduces to  $lnf_i[f]$ .

Setup:



Setup:

•  $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$ 

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- +  $i\in [n]$  be the target feature to invert.

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S \subseteq \{-1,1\}^n \times \{-1,1\}$  training set used to learn f.

The MI-Attack World

The Simulated World

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

The MI-Attack World	The Simulated World
	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S$	Nature samples $(x, b_x) \sim S$

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S$	Nature samples $(x, b_x) \sim S$
Nature presents $x_{-i}$ , $y = f(x)$	Nature presents $x_{-i}$

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S$	Nature samples $(x, b_x) \sim S$
Nature presents $x_{-i}$ , $y = f(x)$	Nature presents $x_{-i}$
Adversary: $A^f(x_{-i}, y)$	Adversary: $A^*(x_{-i})$

Setup:

- $f: \{-1, 1\}^n \mapsto \{-1, 1\}.$
- $i \in [n]$  be the target feature to invert.
- +  $S\subseteq\{-1,1\}^n\times\{-1,1\}$  training set used to learn f.

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S$	Nature samples $(x, b_x) \sim S$
Nature presents $x_{-i}$ , $y = f(x)$	Nature presents $x_{-i}$
Adversary: $A^f(x_{-i}, y)$	Adversary: $A^*(x_{-i})$

$$Adv(A, A^*) = \Pr_{z \sim S}[A^f(x_{-i}, y) = x_i] - \Pr_{z \sim S}[A^*(x_{-i}) = x_i]$$

**Idea**: Measure the additional invertibility (advantage) of being able to access the model with model output.

• As  $x_i$  is uniformly random, so  $\mathbb{E}_{x \sim \{-1,1\}^n} [A^*(x_{-i}) = x_i] = \frac{1}{2}$ .

- As  $x_i$  is uniformly random, so  $\mathbb{E}_{x \sim \{-1,1\}^n}[A^*(x_{-i}) = x_i] = \frac{1}{2}$ .
- For  $\mathbb{E}_{x \sim \{-1,1\}^n}[A^f(x_{-i}, y) = x_i]$ , consider

### Noiseless is Easy

- As  $x_i$  is uniformly random, so  $\mathbb{E}_{x \sim \{-1,1\}^n} [A^*(x_{-i}) = x_i] = \frac{1}{2}$ .
- For  $\mathbb{E}_{x \sim \{-1,1\}^n}[A^f(x_{-i}, y) = x_i]$ , consider

#### Algorithm 3 Algorithm A<sub>#</sub>

Input: 
$$x_{-i}, y \in \{-1, 1\}$$
. Oracle access to  $f$ .  
1: function  $A_{\#}(x_{-i}, y)$   
2: Compute  $y' = f(x_1, ..., x_{i-1}, -1, x_{i+1}, ..., x_n)$   
3: return  $(-1)^{1[y'=y]}$ 

- The recovery is correct when  $f(x^{i \to 1}) \neq f(x^{i \to -1})$ .

- The recovery is correct when  $f(x^{i \to 1}) \neq f(x^{i \to -1})$ .
- Otherwise, no additional information is obtained.

- The recovery is correct when  $f(x^{i \to 1}) \neq f(x^{i \to -1})$ .
- Otherwise, no additional information is obtained.

• Let 
$$p = \Pr_{x \sim \{-1,1\}^n} [f(x^{i \to 1}) \neq f(x^{i \to -1})]$$
, then  

$$\Pr_{x \sim \{-1,1\}^n} [A^f_{\#}(x_{-i}, y) = x_i] = (1-p) \cdot \frac{1}{2} + p \cdot 1 = \frac{1}{2} + \frac{p}{2}$$

- The recovery is correct when  $f(x^{i \to 1}) \neq f(x^{i \to -1})$ .
- Otherwise, no additional information is obtained.

• Let 
$$p = \Pr_{x \sim \{-1,1\}^n} [f(x^{i \to 1}) \neq f(x^{i \to -1})]$$
, then  

$$\Pr_{x \sim \{-1,1\}^n} [A^f_{\#}(x_{-i}, y) = x_i] = (1-p) \cdot \frac{1}{2} + p \cdot 1 = \frac{1}{2} + \frac{p}{2}$$

Theorem

$$(\forall A^*) Adv(A_{\#}, A^*) = \frac{\ln \mathbf{f}_i[f]}{2}.$$

- The recovery is correct when  $f(x^{i \to 1}) \neq f(x^{i \to -1})$ .
- Otherwise, no additional information is obtained.

• Let 
$$p = \Pr_{x \sim \{-1,1\}^n} [f(x^{i \to 1}) \neq f(x^{i \to -1})]$$
, then  

$$\Pr_{x \sim \{-1,1\}^n} [A^f_{\#}(x_{-i}, y) = x_i] = (1-p) \cdot \frac{1}{2} + p \cdot 1 = \frac{1}{2} + \frac{p}{2}$$

#### Theorem

$$(\forall A^*) Adv(A_{\#}, A^*) = \frac{\ln \mathbf{f}_i[f]}{2}.$$

• This is in fact *optimal* given the information the adversary has.

Theorem 
$$(\forall A, \forall A^*) \ Adv(A, A^*) \leq \frac{\ln f_i[f]}{2}.$$

## Noisy Case: $\rho$ -Independent Perturbation Model

The MI-Attack World

The Simulated World

## Noisy Case: $\rho$ -Independent Perturbation Model

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$	Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{ ho}(x)$	Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$
Nature presents $\widetilde{x}_{-i}$ , $y = f(x)$	Nature presents $\widetilde{x}_{-i}$

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x,b_x)\sim S, \widetilde{x}\sim N_ ho(x)$	Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$
Nature presents $\widetilde{x}_{-i}, \ y = f(x)$	Nature presents $\widetilde{x}_{-i}$
Adversary: $A^f(\widetilde{x}_{-i}, y)$	Adversary: $A^*(\widetilde{x}_{-i})$

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{ ho}(x)$	Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$
Nature presents $\widetilde{x}_{-i}, \ y = f(x)$	Nature presents $\widetilde{x}_{-i}$
Adversary: $A^f(\widetilde{x}_{-i}, y)$	Adversary: $A^*(\widetilde{x}_{-i})$

Key: The auxiliary information is noisy – the adversary gets  $\tilde{x}_{-i}$ .

The MI-Attack World	The Simulated World
Goal: recover $x_i$	Goal: recover $x_i$
Nature samples $(x,b_x)\sim S, \widetilde{x}\sim N_ ho(x)$	Nature samples $(x, b_x) \sim S, \widetilde{x} \sim N_{\rho}(x)$
Nature presents $\widetilde{x}_{-i}$ , $y = f(x)$	Nature presents $\widetilde{x}_{-i}$
Adversary: $A^f(\widetilde{x}_{-i}, y)$	Adversary: $A^*(\widetilde{x}_{-i})$

Key: The auxiliary information is noisy – the adversary gets  $\tilde{x}_{-i}$ .

What is model invertibility then?

• Consider the same algorithm  $A_{\#}$  again

Algorithm 4 Algorithm A<sub>#</sub>

Input: 
$$\widetilde{x}_{-i}, y \in \{-1, 1\}$$
. Oracle access to  $f$ .  
1: function  $A^{\#}(\widetilde{x}_{-i}, y)$   
2: Compute  $y' = f(\widetilde{x}_1, \dots, \widetilde{x}_{i-1}, -1, \widetilde{x}_{i+1}, \dots, \widetilde{x}_n)$   
3: return  $(-1)^{\mathbb{I}[y'=y]}$ 

Instead of receiving  $x_{-i}$ , it gets now  $\tilde{x}_{-i}$ .

• Invertibility becomes "stable influence."

- Invertibility becomes "stable influence."
- · Recall that

#### Definition (Stable Influence)

Let  $0 \le \rho \le 1$ . The  $\rho$ -stable influence of f at i, denoted as  $\operatorname{Inf}_{i}^{(\rho)}[f]$ , is defined to be  $\operatorname{Inf}_{i}^{(\rho)}[f] = \operatorname{Stab}_{\rho}[\operatorname{D}_{i} f] = \mathbb{E}_{x \sim \{-1,1\}^{n}} \left[\operatorname{D}_{i} f(x) \operatorname{D}_{i} f(y)\right].$   $y \sim N_{\rho}(x)$ 

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

1.  $D_i f(x) D_i f(\tilde{x}) > 0$ : "Good,"  $A_{\#}$  infers  $x_i$  correctly as before.

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

- 1.  $D_i f(x) D_i f(\tilde{x}) > 0$ : "Good,"  $A_{\#}$  infers  $x_i$  correctly as before.
- 2.  $D_i f(x) D_i f(\tilde{x}) = 0$ : "Random guessing," the information is "erased," and  $A_{\#}$  is "essentially" doing random guessing.

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

- 1.  $D_i f(x) D_i f(\tilde{x}) > 0$ : "Good,"  $A_{\#}$  infers  $x_i$  correctly as before.
- 2.  $D_i f(x) D_i f(\tilde{x}) = 0$ : "Random guessing," the information is "erased," and  $A_{\#}$  is "essentially" doing random guessing.
- 3.  $D_i f(x) D_i f(\tilde{x}) < 0$ : "Bad," the information is "reversed,"  $A_{\#}$  always gets it wrong!

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

- 1.  $D_i f(x) D_i f(\tilde{x}) > 0$ : "Good,"  $A_{\#}$  infers  $x_i$  correctly as before.
- 2.  $D_i f(x) D_i f(\tilde{x}) = 0$ : "Random guessing," the information is "erased," and  $A_{\#}$  is "essentially" doing random guessing.
- 3.  $D_i f(x) D_i f(\tilde{x}) < 0$ : "Bad," the information is "reversed,"  $A_{\#}$  always gets it wrong!

#### Theorem

For the same 
$$A_{\#}$$
,  $(\forall A^*)$   $Adv(A_{\#}, A^*) \leq \frac{\ln \mathbf{f}_i^{(p)}[f]}{2}$ .

Let  $\widetilde{x} \sim N_{\rho}(x)$ . For  $A_{\#}$ , intuitively, there are three cases:

- 1.  $D_i f(x) D_i f(\tilde{x}) > 0$ : "Good,"  $A_{\#}$  infers  $x_i$  correctly as before.
- 2.  $D_i f(x) D_i f(\tilde{x}) = 0$ : "Random guessing," the information is "erased," and  $A_{\#}$  is "essentially" doing random guessing.
- 3.  $D_i f(x) D_i f(\tilde{x}) < 0$ : "Bad," the information is "reversed,"  $A_{\#}$  always gets it wrong!

#### Theorem

For the same  $A_{\#}$ ,  $(\forall A^*) Adv(A_{\#}, A^*) \leq \frac{\ln \mathbf{f}_i^{(\rho)}[f]}{2}$ .

Is  $A_{\#}$  optimal (as in the noiseless case)?

- Subtlety: If  $D_i f(x) D_i f(\tilde{x}) = 0$ ,  $A_{\#}$  essentially does random guessing, but one can do better..

- Subtlety: If  $D_i f(x) D_i f(\tilde{x}) = 0$ ,  $A_{\#}$  essentially does random guessing, but one can do better..
- For example, consider  $OR_n(x_1, \ldots, x_n) = \bigvee_{i=1}^n x_i$ .
  - The value is -1 if any input bit is -1.

- Subtlety: If  $D_i f(x) D_i f(\tilde{x}) = 0$ ,  $A_{\#}$  essentially does random guessing, but one can do better..
- For example, consider  $OR_n(x_1, \ldots, x_n) = \bigvee_{i=1}^n x_i$ .
  - The value is -1 if any input bit is -1.
- If we "see" that y = 1, then x = 1.

- Subtlety: If  $D_i f(x) D_i f(\tilde{x}) = 0$ ,  $A_{\#}$  essentially does random guessing, but one can do better..
- For example, consider  $OR_n(x_1, \ldots, x_n) = \bigvee_{i=1}^n x_i$ .
  - The value is -1 if any input bit is -1.
- If we "see" that y = 1, then x = 1.
- That is, one can use the structure of the model *f* to "denoise."

- Subtlety: If  $D_i f(x) D_i f(\tilde{x}) = 0$ ,  $A_{\#}$  essentially does random guessing, but one can do better..
- For example, consider  $OR_n(x_1, \ldots, x_n) = \bigvee_{i=1}^n x_i$ .
  - The value is -1 if any input bit is -1.
- If we "see" that y = 1, then x = 1.
- That is, one can use the structure of the model *f* to "denoise."
- In fact for OR<sub>n</sub>, in noisy model one can always achieve advantage  $\frac{\ln f_i[\text{OR}_n]}{2} = 2^{-n}, \text{ while } \frac{\ln f_i^{(\rho)}[\text{OR}_n]}{2} = \rho^{n-1}2^{-n}.$

### **Open Question**

### **Open Question**

For any  $A', A^*$ ,  $Adv(A', A^*) \le Adv(A_{\#}, A^*) + o_n(1)$ ?



	Noiseless	Noisy Model
Invertibility	Influence	Stable Influence

	Noiseless	Noisy Model
Invertibility	Influence	Stable Influence

• As  $\rho \to 0$ ,  $\mathbf{Inf}_i^{(\rho)}[\mathsf{OR}_n]$  is exponentially smaller than  $\mathbf{Inf}_i[\mathsf{OR}_n]$ .

	Noiseless	Noisy Model
Invertibility	Influence	Stable Influence

- As  $\rho \to 0$ ,  $\inf_{i}^{(\rho)}[OR_n]$  is exponentially smaller than  $\inf_{i}[OR_n]$ .
- But  $Inf_i[OR_n]$  is "exponentially small:"  $2^{1-n}$ . Not very interesting...

	Noiseless	Noisy Model
Invertibility	Influence	Stable Influence

- As  $\rho \to 0$ ,  $\ln \mathbf{f}_i^{(\rho)}[OR_n]$  is exponentially smaller than  $\ln \mathbf{f}_i[OR_n]$ .
- But  $Inf_i[OR_n]$  is "exponentially small:"  $2^{1-n}$ . Not very interesting...
- A more interesting phenomenon termed "invertibility interference."

• Consider the parity function  $\chi_n(x) = \prod_{i=1}^n x_i$ .

- Consider the parity function  $\chi_n(x) = \prod_{i=1}^n x_i$ .
- $\ln \mathbf{f}_i[\chi_n] = 1 \text{most "invertible" in the noiseless model.}$

- Consider the parity function  $\chi_n(x) = \prod_{i=1}^n x_i$ .
- $\ln \mathbf{f}_i[\chi_n] = 1 \text{most "invertible" in the noiseless model.}$
- $\ln \mathbf{f}_i^{(\rho)}[\chi_n] = \rho^{n-1} \text{highly "non-invertible" in the noisy model.}$

- Consider the parity function  $\chi_n(x) = \prod_{i=1}^n x_i$ .
- $\ln f_i[\chi_n] = 1 most$  "invertible" in the noiseless model.
- $\ln \mathbf{f}_i^{(\rho)}[\chi_n] = \rho^{n-1} \text{highly "non-invertible" in the noisy model.}$
- Why? "Influential" coordinates interfere with each other to render the model "non-invertible" when little noise present.

#### Theorem

Suppose that  $h : \{-1, 1\}^n \mapsto \{-1, 1\}$  has t coordinates with influence 1. Let  $0 < \rho \leq 1$ , then for any  $i \in [n]$ ,  $\inf_i^{(\rho)}[h] \leq \rho^{t-1} \inf_i[h]$ .

#### Theorem

Suppose that  $h : \{-1, 1\}^n \mapsto \{-1, 1\}$  has t coordinates with influence 1. Let  $0 < \rho \leq 1$ , then for any  $i \in [n]$ ,  $\inf_i^{(\rho)}[h] \leq \rho^{t-1} \inf_i[h]$ .

#### **Open Question**

If, instead of having coordinates of influence 1, we are only guaranteed that individual influence is lower bounded by  $1 - \delta$  for some  $\delta > 0$ , how fast will the stable influence decay with respect to  $\delta$ ?

# Thanks!

?

