

Clarkson's Theorem

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Let $P \subseteq \mathbb{R}^2$ be a planar point set of size n . The set $S_{\leq k}$ of $(\leq k)$ -sets of P is defined as

$$S_{\leq k} := \{Q \subseteq P \mid |Q| \leq k \text{ and } Q = P \cap h, h \text{ open halfplane}\}.$$

Clarkson's theorem gives an upper bound on the number of possible $(\leq k)$ -sets.

Theorem 1. *We have $|S_{\leq k}| = O(nk)$.*

Proof. We may assume that $2 \leq k \leq n - 2$, since otherwise the theorem clearly holds.

We begin with a definition: Let $0 \leq \ell \leq k$. A pair $(p, q) \in P^2$ of distinct points in P is called ℓ -edge if and only if $|P \cap h_{\vec{pq}}^+| = \ell$. Here, $h_{\vec{pq}}^+$ denotes the open halfplane to the left of the oriented line \vec{pq} . Let $L_{\leq k}$ be the set of all $(\leq k)$ -edges.

We have $|S_{\leq k}| \leq 2|L_{\leq k}|$. We can assign to each ℓ -edge (p, q) one ℓ - and one $(\ell + 1)$ -set, namely the ℓ -set $P \cap h_{\vec{pq}}^+$, and the $(\ell + 1)$ -set that is cut off from P after slightly rotating \vec{pq} clockwise around p . Every $(\leq k)$ -set Q can be obtained this way. To see this, take a line g that bounds the halfplane defining Q . Translate g away from Q until it hits a point from P , then rotate g counterclockwise until it hits another point from P .

Let $R \subseteq P$ a random subset of P containing each point $p \in P$ independently with probability $1/k$. We consider the set $E(\text{CH}(R))$ of the edges on the convex hull of R , and we bound the size of $E(\text{CH}(R))$ in two different ways.

On the one hand, we have

$$\mathbf{E}[|E(\text{CH}(R))|] \leq \mathbf{E}[|R|] = n/k,$$

since the convex hull of R has at most $|R|$ edges, and each point from P was chosen with probability $1/k$.

Now let $(p, q) \in P^2$ be a pair of distinct points in P , and let $I_{(p,q)}$ be the indicator random variable for the event that (p, q) defines a (clockwise) edge on $\text{CH}(R)$. Then,

$$\mathbf{E}[|E(\text{CH}(R))|] = \sum_{(p,q) \in P^2} \mathbf{E}[I_{(p,q)}] \geq \sum_{(p,q) \in L_{\leq k}} \mathbf{E}[I_{(p,q)}],$$

by linearity of expectation. For a $(\leq k)$ -edge (p, q) we have that $\mathbf{E}[I_{(p,q)}]$ is precisely the probability of the event $(p, q) \in E(\text{CH}(R))$. For this event to happen, we must have (i) $p, q \in R$; and (ii) $R \cap h_{\vec{pq}}^+ = \emptyset$. The probability for this is at least $k^{-2}(1 - 1/k)^k$, since $|P \cap h_{\vec{pq}}^+| \leq k$ and since the points in R were chosen independently.

It follows that

$$\mathbf{E}[|E(\text{CH}(R))|] \geq \sum_{(p,q) \in L_{\leq k}} \mathbf{E}[I_{(p,q)}] \geq \sum_{(p,q) \in L_{\leq k}} k^{-2}(1 - 1/k)^k \geq |L_{\leq k}|/4k^2,$$

as $k \geq 2$. Hence, $|L_{\leq k}| \leq 4nk$ and $|S_{\leq k}| \leq 8nk$. □