

# Supplementary Material

## Learning Gaussian Graphical Models via Multiplicative Weights

(Anamay Chaturvedi and Jonathan Scarlett, AISTATS 2020)

All citations below are to the reference list in the main document.

### A Comparison of Runtimes

Recall that  $p$  denotes the number of nodes,  $d$  denotes the maximal degree,  $\kappa$  denotes the minimum normalized edge strength, and  $m$  denotes the number of samples. The runtimes of some existing algorithms in the literature for Gaussian graphical model selection (see Section 1.1 for an overview) are outlined as follows:

- The only algorithms with assumption-free sample complexity bounds depending only on  $(p, d, \kappa)$  have a high runtime of  $p^{O(d)}$ , namely,  $O(p^{2d+1})$  in [Misra et al., 2017], and  $O(p^{d+1})$  in [Kelner et al., 2019, Thm. 11].
- A greedy method in [Kelner et al., 2019, Thm. 7] has runtime  $O((d \log \frac{1}{\kappa})^3 mp^2)$ . The sample complexity for this algorithm is  $O(\frac{d}{\kappa^2} \cdot \log \frac{1}{\kappa} \cdot \log n)$ , but this result is restricted to attractive graphical models.
- To our knowledge,  $\ell_1$ -based methods [Cai et al., 2011, 2016, d’Aspremont et al., 2008, Meinshausen et al., 2006, Ravikumar et al., 2011, Wang et al., 2016, Yuan and Lin, 2007] such as Graphical Lasso and CLIME do not have precise time complexities stated, perhaps because this depends strongly on the optimization algorithm used. We expect that a general-purpose solver would incur  $O(p^3)$  time, and we note that [Kelner et al., 2019, Table 2] indeed suggests that these approaches are slower.
- In practice, we expect BigQUIC [Hsieh et al., 2013] to be one of the most competitive algorithms in terms of runtime, but no sample complexity bounds were given for this algorithm.
- Under the local separation condition and a walk-summability assumption, the algorithm of [Anandkumar et al., 2012] yields a runtime of  $O(p^{2+\eta})$ , where  $\eta > 0$  is an integer specifying the local separation condition.

Hence, we see that our runtime of  $O(mp^2)$  is competitive among the existing works – it is faster than other algorithms for which sample complexity bounds have been established.

### B Proof of Lemma 2 (Properties of Multivariate Gaussians)

We restate the lemma for ease of reference.

**Lemma 2.** *Given a zero-mean multivariate Gaussian  $X = (X_1, \dots, X_p)$  with inverse covariance matrix  $\Theta = [\theta_{ij}]$ , and given  $T$  independent samples  $(X^1, \dots, X^T)$  with the same distribution as  $X$ , we have the following:*

1. For any  $i \in [p]$ , we have  $X_i = \eta_i + \sum_{j \neq i} (-\frac{\theta_{ij}}{\theta_{ii}}) X_j$ , where  $\eta_i$  is a Gaussian random variable with variance  $\frac{1}{\theta_{ii}}$ , independent of all  $X_j$  for  $j \neq i$ .
2.  $\mathbb{E}[X_i | X_{\bar{i}}] = \sum_{j \neq i} (-\frac{\theta_{ij}}{\theta_{ii}}) X_j = w^i \cdot X_{\bar{i}}$ , where  $w^i = (-\frac{\theta_{ij}}{\theta_{ii}})_{j \neq i} \in \mathbb{R}^n$  (with  $n = p - 1$ ).
3. Let  $\lambda$  and  $\nu_{\max}$  be defined as in (4) and (6), set  $B := \sqrt{2 \log \frac{2pT}{\delta}}$ , and define  $(\tilde{x}^t, \tilde{y}^t) := \frac{1}{B \sqrt{\nu_{\max}(\lambda+1)}} (x^t, y^t)$ , where  $(x^t, y^t) = (X_i^t, X_{\bar{i}}^t)$  for an arbitrary fixed coordinate  $i$ . Then, with probability at least  $1 - \delta$ ,  $\tilde{y}^t$  and all entries of  $\tilde{x}^t$  ( $t = 1, \dots, T$ ) have absolute value at most  $\frac{1}{\sqrt{\lambda+1}}$ .

*Proof.* The first claim is standard in the literature (e.g., see [Zhou et al., 2011, Eq. (4)]), and the second claim follows directly from the first.

For the third claim, let  $N$  be a Gaussian random variable with mean 0 and variance 1. We make use of the standard (Chernoff) tail bound

$$\mathbb{P}(|N| > x) \leq 2e^{-x^2/2}. \quad (41)$$

By scaling the standard Gaussian distribution, recalling the definition of  $\nu_{\max}$  in (6), and using  $B = \sqrt{2 \log \frac{2pT}{\delta}}$ , it follows that

$$\mathbb{P}(|x_i^t| > \sqrt{\nu_{\max}}B) \leq \mathbb{P}\left(|N| > \sqrt{2 \log \frac{2pT}{\delta}}\right) \quad (42)$$

$$\leq 2 \exp\left(-\log \frac{2pT}{\delta}\right) \quad (43)$$

$$\leq \frac{\delta}{pT}, \quad (44)$$

and hence

$$\mathbb{P}\left(|x_i^t| > \frac{1}{\sqrt{\lambda+1}}\right) \leq \frac{\delta}{pT}. \quad (45)$$

The same high probability bound holds similarly for  $\tilde{y}^t$ . By taking the union bound over these  $p$  events, and also over  $t = 1, \dots, T$ , we obtain the desired result.  $\square$

## C Establishing Lemma 4 (Martingale Concentration Bound)

Here we provide additional details on attaining Lemma 4 from a more general result in [van de Geer, 1995]. While the latter concerns continuous-time martingales, we first state some standard definitions for discrete-time martingales. Throughout the appendix, we distinguish between discrete time and continuous time by using notation such as  $M_t, \mathcal{F}_t$  for the former, and  $\tilde{M}_t, \tilde{\mathcal{F}}_t$  for the latter.

**Definition 10.** *Given a discrete-time martingale  $\{M_t\}_{t=0,1,\dots}$  with respect to a filtration  $\{\mathcal{F}_t\}_{t=0,1,\dots}$ , we define the following:*

1. The compensator of  $\{M_t\}$  is defined to be

$$V_t = \sum_{j=1}^t \mathbb{E}[M_j - M_{j-1} | \mathcal{F}_{j-1}]. \quad (46)$$

2. A discrete-time process  $\{W_t\}_{t=1,2,\dots}$  defined on the same probability space as  $\{M_t\}$  is said to be predictable if  $W_t$  is measurable with respect to  $\mathcal{F}_{t-1}$ .
3. We say that  $\{M_t\}$  is locally square integrable if there exists a sequence of stopping times  $\{\tau_k\}_{k=1}^\infty$  with  $\tau_k \rightarrow \infty$  such that  $\mathbb{E}[M_{\tau_k}^2] < \infty$  for all  $k$ .

In the continuous-time setup of [van de Geer, 1995, Lemma 2.2], the preceding definitions are replaced by generalized notions, e.g., see [Liptser and Shiriyayev, 1989]. Note that the notion of a compensator in the continuous-time setting is much more technical, in contrast with the explicit formula (46) for discrete time.

The setup of [van de Geer, 1995] is as follows: Let  $\{\tilde{M}_t\}_{t \geq 0}$  be a locally square integrable continuous-time martingale with respect to a filtration  $\{\tilde{\mathcal{F}}_t\}_{t \geq 0}$  satisfying right-continuity ( $\tilde{\mathcal{F}}_t = \bigcap_{s>t} \tilde{\mathcal{F}}_s$ ) and completeness ( $\mathcal{F}_0$  includes all sets of null probability). For each  $t > 0$ , the martingale jump is defined as  $\Delta \tilde{M}_t = \tilde{M}_t - \tilde{M}_{t-}$ , where  $t_-$  represents an infinitesimal time instant prior to  $t$ . For each integer  $m \geq 2$ , a higher-order variation process  $\{\sum_{s \leq t} |\Delta \tilde{M}_s|^m\}$  is considered, and its compensator is denoted by  $\tilde{V}_{m,t}$ . Then, we have the following.

**Lemma 11.** [van de Geer, 1995, Lemma 2.2] *Under the preceding setup for continuous-time martingales, suppose that for all  $t \geq 0$  and some  $0 < K < \infty$ , it holds that*

$$\tilde{V}_{m,t} \leq \frac{m!}{2} K^{m-2} \tilde{R}_t, \quad m = 2, 3, \dots, \quad (47)$$

for some predictable process  $\tilde{R}_t$ . Then, for any  $a, b > 0$ , we have

$$\mathbb{P}(\tilde{M}_t \geq a \text{ and } \tilde{R}_t \leq b^2 \text{ for some } t) \leq \exp\left(-\frac{a^2}{2aK + b^2}\right). \quad (48)$$

While Lemma 11 is stated for continuous-time martingales, we obtain the discrete-time version in Lemma 4 by considering the choice  $\tilde{M}_t = M_{\lfloor t \rfloor}$ , where  $\{M_t\}_{t=0,1,\dots}$  is the discrete-time martingale. Due to the floor operation, the required right-continuity condition on the continuous-time martingale holds. Moreover, the definition of a compensator in (46) applied to the higher-order variation process with parameter  $m$  yields

$$V_{m,t} = \sum_{j=1}^t \mathbb{E}[\Delta M_j^m | \mathcal{F}_{j-1}] \quad (49)$$

with  $\Delta M_t = M_t - M_{t-1}$ , in agreement with the statement of Lemma 4. Finally, since we assumed that  $\mathbb{E}[M_t^2] < \infty$  for all  $t$  in Lemma 4, the locally square integrable condition follows by choosing the trivial sequence of stopping times,  $\tau_k = k$ .

## D Proof of Lemma 5 (Concentration of $\sum_j Z^j$ )

Lemma 5 is restated as follows.

**Lemma 5.**  $|\sum_{j=1}^T Z^j| = O\left(\sqrt{T \log \frac{1}{\delta}}\right)$  with probability at least  $1 - \delta$ .

*Proof.* Recall that  $\mathbb{E}_{t-1}[\cdot]$  denotes expectation conditioned on the history up to index  $t-1$ . Using the notation of Lemma 4, we let  $M_t = \sum_{j \leq t} Z^j$ , which yields  $\Delta M_t = Z^t$ . The definition of  $Z^t$  in (15) ensures that  $\mathbb{E}_{t-1}[Z^t] = 0$ , so that  $M_t$  is a martingale. In addition, we have

$$V_{m,t} = \sum_{j=1}^t \mathbb{E}_{j-1}[|\Delta M_j|^m] = \sum_{j=1}^t \mathbb{E}_{j-1}[|Z^j|^m]. \quad (50)$$

To use Lemma 4, we need to bound  $\sum_{j=1}^t \mathbb{E}_{j-1}[|Z^j|^m]$  for some appropriate choices of  $K$  and  $R_t$  in (12). The conditional moments of  $|Z^j|$  are the central conditional moments of  $Q^j$ :

$$\mathbb{E}_{j-1}[|Z^j|^m] = \mathbb{E}_{j-1}[|Q^j - \mathbb{E}_{j-1}[Q^j]|^m] \quad (51)$$

$$\leq \mathbb{E}_{j-1}[2^m(|Q^j|^m + |\mathbb{E}_{j-1}[Q^j]|^m)] \quad (52)$$

$$\leq 2^{m+1} \mathbb{E}_{j-1}[|Q^j|^m], \quad (53)$$

where (51) follows from the definition of  $Z^j$  in (15), (52) uses  $|a - b| \leq 2 \max\{|a|, |b|\}$ , and (53) follows from Jensen's inequality ( $|\mathbb{E}[Q^j]|^m \leq \mathbb{E}[|Q^j|^m]$ ). Furthermore, we have that

$$\mathbb{E}_{j-1}[|Q^j|^m] = \mathbb{E}_{j-1}[(\lambda p^j \cdot \tilde{x}^j - \tilde{y}^j)(p^j - w/\lambda) \cdot \tilde{x}^j]^m \quad (54)$$

$$\leq \mathbb{E}_{j-1}[|(\lambda p^j \cdot \tilde{x}^j - \tilde{y}^j)|^{2m}]^{1/2} \mathbb{E}_{j-1}[(p^j - w/\lambda) \cdot \tilde{x}^j]^{2m}]^{1/2}, \quad (55)$$

where (54) uses the definition of  $Q^j$  in (14), and (55) follows from the Cauchy-Schwartz inequality. Both of the averages in (55) contain Gaussian random variables (with  $p^j$  fixed due to the conditioning); we proceed by

establishing an upper bound on the variances. Since  $(\tilde{x}^j, \tilde{y}^j) = \frac{1}{B\sqrt{\nu_{\max}(\lambda+1)}}(x^j, y^j)$ , the definition of  $\nu_{\max}$  (see (6)) implies that each coordinate has a variance of at most  $(\frac{1}{B\sqrt{\lambda+1}})^2$ . Then, using that  $\sum_i p_i^j = 1$ , we have

$$\text{Var}(\lambda p^j \cdot \tilde{x}^j - \tilde{y}^j) \leq (\lambda + 1)^2 \max_{z \in \{\tilde{x}_1^j, \dots, \tilde{x}_n^j \tilde{y}^j\}} \text{Var}(z) \quad (56)$$

$$\leq \frac{\lambda + 1}{B^2}, \quad (57)$$

and similarly, using  $\sum_i p_i^j = 1$  and  $\|w\| = \lambda$  (see Footnote 2),

$$\text{Var}((p^j - w/\lambda) \cdot \tilde{x}^j) \leq \frac{4}{(\lambda + 1)B^2}. \quad (58)$$

Next, we use the standard fact that if  $N$  is a Gaussian random variable with mean 0 and variance  $\sigma$ , then

$$\mathbb{E}[N^p] = \begin{cases} 0 & \text{if } p \text{ is odd} \\ \sigma^p (p-1)!! & \text{if } p \text{ is even.} \end{cases} \quad (59)$$

It then follows from (53) and (57)–(59) that

$$\mathbb{E}_{j-1}[|Z^j|^m] \leq 2^{m+1} \mathbb{E}_{j-1}[|(\lambda p^j \cdot \tilde{x}^j - \tilde{y}^j)|^{2m}]^{1/2} \mathbb{E}_{j-1}[|(p^j - w/\lambda) \cdot \tilde{x}^j|^{2m}]^{1/2} \quad (60)$$

$$\leq 2^{m+1} \left( \left( \frac{\lambda + 1}{B^2} \right)^{2m} (2m-1)!! \left( \frac{4}{(\lambda + 1)B^2} \right)^{2m} (2m-1)!! \right)^{1/2} \quad (61)$$

$$= 2^{m+1} \frac{4^m}{B^{4m}} (2m-1)!! \quad (62)$$

$$= 2^{m+1} \frac{4^m}{B^{4m}} (1 \cdot 3 \cdot \dots \cdot (2m-1)) \quad (63)$$

$$\leq 2^{m+1} \frac{4^m}{B^{4m}} (2 \cdot 4 \cdot \dots \cdot 2m) \quad (64)$$

$$= 2 \cdot 4^m \frac{4^m}{B^{4m}} m! \quad (65)$$

$$= \frac{m!}{2} \left( \frac{16}{B^4} \right)^{m-2} \frac{2^{10}}{B^8}, \quad (66)$$

and summing over  $j = 1, \dots, t$  gives

$$\sum_{j=1}^t \mathbb{E}_{k-1}[|Z^j|^m] \leq \frac{m!}{2} \left( \frac{16}{B^4} \right)^{m-2} \frac{2^{10}t}{B^8}. \quad (67)$$

Hence, using the notation of Lemma 4, it suffices to set  $K = \frac{16}{B^4}$  and  $R_t = \frac{2^{10}t}{B^8}$ . Plugging everything in, we get

$$\mathbb{P} \left( \sum_{j=1}^T Z^j > a \right) < \exp \left( - \frac{a^2}{32a \frac{1}{B^4} + 2^{10} \frac{T}{B^8}} \right). \quad (68)$$

Let  $a = 2^{10} \sqrt{T \log \frac{1}{\delta}}$ . Then, since  $B = \sqrt{2 \log \frac{2pT}{\delta}}$  is always greater than  $\sqrt{\log \frac{1}{\delta}}$ , we obtain

$$\mathbb{P} \left( \sum_{j=1}^t Z^j > 2^{10} \sqrt{T \log \frac{1}{\delta}} \right) \leq \frac{\delta}{2}. \quad (69)$$

By replacing  $Z^j$  by  $-Z^j$  above, we get a symmetric lower bound on  $\sum_j Z^j$ , as all the moments used above remain the same. Applying the union bound, we get that  $|\sum_{j=1}^T Z^j| = O(\sqrt{T \log \frac{1}{\delta}})$  with probability at least  $1 - \delta$ .  $\square$

## E Proof of Lemma 7 (Concentration of Empirical Risk)

Lemma 7 is restated as follows.

**Lemma 7.** *For  $\gamma > 0$ ,  $\rho \in (0, 1]$ , and fixed  $v \in \mathbb{R}^n$  satisfying  $\|v\|_1 \leq \lambda$ , there is some  $M = O((\lambda + 1)^{\frac{\log(1/\rho)}{\gamma}})$  such that*

$$\mathbb{P}\left(\left|\frac{1}{M}\sum_{j=1}^M\left((v \cdot a^j - b^j)^2 - \Xi\right) - \varepsilon(v)\right| \geq \gamma\right) \leq \rho, \quad (32)$$

where  $\{(a^j, b^j)\}_{j=1}^M$  are the normalized samples defined in Algorithm 2, and  $\Xi = \mathbb{E}[\text{Var}[b^j | a^j]]$ .<sup>7</sup>

*Proof.* We first derive a simple equality:

$$\mathbb{E}[(v \cdot a^j - b^j)^2] = \mathbb{E}[\mathbb{E}[(v \cdot a^j - b^j)^2 | a^j]] \quad (70)$$

$$= \mathbb{E}[(\mathbb{E}[v \cdot a^j - b^j | a^j])^2 + \text{Var}[b^j | a^j]] \quad (71)$$

$$= \mathbb{E}[(v \cdot a^j - w \cdot a^j)^2] + \mathbb{E}[\text{Var}[b^j | a^j]] \quad (72)$$

$$= \varepsilon(v) + \Xi, \quad (73)$$

where (71) uses  $\text{Var}[Z] = \mathbb{E}[Z^2] - (\mathbb{E}[Z])^2$ , (72) uses the second part of Lemma 2, and (73) uses the definitions of  $\varepsilon(v)$  and  $\Xi$ .

In the following, we recall Bernstein's inequality.

**Lemma 12.** [Boucheron et al., 2013, Corollary 2.11] *Let  $Z_1, \dots, Z_n$  be independent real-valued random variables, and assume that there exist positive numbers  $\vartheta$  and  $c$  such that*

$$\sum_{i=1}^n \mathbb{E}[(Z_i)_+^2] \leq \vartheta \quad (74)$$

$$\sum_{i=1}^n \mathbb{E}[(Z_i)_+^q] \leq \frac{q!}{2} \vartheta \cdot c^{q-2}, \quad (75)$$

where  $(x)_+ = \max\{x, 0\}$ . Letting  $S = \sum_{i=1}^n (Z_i - \mathbb{E}[Z_i])$ , we have for all  $t > 0$  that

$$\mathbb{P}(S \geq t) \leq \exp\left(-\frac{t^2}{2(\vartheta + ct)}\right). \quad (76)$$

We would like to use Bernstein's inequality to bound the deviation of

$$\frac{1}{M}\sum_{j=1}^M\left((v \cdot a^j - b^j)^2 - \Xi - \varepsilon(v)\right) \quad (77)$$

from its mean value 0. To do so, we need to find constants  $\vartheta$  and  $c$  as described in the statement of Bernstein's inequality above.

Recall that  $\nu_{\max}$  upper bounds the variance of any marginal variable in each unnormalized sample, and that  $(a^j, b^j)$  are samples normalized by  $B\sqrt{\nu_{\max}(\lambda + 1)}$  with  $B = \sqrt{2 \log \frac{2pT}{\delta}} \geq 1$ . Hence, the entries of  $(a^i, b^i)$  have variance at most  $\frac{1}{\lambda + 1}$ , and since  $\|v\|_1 \leq \lambda$ , this implies that  $v \cdot a^j - b^j$  has variance at most  $\lambda + 1$ .

Using the expression for the moments of a Gaussian distribution (see (59)), it follows that

$$\mathbb{E}[(v \cdot a^j - b^j)^4] \leq 8(\lambda + 1)^2, \quad (78)$$

<sup>7</sup>This quantity is the same for all values of  $j$ .

$$\mathbb{E}[(v \cdot a^j - b^j)^{2m}] \leq (2m-1)!(\lambda+1)^m \quad (79)$$

$$\leq 2^m m! (\lambda+1)^m \quad (80)$$

$$= \frac{m!}{2} (8(\lambda+1)^2) (2(\lambda+1))^{m-2}, \quad (81)$$

where (80) is established in the same way as (65). Since  $(v \cdot a^j - b^j)^2$  is a non-negative random variable, the non-central moments bound the central moments from above. Hence, it suffices to let  $\vartheta = 8(\lambda+1)^2$  and  $c = 2(\lambda+1)$ , and we obtain from Bernstein's inequality that

$$\mathbb{P} \left( \left| \sum_{j=1}^M \left( (v \cdot a^j - b^j)^2 - \Xi - \varepsilon(v) \right) \right| \geq \gamma M \right) \leq \exp \left( \frac{-\gamma^2 M^2}{2(8(\lambda+1)^2 + 2(\lambda+1)\gamma M)} \right). \quad (82)$$

To simplify the notation, we let  $M_0$  be such that  $M = (\lambda+1)M_0$ , which yields

$$\mathbb{P} \left( \left| \frac{1}{M} \sum_{j=1}^M \left( (v \cdot a^j - b^j)^2 - \Xi - \varepsilon(v) \right) \right| \geq \gamma \right) \leq \exp \left( \frac{-\gamma^2 M_0^2}{16 + 2\gamma M_0} \right). \quad (83)$$

If  $\gamma M_0 \geq 1$ , then the right hand side is less than or equal to  $\exp\left(\frac{-\gamma M_0}{18}\right)$ . Otherwise, if  $\gamma M_0 < 1$ , then the right hand side is less than  $\exp\left(\frac{-\gamma^2 M_0^2}{18}\right)$ . It follows that to have a deviation of  $\gamma$  with probability at most  $\rho$ , it suffices to set  $M_0 = \frac{18 \log(1/\rho)}{\gamma}$ . Recalling that  $M = (\lambda+1)M_0$ , it follows that with  $M = 18(\lambda+1) \frac{\log(1/\rho)}{\gamma}$ , we attain the desired target probability  $\rho$ .  $\square$

## F Proof of Lemma 8 (Low Risk Implies an $\ell_\infty$ Bound)

Lemma 8 is restated as follows, and refers to the setup described in Section 4.

**Lemma 8.** *Under the preceding setup, if we have  $\varepsilon(v) \leq \epsilon$ , then we also have  $\|v - w\|_\infty \leq \sqrt{\epsilon \theta_{\max}}$ , where  $\theta_{\max}$  is a uniform upper bound on the diagonal entries of  $\Theta$ .*

*Proof.* Recall that  $\varepsilon(v) = \mathbb{E}[(v - w) \cdot X_{\bar{i}}]^2$ , where  $w = \left(\frac{-\theta_{ij}}{\theta_{ii}}\right)_{j \neq i}$  is the neighborhood weight vector of the node  $i$  under consideration, and  $X_{\bar{i}} = (X_j)_{j \neq i}$ . To motivate the proof, note from Lemma 2 that  $X_i = \eta_i + \sum_{j \neq i} (-\theta_{ij}/\theta_{ii}) X_j$ , where  $\eta_i$  is an  $\mathcal{N}(0, \frac{1}{\theta_{ii}})$  random variable independent of  $\{X_j\}_{j \neq i}$ , from which it follows that  $\text{Var}(X_i) \geq \text{Var}(\eta_i) = 1/\theta_{ii}$ . In the following, we apply similar ideas to  $(v - w) \cdot X_{\bar{i}}$ .

Specifically, for an arbitrary index  $i^* \neq i$ , we can lower bound the expected risk  $\varepsilon(v)$  as follows:

$$\begin{aligned} & \mathbb{E}[(v - w) \cdot X_{\bar{i}}]^2 \\ &= \text{Var}((v - w) \cdot X_{\bar{i}}) \end{aligned} \quad (84)$$

$$= \text{Var} \left( \sum_{j \neq i} (v_j - w_j) X_j \right) \quad (85)$$

$$= \text{Var} \left( (v_{i^*} - w_{i^*}) X_{i^*} + \sum_{j \notin \{i, i^*\}} (v_j - w_j) X_j \right) \quad (86)$$

$$= \text{Var} \left( (v_{i^*} - w_{i^*}) \eta_{i^*} - (v_{i^*} - w_{i^*}) \frac{\theta_{i^*i}}{\theta_{i^*i^*}} X_i + \sum_{j \notin \{i, i^*\}} \left( (v_j - w_j) - (v_{i^*} - w_{i^*}) \frac{\theta_{i^*j}}{\theta_{i^*i^*}} \right) X_j \right) \quad (87)$$

$$= \text{Var}((v_{i^*} - w_{i^*}) \eta_{i^*}) + \text{Var} \left( - (v_{i^*} - w_{i^*}) \frac{\theta_{i^*i}}{\theta_{i^*i^*}} X_i + \sum_{j \notin \{i, i^*\}} \left( (v_j - w_j) - (v_{i^*} - w_{i^*}) \frac{\theta_{i^*j}}{\theta_{i^*i^*}} \right) X_j \right) \quad (88)$$

$$\geq \text{Var}((v_{i^*} - w_{i^*}) \eta_{i^*}) \quad (89)$$

$$= |v_{i^*} - w_{i^*}|^2 \text{Var}(\eta_{i^*}), \quad (90)$$

where (84) follows since  $\mathbb{E}[X_{\bar{i}}] = 0$ , (87) follows from the first part of Lemma 2 applied to node  $i^*$ , and (88) uses the independence of  $\eta_{i^*}$  and  $X_{\bar{i}^*}$ . Since  $\text{Var}(\eta_{i^*}) = \frac{1}{\theta_{i^*i^*}}$  and  $\varepsilon(v) \leq \epsilon$ , this gives  $|v_{i^*} - w_{i^*}| \leq \sqrt{\epsilon\theta_{i^*i^*}} \leq \sqrt{\epsilon\theta_{\max}}$ . Then, since this holds for all  $i^* \neq i$ , we deduce that  $\|v - w\|_\infty \leq \sqrt{\epsilon\theta_{\max}}$ , as desired.  $\square$