## Appendix

## A. Proofs

## A.1. Proofs for Section 2

Proof of Theorem 1. Consider the distribution where there are two connected components, one for $\mathcal{X}_{+}$and one for $\mathcal{X}_{-}$, each with mixture probability $p=\frac{1}{2}$. Thus, Assumption 1 holds and we are choose to free the other parameters of the distribution in any way that satisfies Assumption 2 and 3 (e.g. a mixture of uniform density functions satisfies these assumptions). Now note that with probability $\frac{1}{2}$, the final point that is label queried by the passive learner will be positive and, thus, the passive algorithm will need to query all of the points with probability $\frac{1}{2}$ in order to retrieve all positive points. In such an event, the excess query cost is at least $\frac{1}{2} \cdot n$.

## A.2. Proofs for Section 3

Much of our technical results require the following uniform high-probability guarantee that balls of sufficient probability mass contain an example:
Lemma 3. Let $0<\delta<1$ and $\mathcal{F}$ be some distribution over $\mathbb{R}^{D}$ and $X$ be a sample of size $n$ drawn i.i.d. from $\mathcal{F}$. There exists universal constant $C_{0}$ such that the following holds with probability at least $1-\delta$ uniformly over all balls $B \in \mathbb{R}^{D}$ :

$$
\mathcal{F}(B) \geq \frac{C_{0} \cdot D \cdot \log (2 / \delta) \log n}{n} \Rightarrow|B \cap X|>0
$$

Proof. This follows by Lemma 7 of Chaudhuri \& Dasgupta (2010).
The following result bounds the volume of the $\epsilon$-neighborhood around $\mathcal{X}_{+}$, which will be used later to bound the excess number of points queried around $\mathcal{X}_{+}$. The result says that the volume of the $\epsilon$-neighboorhood around $\mathcal{X}_{+}$(and not including $\mathcal{X}_{+}$) is linear in $\epsilon$.
Lemma 4. Suppose Assumption 2 holds. Then there exists constants $r_{1}, C_{+}^{\prime}>0$ depending only on $\mathcal{X}_{+}$such that for all $0<\epsilon<r_{1}$, we have

$$
\operatorname{Vol}\left(B\left(\mathcal{X}_{+}, \epsilon\right) \backslash \mathcal{X}_{+}\right) \leq C_{+}^{\prime} \cdot \epsilon
$$

where $B\left(\mathcal{X}_{+}, \epsilon\right):=\left\{x \in \mathbb{R}^{D}: \inf _{x^{\prime} \in \mathcal{X}_{+}}\left|x-x^{\prime}\right| \leq \epsilon\right\}$.
Proof of Lemma 4. This follows from Gorin (1983). To see this, the equation on page 159 of Gorin (1983) states that if $M$ and $N$ are respectively $d$-dimensional and $(d+k)$-dimensional compact smooth Riemannian manifolds and $f: M \rightarrow N$ is a smooth isometric embedding, then we have

$$
\operatorname{Vol}(B(f(M), \epsilon))=V_{k} \cdot \epsilon^{k} \cdot \operatorname{Vol}(M)+O\left(\epsilon^{k+1}\right)
$$

where $V_{k}$ is the volume of a $k$-dimensional ball. Here, we take $M=\mathcal{X}_{+}$and $N=B\left(\mathcal{X}_{+}, r_{1}\right)$ for some $r_{1}>0$. Then, we have $k=0$ and taking $f$ to be the identity function, we have

$$
\operatorname{Vol}\left(B\left(\mathcal{X}_{+}, \epsilon\right)\right)=\operatorname{Vol}\left(\mathcal{X}_{+}\right)+O(\epsilon)
$$

and the result follows immediately.
Proof of Theorem 2. By Hoeffding's inequality, out of the initial $m$ examples that Algorithm 1 label queries, we have with probability at least $1-\delta / 2$ that at least $p-\sqrt{\frac{1}{2 m} \cdot \log (2 / \delta)}$ fraction of them are positively labeled, since the example being positive follows a Bernoulli distribution with probability $p$. Then by the condition on $m$, we have that at least $p / 2$ fraction of the points are positively labeled.
Take

$$
\epsilon=\left(\frac{2 \cdot C_{0} \cdot D \cdot \log (4 / \delta) \log (p \cdot m / 2)}{p^{2} \cdot \lambda_{0} \cdot C_{+} \cdot v_{D} \cdot m}\right)^{1 / D}, \quad M_{0}=\max \left\{\frac{2 \cdot C_{0} \cdot D(\log (p \cdot / 2)+1)}{p^{2} \cdot \lambda_{0} \cdot C_{+} \cdot v_{D} \cdot \min \left\{r_{0}, r_{1}\right\}^{D}}, 2 e\right\}
$$

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where $v_{D}$ is the volume of a unit ball in $\mathbb{R}^{D}$. Then, we have that the condition on $m$ and $M_{0}$ guarantees that $\epsilon<\min \left\{r_{0}, r_{1}\right\}$.
Let $x \in \mathcal{X}_{+}$. We have that the probability mass of positive examples in $B(x, \epsilon)$ w.r.t. $\mathcal{P}$ is:

$$
\begin{aligned}
p \cdot \mathcal{P}_{+}(B(x, \epsilon)) & \geq p \cdot \lambda_{0} \cdot \operatorname{Vol}\left(B(x, \epsilon) \cap \mathcal{X}_{+}\right) \\
& \geq p \cdot \lambda_{0} \cdot C_{+} \cdot \operatorname{Vol}(B(x, \epsilon)) \\
& \geq p \cdot \lambda_{0} \cdot C_{+} \cdot v_{D} \cdot \epsilon^{D} \\
& \geq \frac{2 \cdot C_{0} \cdot D \log (4 / \delta) \log (p \cdot m / 2)}{p \cdot m} .
\end{aligned}
$$

Then by Lemma 3, we have with probability at least $1-\delta / 2$ that all the positive examples in $X$ are within $\epsilon$ of one of the positive examples among the initially sampled $m$ examples. Therefore, Algorithm 1 retrieves all of the positive examples.

Now we bound the expected regret:

$$
\begin{aligned}
\mathbb{E}\left[C_{\text {offline }}\right] & \leq(1-p) \cdot m+n \cdot(1-p) \cdot \mathcal{P}_{-}\left(B\left(\mathcal{X}_{+}, \epsilon\right) \backslash \mathcal{X}_{+}\right) \\
& \leq(1-p) \cdot m+n \cdot(1-p) \cdot \lambda_{1} \cdot C_{+}^{\prime} \cdot \epsilon
\end{aligned}
$$

The result follows.
Proof of Theorem 3. Let $\mathcal{P}_{+}$be the uniform distribution on the unit hypercube $[0,1]^{D}$ and $\mathcal{P}_{-}$be the uniform distribution on $[-1,2]^{D}$. In the initial sampling phase of Algorithm 1, at most $m$ of the examples will be positively labeled. Let $\widehat{\mathcal{X}_{+}}=X \cap\left(\cup_{x \in X_{0,+}} B(x, \epsilon)\right)$, the set of points that Algorithm 1 labeled. Then, Theorem 3 b in (Cuevas et al., 1997) shows that for $n$ sufficiently large, with probability at least $1 / 4$, we have

$$
d_{H}\left(\widehat{\mathcal{X}_{+}}, \mathcal{X}_{+}\right) \geq \frac{1}{4}\left(\frac{\log m}{m}\right)^{1 / D}
$$

for any $\epsilon>0$, where $d_{H}(A, B):=\max \left\{\sup _{x \in A} d(x, B), \sup _{x \in B} d(x, A)\right\}$ is the Hausdorff distance. Therefore, we have (in the case of taking $\epsilon \rightarrow 0$ ):

$$
d_{H}\left(X_{0,+}, \mathcal{X}_{+}\right) \geq \frac{1}{4}\left(\frac{\log m}{m}\right)^{1 / D}
$$

Since $X_{0,+} \subseteq \mathcal{X}_{+}$, it follows that $d_{H}\left(X_{0,+}, \mathcal{X}_{+}\right)=\sup _{x \in \mathcal{X}_{+}} d\left(x, X_{0,+}\right)$. Therefore, we need $\epsilon \geq \frac{1}{4}\left(\frac{\log m}{m}\right)^{1 / D}$ in order for Algorithm 1 to recover all of the positive examples. Thus, the expected regret is at least (for some $C>0$ )

$$
\mathbb{E}\left[C_{\text {offline }}\right] \geq(1-p) \cdot m+C \cdot\left(\frac{\log m}{m}\right)^{1 / D} \cdot n
$$

as desired.

## A.3. Proofs for Section 4

Proof of Lemma 1. By Hoeffding's inequality, out of the initial $m$ examples that Algorithm 2 label queries, we have with probability at least $1-\delta / 2$ that at least $p-\sqrt{\frac{1}{2 m} \cdot \log (2 / \delta)}$ are fraction of them are positively labeled, since the example being positive follows a Bernoulli distribution with probability $p$. Then by the condition on $m$, we have that at least $p / 2$ fraction of the points are positively labeled and thus we have at least $m p / 2$ positive examples.

Then, we have that out of these $m p / 2$ examples, the probability that none of them are in $\mathcal{X}_{+, i}$ for each $i \in[c]$ is at most

$$
\left(1-\mathcal{P}_{+}\left(\mathcal{X}_{+, i}\right)\right)^{m p / 2} \leq(1-q)^{m p / 2} \leq \frac{\delta}{2 c}
$$

The result follows by union bound.

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Proof of Lemma 2. Let $x, x^{\prime} \in \mathcal{X}_{+, i}$. There exists a path $x=x_{1} \rightarrow x_{2} \rightarrow \ldots \rightarrow x_{q}=x^{\prime}$ in $\mathcal{X}_{+, i}$ such that $\left\|x_{j}-x_{j+1}\right\| \leq$ $\epsilon / 3$. We also have that the probability mass of positive examples in $B\left(x_{j}, \epsilon / 3\right)$ w.r.t. $\mathcal{P}$ is:

$$
p \cdot \mathcal{P}_{+}\left(B\left(x_{j}, \epsilon / 2\right)\right) \geq p \cdot C_{+} \cdot \lambda_{0} \cdot v_{D} \cdot(\epsilon / 3)^{D} \geq \frac{C_{0} \cdot D \log (2 / \delta) \cdot \log n}{n}
$$

Therefore by Lemma 3, there exists $x_{j}^{\prime} \in B\left(x_{j}, \epsilon / 3\right) \cap X_{+}$, where $X_{+}$are the positive examples in $X$. Hence, by triangle inequality, there exists path $x=x_{1}^{\prime} \rightarrow x_{2}^{\prime} \rightarrow \ldots \rightarrow x_{q}^{\prime}=x^{\prime}$ all in $X_{+}$where $\left\|x_{j}^{\prime}-x_{j+1}^{\prime}\right\| \leq\left\|x_{j}^{\prime}-x_{j}\right\|+\| x_{j}-$ $x_{j+1}\|+\| x_{j+1}-x_{j+1}^{\prime} \| \leq \epsilon$ implying that $\mathcal{X}_{+, i} \cap X_{+}$is connected in the $\epsilon$-neighborhood graph of $X_{+}$. The result follows immediately.

Finally, we combine these two results to show the final excess query cost guarantee for Algorithm 2.
Proof of Theorem 4. Take

$$
N_{0}=\frac{3^{D} \cdot C_{0} \cdot D}{\min \left\{r_{0}, r_{1}\right\}^{D} \cdot p \cdot C_{+} \cdot \lambda_{0} \cdot v_{D}}
$$

By Lemma 1, there exists at least one positive example in the initial $m$ samples from each connected component of $\mathcal{X}_{+}$. Define

$$
\epsilon:=3\left(\frac{C_{0} \cdot D \log (4 / \delta) \cdot \log n}{p \cdot C_{+} \cdot \lambda_{0} \cdot v_{D} \cdot n}\right)^{1 / D}
$$

We have that the condition on $n$ implies that $\epsilon \leq \min \left\{r_{0}, r_{1}\right\}$. By Lemma 2, we have that all of the positive examples of each connected component of $\mathcal{X}_{+}$are in the same CC of the $\epsilon$-neighborhood graph of the positive examples. Therefore, when the algorithm terminates, the set of examples it will select from will be contained in $B\left(\mathcal{X}_{+}, \epsilon\right)$.
Therefore, we have

$$
\mathbb{E}\left[C_{\text {exp-commit }}\right] \leq(1-p) \cdot m+C_{+}^{\prime} \cdot \epsilon \cdot n \leq(1-p) \cdot m+C \cdot\left((\log (4 / \delta) \cdot \log n)^{1 / D} \cdot n^{(D-1) / D}\right.
$$

for some $C$ depending on $\mathcal{P}$, as desired.

## B. Additional Experiment Plots

In Table 3, we show the full results for the Letters dataset for the area under curve metrics. We see that in all cases, our method outperforms outright. In Table 4, we show the full results for CelebA. We see that our method is competitive for 32 out of the 40 tasks.

| Dataset | Label | O | O-LS | A-LS | O-RS | A-RS | O-IF | A-IF | O-RC | A-RC | EC (Ours) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Letters | A | 91.14 | 52.52 | 52.49 | 84.81 | 89.02 | 59.52 | 84.69 | 64.27 | 86.87 | 97.12 |
|  | B | 83.41 | 52.73 | 52.57 | 75.95 | 82.41 | 56.13 | 75.89 | 61.38 | 80.18 | 93.76 |
|  | C | 84.48 | 56.19 | 56.08 | 75.92 | 83.78 | 55.78 | 78.68 | 59.21 | 81.92 | 94.2 |
|  | D | 83.51 | 52.57 | 52.45 | 76.14 | 82.23 | 56.09 | 76.03 | 61.51 | 78.83 | 93.73 |
|  | E | 78.6 | 52.85 | 52.99 | 74.84 | 81.41 | 55.77 | 73.7 | 60.17 | 77.27 | 89.5 |
|  | F | 83.63 | 53.4 | 53.41 | 78.72 | 83.16 | 57.44 | 77.83 | 64.69 | 80.88 | 94.0 |
|  | G | 82.23 | 52.72 | 52.67 | 75.76 | 81.88 | 58.15 | 74.58 | 63.45 | 79.79 | 92.35 |
|  | H | 69.07 | 52.4 | 52.63 | 61.19 | 69.51 | 53.95 | 64.15 | 52.03 | 66.94 | 81.78 |
|  | I | 84.26 | 52.37 | 52.59 | 73.96 | 80.72 | 55.95 | 77.56 | 59.11 | 83.83 | 93.89 |
|  | J | 84.24 | 54.45 | 54.35 | 75.16 | 81.72 | 56.22 | 77.06 | 58.58 | 81.61 | 94.28 |
|  | K | 72.7 | 52.57 | 52.62 | 65.15 | 72.03 | 55.46 | 68.12 | 51.45 | 74.38 | 89.4 |
|  | L | 80.98 | 52.59 | 52.67 | 72.74 | 79.87 | 55.44 | 77.73 | 59.16 | 78.57 | 93.18 |
|  | M | 81.83 | 54.41 | 54.25 | 77.48 | 83.23 | 59.55 | 78.61 | 60.62 | 81.67 | 93.07 |
|  | N | 75.15 | 52.67 | 52.68 | 68.58 | 75.31 | 55.37 | 69.43 | 55.85 | 73.74 | 89.8 |
|  | O | 87.71 | 52.51 | 52.62 | 80.88 | 86.08 | 57.88 | 78.08 | 67.62 | 82.71 | 94.69 |
|  | P | 84.76 | 52.7 | 52.7 | 79.24 | 84.47 | 57.63 | 79.18 | 64.81 | 83.19 | 93.47 |
|  | Q | 82.25 | 53.22 | 53.08 | 74.97 | 80.09 | 55.71 | 74.66 | 60.49 | 79.55 | 92.18 |
|  | R | 82.68 | 52.59 | 52.48 | 76.72 | 82.47 | 56.32 | 75.3 | 61.37 | 79.26 | 92.77 |
|  | S | 76.51 | 52.83 | 52.91 | 70.42 | 75.85 | 55.24 | 72.01 | 59.66 | 75.84 | 89.27 |
|  | T | 83.47 | 55.32 | 55.35 | 76.38 | 82.95 | 56.87 | 79.06 | 62.87 | 81.96 | 93.34 |
|  | U | 78.05 | 52.91 | 53.05 | 73.72 | 81.28 | 55.86 | 72.81 | 58.56 | 76.41 | 92.7 |
|  | V | 89.82 | 54.61 | 54.59 | 82.3 | 87.65 | 59.74 | 82.18 | 61.01 | 85.18 | 96.19 |
|  | W | 90.15 | 55.25 | 55.29 | 84.57 | 89.11 | 59.37 | 82.04 | 63.1 | 86.65 | 96.79 |
|  | X | 80.18 | 52.63 | 52.6 | 74.84 | 80.68 | 55.92 | 73.04 | 62.48 | 78.1 | 92.93 |
|  | Y | 81.73 | 54.65 | 54.7 | 72.3 | 79.67 | 57.49 | 76.38 | 53.48 | 79.33 | 92.46 |
|  | Z | 83.78 | 54.6 | 54.54 | 76.89 | 84.18 | 56.46 | 78.56 | 60.14 | 82.64 | 93.35 |

Table 3. Letters: Area under curve metric.

| Label | O | O-LS | A-LS | O-RS | A-RS | O-IF | A-IF | O-RC | A-RC | EC (Ours) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5-o-Clock-Shadow | 20.17 | 15.54 | 15.55 | 17.31 | 18.28 | 17.89 | 18.85 | 20.25 | 22.49 | 23.85 |
| Arched-Eyebrows | 19.29 | 15.52 | 15.51 | 17.54 | 17.69 | 18.34 | 18.65 | 19.87 | 19.76 | 20.93 |
| Attractive | 18.4 | 19.19 | 18.32 | 18.6 | 18.7 | 18.86 | 18.89 | 18.03 | 17.93 | 18.61 |
| Bags-Under-Eyes | 17.06 | 15.5 | 15.51 | 16.58 | 17.54 | 17.26 | 17.11 | 17.3 | 16.33 | 16.97 |
| Bangs | 20.88 | 15.52 | 15.52 | 20.6 | 19.97 | 20.49 | 21.43 | 22.12 | 19.9 | 22.08 |
| Bald | 18.52 | 15.48 | 15.49 | 16.6 | 16.98 | 15.48 | 17.92 | NA | NA | 28.34 |
| Big-Lips | 15.81 | 15.51 | 15.5 | 15.44 | 14.91 | 15.26 | 15.46 | 15.61 | 15.68 | 16.35 |
| Big-Nose | 16.58 | 15.53 | 15.53 | 15.92 | 16.58 | 15.95 | 16.23 | 16.88 | 16.12 | 17.23 |
| Black-Hair | 22.87 | 15.49 | 15.5 | 20.8 | 19.43 | 21.35 | 22.52 | 21.23 | 20.5 | 24.6 |
| Blond-Hair | 36.36 | 15.67 | 15.67 | 31.41 | 34.64 | 35.92 | 37.84 | 37.31 | 39.24 | 41.72 |
| Blurry | 16.75 | 15.44 | 15.48 | 16.38 | 16.06 | 16.48 | 16.66 | 16.66 | 16.81 | 17.79 |
| Brown-Hair | 20.88 | 15.48 | 15.48 | 18.82 | 20.45 | 20.64 | 21.34 | 20.94 | 20.77 | 21.57 |
| Bushy-Eyebrows | 19.11 | 15.49 | 15.49 | 17.87 | 18.12 | 18.18 | 18.7 | 18.88 | 19.2 | 21.3 |
| Chubby | 16.77 | 15.51 | 15.5 | 15.98 | 16.2 | 15.83 | 16.15 | 16.51 | 19.01 | 18.63 |
| Double-Chin | 18.05 | 15.53 | 15.52 | 16.73 | 17.31 | 16.17 | 17.28 | 16.66 | 22.57 | 22.3 |
| Eyeglasses | 16.48 | 15.51 | 15.5 | 16.14 | 17.12 | 16.48 | 16.78 | 17.74 | 15.88 | 15.44 |
| Goatee | 17.57 | 15.49 | 15.51 | 16.86 | 16.69 | 16.22 | 16.93 | 16.98 | 17.98 | 19.81 |
| Gray-Hair | 23.19 | 15.53 | 15.55 | 19.55 | 21.77 | 17.01 | 23.12 | 20.84 | 32.86 | 31.66 |
| Heavy-Makeup | 21.42 | 15.85 | 15.58 | 20.49 | 20.5 | 21.24 | 21.94 | 21.29 | 20.4 | 22.45 |
| High-Cheekbones | 19.17 | 15.38 | 15.17 | 18.59 | 18.79 | 18.91 | 18.89 | 19.78 | 18.77 | 20.07 |
| Male | 18.64 | 15.49 | 15.57 | 20.09 | 19.36 | 19.14 | 18.8 | 18.43 | 16.5 | 19.24 |
| Mouth-Slightly-Open | 17.37 | 15.66 | 15.56 | 17.43 | 17.34 | 17.42 | 17.16 | 17.79 | 17.2 | 17.77 |
| Mustache | 17.16 | 15.5 | 15.48 | 16.77 | 16.46 | 15.98 | 16.99 | 16.64 | 17.05 | 18.13 |
| Narrow-Eyes | 15.48 | 15.49 | 15.5 | 15.66 | 15.95 | 15.78 | 15.78 | 15.68 | 15.36 | 14.83 |
| No-Beard | 15.82 | 16.38 | 16.37 | 15.79 | 15.93 | 16.0 | 15.99 | 16.08 | 15.98 | 15.83 |
| Oval-Face | 18.02 | 15.5 | 15.5 | 16.76 | 17.11 | 17.18 | 17.37 | 18.56 | 18.37 | 19.22 |
| Pale-Skin | 18.15 | 15.45 | 15.48 | 19.89 | 19.15 | 16.95 | 20.95 | 18.75 | 17.94 | 19.55 |
| Pointy-Nose | 18.55 | 15.49 | 15.49 | 17.03 | 17.22 | 17.23 | 17.75 | 18.69 | 18.8 | 19.97 |
| Receding-Hairline | 18.5 | 15.5 | 15.5 | 17.09 | 17.47 | 16.67 | 17.55 | 18.61 | 22.2 | 22.38 |
| Rosy-Cheeks | 23.96 | 15.5 | 15.51 | 19.77 | 22.49 | 18.5 | 22.3 | 22.45 | 32.64 | 34.69 |
| Sideburns | 18.32 | 15.51 | 15.52 | 17.23 | 17.19 | 16.9 | 17.39 | 17.41 | 19.36 | 21.77 |
| Smiling | 19.46 | 16.01 | 15.52 | 18.97 | 19.05 | 19.31 | 19.15 | 19.89 | 18.81 | 20.15 |
| Straight-Hair | 16.1 | 15.5 | 15.5 | 15.97 | 16.25 | 15.91 | 16.02 | 16.24 | 16.13 | 16.29 |
| Wavy-Hair | 21.0 | 15.52 | 15.52 | 20.18 | 20.09 | 20.56 | 20.97 | 21.13 | 20.42 | 21.88 |
| Wearing-Earrings | 18.75 | 15.47 | 15.47 | 16.93 | 17.87 | 17.75 | 18.37 | 19.28 | 20.0 | 20.62 |
| Wearing-Hat | 18.32 | 15.48 | 15.49 | 18.33 | 18.27 | 17.44 | 19.95 | 20.12 | 15.16 | 16.82 |
| Wearing-Lipstick | 20.42 | 19.0 | 17.34 | 19.64 | 19.65 | 20.49 | 20.8 | 20.21 | 19.51 | 21.06 |
| Wearing-Necklace | 18.99 | 15.48 | 15.49 | 16.72 | 17.85 | 17.54 | 18.15 | 19.44 | 21.28 | 21.48 |
| Wearing-Necktie | 19.56 | 15.52 | 15.51 | 17.51 | 18.13 | 16.53 | 18.36 | 18.45 | 23.79 | 24.46 |
| Young | 15.51 | 16.22 | 16.21 | 15.54 | 15.63 | 15.55 | 15.55 | 15.56 | 15.51 | 15.57 |

Table 4. CelebA: Area under curve metric. We note that for Bald, there were no results for the Robust Covariance metrics. This is because due to the low rate of positive examples, it was not possible to tune Robust Covariance's hyper-parameters via cross-validation on the initial sample.

