

## Full Length Article



## Poidge-convexity in triangular lattices

Bo Wang<sup>a,c</sup>, Liping Yuan<sup>a,b,c,\*</sup>, Tudor Zamfirescu<sup>a,b,d,e</sup><sup>a</sup> School of Mathematical Sciences, Hebei Normal University, 050024 Shijiazhuang, PR China<sup>b</sup> Hebei International Joint Research Center for Mathematics and Interdisciplinary Science, 050024 Shijiazhuang, PR China<sup>c</sup> Hebei Key Laboratory of Computational Mathematics and Applications, 050024 Shijiazhuang, PR China<sup>d</sup> Fachbereich Mathematik, TU Dortmund, 44221 Dortmund, Germany<sup>e</sup> Roumanian Academy, Bucharest, Romania

## ARTICLE INFO

MSC:  
52A01  
52C20

## Keywords:

Poidge-convexity  
Triangular lattice  
Grid graph

## ABSTRACT

We adapt the notion of poidge-convexity to subgraphs of the triangular lattice, and investigate the poidge-convexity of graphs belonging to two families of such subgraphs: paths and grid graphs.

## 1. Introduction

At the 1974 meeting on convexity in Oberwolfach, the third author proposed the investigation of the following very general kind of convexity, also see [1]. Let  $\mathcal{F}$  be a family of sets in  $\mathbb{R}^d$ . A set  $M \subset \mathbb{R}^d$  is called  $\mathcal{F}$ -convex if for any pair of distinct points  $x, y \in M$  there is a set  $F \in \mathcal{F}$  such that  $x, y \in F$  and  $F \subset M$ . Usual convexity, affine linearity, arc-wise connectedness, polygonal connectedness, are just some examples of  $\mathcal{F}$ -convexity (for suitably chosen families  $\mathcal{F}$ ).

For  $M \subset \mathbb{R}^2$ ,  $\text{conv } M$  and  $\text{bd} M$  denote the convex hull and the boundary of  $M$ , respectively.

In 1980, Blind, Valette and the third author [1] first investigated rectangular convexity, the case when  $\mathcal{F}$  is the family of all (2-dimensional) rectangles, which was also studied by Böröczky, Jr. [2]. In 2014, the third author [10] studied the right convexity, the case with  $\mathcal{F}$  consisting of all (2-dimensional) right triangles. Later, the last two authors [8,9] investigated the  $rt$ -convexity, which is a discrete generalization of the right convexity. Recently, the authors [7] studied poidge-convexity, another generalization of the right convexity.

But  $\mathcal{F}$  can be taken to be any family of sets, not just subsets of  $\mathbb{R}^d$ . For a subset of the vertex set of a graph, the  $g$ -convexity investigated by Farber and Jamison [6], the  $T$ -convexity studied by Changat and Mathew [4], Duchet's  $M$ -convexity [5], and the  $P_3$ -convexity considered by Centeno et al. [3] can also be regarded as examples of  $\mathcal{F}$ -convexity for suitable families  $\mathcal{F}$ .

Planar lattices, square, triangular or hexagonal, play a crucial role in many real life problems. There is no need of enumerating them here. These problems are converted into mathematical ones. In this mathematical paper, we propose a mathematical problem, and solve it, hoping that it will be useful one day.

A *lattice graph* is a finite subgraph of the infinite planar triangular lattice  $\mathcal{T}$  in  $\mathbb{R}^2$ . The triangles in  $\mathcal{T}$  are equilateral, of side 1.

\* Corresponding author at: School of Mathematical Sciences, Hebei Normal University, 050024 Shijiazhuang, PR China.  
E-mail address: [lp yuan@hebtu.edu.cn](mailto:lp yuan@hebtu.edu.cn) (L. Yuan).

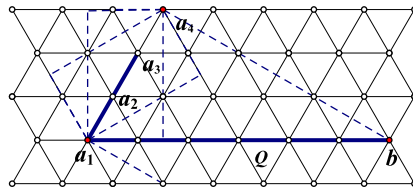


Fig. 1.  $a_1 \in V(P)$  and  $a_3a_4 \notin E(P)$ .

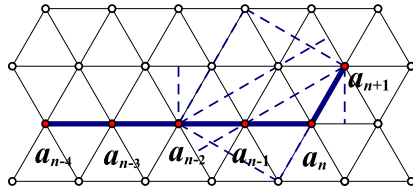


Fig. 2.  $\angle a_{n-1}a_n a_{n+1} = \frac{2\pi}{3}$ .

A path in a lattice graph is called *linear* if all its vertices are on the same lattice line.

Now, we adapt to lattice graphs the notion of *poidge-convexity*, introduced for planar sets in [7].

A graph in  $\mathcal{T}$  consisting of an isolated vertex  $v$  of  $\mathcal{T}$  plus a linear path  $P$  in  $\mathcal{T}$ , such that  $v$  and the endvertices of  $P$  are the vertices of a non-degenerate right triangle, is called a *poidge*. A lattice graph  $G$  is *poidge-convex* if, for any pair  $x, y \in V(G)$ , there is a poidge with  $x, y$  as vertices, subgraph of  $G$ .

For any vertices  $a, b, c$  of  $\mathcal{T}$ , we denote by  $\widehat{abc}$  the angle of  $abc$  at  $b$ , and by  $\angle abc$  its measure. Also, we denote by  $\delta(a, b)$  the length of the shortest path in  $\mathcal{T}$  between  $a$  and  $b$ .

For any two vertices  $a, b$ , we denote by  $\overline{ab}$  the lattice line through  $a, b$ , and by  $ab$  the line-segment from  $a$  to  $b$ . If  $a, b$  are neighbours in  $\mathcal{T}$ , then  $ab \in E(\mathcal{T})$ . Put  $a_1a_2 \dots a_n = a_1a_2 \cup a_2a_3 \cup \dots \cup a_{n-1}a_n$  and  $[a_1a_2 \dots a_n] = a_1a_2 \dots a_n \cup a_na_1$ .

For any two parallel lattice lines  $L_1$  and  $L_2$ , we define the intrinsic distance  $\rho(L_1, L_2)$  to be the Euclidean distance between them multiplied by  $\frac{2}{\sqrt{3}}$ . Let  $\pi(a, L)$  be the orthogonal projection of  $a$  onto the line  $L$ .

Let  $C$  be a (finite) cycle in  $\mathcal{T}$ . We denote by  $\text{set}(C)$  the set of all points lying on the edges and at the vertices of  $C$ . All vertices and edges lying on  $C$  or in the bounded component of the complement of  $\text{set}(C)$  in  $\mathbb{R}^2$  form a graph called *grid graph* (of *boundary*  $C$ ).

In this paper, our goal is to investigate the poidge-convexity of paths and grid graphs in  $\mathcal{T}$ . We achieve full characterizations.

## 2. Poidge-convex paths

First, we discover that linear paths of a poidge-convex path cannot be too long.

**Lemma 1.** *Suppose that  $a_1, a_2, a_3$  and  $a_4$  are four consecutive vertices on a lattice line, and  $P$  is a poidge-convex path, such that the length of any linear subpath of  $P$  is at most 3. If  $a_1 \in V(P)$  and at least one of the three edges  $a_1a_2, a_2a_3$  and  $a_3a_4$  is not in  $E(P)$ , then  $a_4 \notin V(P)$  (see Fig. 1).*

**Proof.** By the poidge-convexity of  $P$ , if  $a_4 \in V(P)$ , then there exists a poidge in  $P$  containing  $a_1$  and  $a_4$ . Then there exists a linear subpath  $Q$  of  $P$ , such that  $Q$  only contains one of the two vertices  $a_1$  and  $a_4$ , say  $a_1$ , and  $Q \cup \{a_4\}$  is a poidge. However,  $Q$  is a subpath of length 6, which is not possible, see Fig. 1.  $\square$

**Lemma 2.** *If  $P$  is a poidge-convex path, then the length of any linear subpath of  $P$  is at most two.*

**Proof.** Let  $P$  be a poidge-convex path. Suppose without loss of generality that  $a_0a_1 \dots a_{n-1}a_n$  ( $n \geq 3$ ) is a maximal linear subpath of  $P$ .

If  $\angle a_{n-1}a_n a_{n+1} = \frac{2\pi}{3}$ , as shown in Fig. 2, then there is no poidge in  $P$  containing  $a_{n-2}$  and  $a_{n+1}$ . Hence,  $\angle a_{n-1}a_n a_{n+1} = \frac{\pi}{3}$ .

Case 1.  $n = 3$ , as shown in Fig. 3. By the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $a_0$  and  $a_4$ . This implies that  $a_0b_1b_2 \subset P$ , see Fig. 3. By the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $a_3$  and  $b_1$ . Therefore, we have  $a_4b_3 \in E(P)$ . Then  $b_2b_3 \notin E(P)$ . The poidge in  $P$  containing  $b_2$  and  $b_3$  must be  $a_4b_3b_4 \cup \{b_2\}$  or  $b_1b_2b_4 \cup \{b_3\}$ . So,  $b_4 \in V(P)$ , but only one of the two edges  $b_2b_4$  and  $b_3b_4$  belongs to  $P$ . This contradicts Lemma 1.

Case 2.  $n > 3$ , as shown in Fig. 4. Then no poidge in  $P$  contains  $a_{n-3}$  and  $a_{n+1}$ .

Hence, the length of any linear subpath of  $P$  is at most two.  $\square$

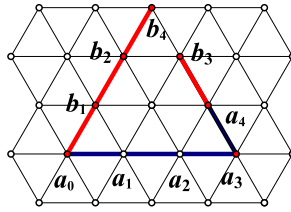


Fig. 3.  $n = 3$ .

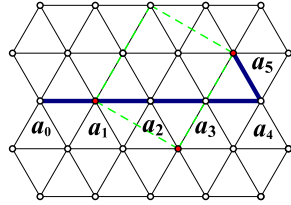


Fig. 4.  $n > 3$ .

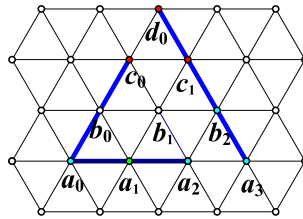


Fig. 5.  $b_2 \in V(P)$ .

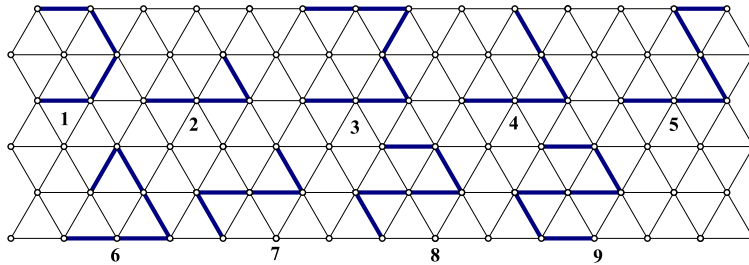


Fig. 6. The 9 kinds of poidge-convex paths in  $\mathcal{T}$ .

**Lemma 3.** Let  $P \subset \mathcal{T}$  be a poidge-convex path such that the length of a maximal linear subpath of  $P$  is 2. Suppose that  $a_0a_1a_2$  is a linear subpath of  $P$ . If  $u \in V(P)$  is a neighbour of  $a_2$  in  $\mathcal{T}$  distinct from  $a_1$ , then  $\angle ua_2a_1 = \frac{\pi}{3}$ .

**Proof.** Suppose, on the contrary, that  $u = b_2 \in V(P)$ , see Fig. 5. Due to the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $a_0$  and  $b_2$ , which means that  $a_0b_0c_0 \subset P$ . Since there exists a poidge in  $P$  containing  $a_1$  and  $b_2$ , and according to Lemma 1, we have  $b_2c_1 \in E(P)$ . Then  $a_2b_2 \notin E(P)$  or  $c_0c_1 \notin E(P)$ . Assume without loss of generality that  $c_0c_1 \notin E(P)$ . Since there exists a poidge in  $P$  containing  $c_0$  and  $c_1$ ,  $c_1d_0 \in E(P)$ . This contradicts Lemma 1; Lemma 3 is proven.  $\square$

Now, we present all poidge-convex paths of  $\mathcal{T}$ .

**Theorem 1.** There are precisely 9 kinds of poidge-convex paths in  $\mathcal{T}$  (see Fig. 6).

**Proof.** Let  $P$  be a poidge-convex path and label some vertices of  $\mathcal{T}$  as shown in Fig. 7. Due to Lemma 2, the length of a maximal linear subpath of  $P$  is at most 2.

First, suppose the length of a maximal linear subpath of  $P$  is 1.

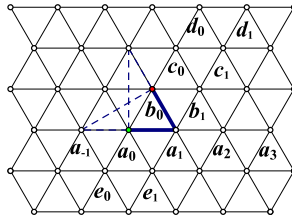


Fig. 7.  $a_1 b_0 \in E(P)$ .

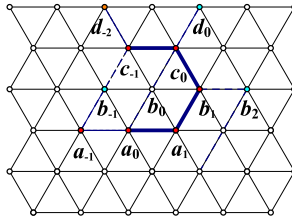


Fig. 8.  $a_1 b_1 \in E(P)$ .

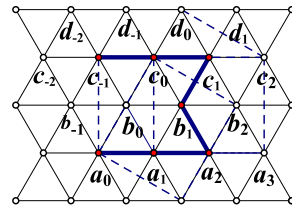


Fig. 9.  $b_1 c_1 \in E(P)$ .

Assume without loss of generality that  $a_0 a_1 \in E(P)$  and  $a_0$  is an endvertex of  $P$ . The angle between any two incident edges is  $\frac{2\pi}{3}$ . Otherwise, suppose  $a_1 b_0 \in E(P)$ , see Fig. 7. Since there exists a poidge in  $P$  containing  $a_0$  and  $b_0$ , there must exist a linear subpath of length 2 containing  $a_0$  or  $b_0$ , contradicting our assumption.

Suppose without loss of generality that  $a_1 b_1 \in E(P)$ , see Fig. 8. If  $b_1 b_2 \in E(P)$ , then, by the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $a_0$  and  $b_2$ , which implies that there exists a linear subpath of length 2 containing  $a_0$  or  $b_2$ , a contradiction. So  $b_1 c_0 \in E(P)$ . Since  $a_0 b_0 \notin E(P)$ , by Lemma 1,  $d_0 \notin V(P)$ . Then  $c_0 c_{-1} \in E(P)$ . Since  $a_1 b_0 \notin E(P)$ , by Lemma 1,  $d_{-2} \notin V(P)$ . Suppose that  $b_{-1} \in V(P)$ . Due to the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $a_0$  and  $b_{-1}$ . Because  $a_0 b_{-1} \notin E(P)$ , this poidge must include a linear subpath of length 2 containing  $a_0$  or  $b_{-1}$ , a contradiction. Thus,  $c_{-1}$  is an endvertex of  $P$ , and it is clear that  $P = a_0 a_1 b_1 c_0 c_{-1}$  is poidge-convex (Type 1).

Second, suppose the length of a maximal linear subpath of  $P$  is 2.

Assume without loss of generality that  $a_0 a_1 a_2$  is a linear subpath of length 2, and  $b_1 a_2$  belongs to  $E(P)$ , by Lemma 3. According to Lemma 3, we have  $b_2 \notin V(P)$ .

Case 1.  $a_0$  is an endvertex of  $P$ . If  $b_1$  is also an endvertex, it is clear that  $P = a_0 a_1 a_2 b_1$  is a poidge-convex path, see Fig. 6 (Type 2). Now, suppose that  $b_1$  is not an endvertex of  $P$ . Assume  $b_2 \in V(P)$ . Since  $a_0 b_0 \notin E(P)$  and  $a_2 b_2 \notin E(P)$ , we have no poidge in  $P$  containing  $a_0$  and  $b_2$ . Hence,  $b_2 \notin V(P)$ . Then  $b_1 c_1 \in E(P)$ , or  $b_1 c_0 \in E(P)$ , or  $b_1 b_0 \in E(P)$ .

Subcase 1.1.  $b_1 c_1 \in E(P)$ , see Fig. 9. Then at most one of the two edges  $c_1 c_2$  and  $c_1 c_0$  belongs to  $E(P)$ . Assuming  $c_1 c_2 \in E(P)$ , there is no poidge in  $P$  containing  $a_1$  and  $c_2$ , because  $a_2 a_3 \notin E(P)$ . Hence,  $c_1 c_2 \notin E(P)$ . Since  $a_0 b_0, a_1 b_0, a_1 b_1, b_1 c_0$  are not in  $E(P)$ , according to Lemma 1 we conclude that  $d_0, d_{-2}, d_1, d_{-1}$  do not belong to  $V(P)$ . If  $P$  has an endvertex at  $c_1$ , then it includes no poidge containing  $b_1$  and  $c_1$ . Hence,  $c_1 c_0 \in E(P)$ . By the poidge-convexity of  $P$ , there exists a poidge in  $P$  containing  $b_1$  and  $c_0$ . This yields  $c_{-1} c_0 \in E(P)$ . Since  $c_0 b_0, a_1 b_0, a_1 b_1$  and  $b_1 c_0$  are not in  $E(P)$ , we have no poidge in  $P$  containing  $b_1$  and  $b_0$ . This implies that  $b_0 \notin V(P)$ . Thus,  $c_{-1}$  is an endvertex of  $P$ . Obviously,  $P = a_0 a_1 a_2 b_1 c_1 c_0 c_{-1}$  is a poidge-convex path (Type 3).

Subcase 1.2.  $b_1 c_0 \in E(P)$ , see Fig. 10. If  $c_0$  is an endvertex of  $P$ , then it is easily seen that  $P = a_0 a_1 a_2 b_1 c_0$  is a poidge-convex path (Type 4). Otherwise, since  $a_0 b_0 \notin E(P)$  and  $a_2 b_1 c_0$  is a linear subpath of length 2, we have  $d_0 \notin V(P)$ ,  $c_0 d_{-1} \notin E(P)$  and  $c_0 c_{-1} \notin E(P)$ . Thus  $c_0 c_1 \in E(P)$  or  $c_0 b_0 \in E(P)$ .

Subcase 1.2.1.  $c_0 c_1 \in E(P)$ , see Fig. 10. Since  $a_0 b_0, a_1 b_1$  are not in  $E(P)$ , by Lemma 1,  $d_0$  and  $d_1$  do not belong to  $V(P)$ . Assume  $c_1 c_2 \in E(P)$ . From  $a_0 b_0, a_2 b_2, c_{-1} c_0, c_2 b_3, a_2 a_3 \notin E(P)$ , it follows that no poidge in  $P$  contains  $a_0$  and  $c_2$ . Hence,  $c_1 c_2 \notin E(P)$ . Since  $b_2 \notin V(P)$  and  $c_1 b_1 \notin E(P)$ ,  $c_1$  is an endvertex of  $P$ . It is clear that  $P = a_0 a_1 a_2 b_1 c_0 c_1$  is a poidge-convex path (Type 5).

Subcase 1.2.2.  $c_0 b_0 \in E(P)$ , see Fig. 11. Next, we claim that  $b_0$  is an endvertex of  $P$ . Suppose this is not the case. The existence of  $a_2 b_1 c_0 \subset P$  prohibits  $c_{-1} \notin V(P)$ , by Lemma 3. Then  $b_0 b_{-1} \in E(P)$ . Since there is a poidge in  $P$  containing  $a_0$  and  $b_{-1}$ ,  $b_{-1} b_{-2} \in E(P)$ .

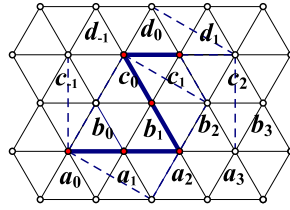


Fig. 10.  $c_0c_1 \in E(P)$ .

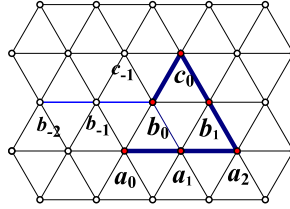


Fig. 11.  $c_0b_0 \in E(P)$ .

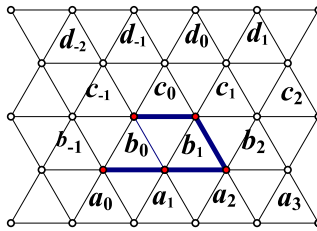


Fig. 12.  $b_1b_0 \in E(P)$ .

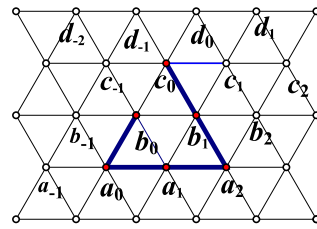


Fig. 13.  $a_0b_0 \in E(P)$ .

This contradicts Lemma 1. Hence,  $b_0$  is an endvertex of  $P$ . It is easy to check that  $P = a_0a_1a_2b_1c_0b_0$  is a poidge-convex path (Type 6).

Subcase 1.3.  $b_1b_0 \in E(P)$ , see Fig. 12. Since  $a_0b_0$  and  $b_1c_0$  are not in  $E(P)$ , according to Lemma 1 we conclude that  $d_0$  and  $d_{-1}$  do not belong to  $V(P)$ . By Lemma 3,  $a_0a_1a_2 \subset P$  yields  $b_{-1} \notin V(P)$ . Hence,  $P$  includes no poidge containing  $b_1$  and  $b_0$ , which is not possible.

Case 2.  $a_0$  is not an endvertex of  $P$ .

Subcase 2.1.  $a_0b_0 \in E(P)$ , see Fig. 13. Then  $b_0b_1 \notin E(P)$ . Since there exists a poidge in  $P$  containing  $b_0$  and  $b_1$ , that poidge must be  $a_0b_0c_0 \cup b_1$  or  $a_2b_1c_0 \cup b_0$ . Suppose without loss of generality that  $b_1c_0 \in E(P)$  and  $b_0c_0 \notin E(P)$ . Then  $d_0 \notin V(P)$ . Since  $a_2b_1c_0$  is a linear subpath of length 2, we have  $c_0d_{-1} \notin E(P)$  and  $d_{-1} \notin V(P)$ . Next, we claim that  $c_0$  is an endvertex of  $P$ . Suppose the claim is false. Since  $a_2b_1c_0 \subset P$ , by Lemma 3,  $c_{-1} \notin V(P)$ . Then  $c_0c_1 \in E(P)$ . Recall that  $b_2 \notin V(P)$ . Since  $d_0, b_2$  and  $c_{-1}$  are not in  $V(P)$  and  $b_0a_1 \notin E(P)$ , no poidge in  $P$  contains  $c_1$  and  $b_0$ , a contradiction.

Thus,  $c_0$  is an endvertex of  $P$ . From the analysis of the symmetric Subcase 1.2.2, it follows that  $b_0$  is an endvertex of  $P$ , and  $P = b_0a_0a_1a_2b_1c_0$  is poidge-convex (Type 6).

Subcase 2.2.  $a_0e_1 \in E(P)$ , see Fig. 14. If  $e_1$  and  $b_1$  are endvertices of  $P$ , then  $P = e_1a_0a_1a_2b_1$ , and  $P$  is poidge-convex (Type 7). Now we suppose that at least one of the vertices  $e_1, b_1$  is not an endvertex of  $P$ . Assume without loss of generality that  $b_1$  is not an endvertex of  $P$ .

Recall that  $b_2 \notin V(P)$ . Since  $e_1 \in V(P)$  and  $e_1a_1 \notin E(P)$ , according to Lemma 1,  $c_1 \notin V(P)$ . Then  $b_1c_0 \in E(P)$  or  $b_1b_0 \in E(P)$ .

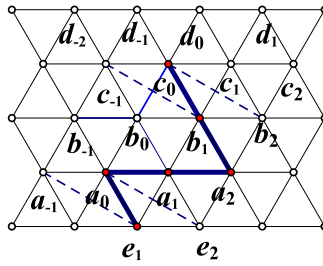


Fig. 14.  $a_0e_1 \in E(P)$  and  $b_1c_0 \in E(P)$ .

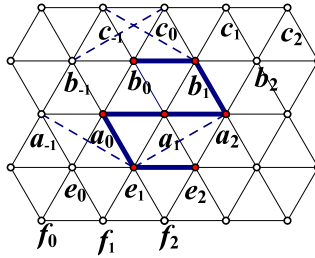


Fig. 15.  $b_1b_0 \in E(P)$ .

Subcase 2.2.1. If  $b_1c_0 \in E(P)$ , then  $c_0d_{-1} \notin E(P)$ , see Fig. 14. Due to  $a_2b_1c_0 \subset P$ , by Lemma 3, we have  $c_{-1} \notin V(P)$ . Since  $a_0a_{-1} \notin E(P)$  and  $a_2 \in V(P)$ , from Lemma 1 we get  $a_{-1} \notin V(P)$ . Next, we claim that  $c_0$  is an endvertex of  $P$ . Suppose the claim were false. Then  $c_0b_0 \in E(P)$  or  $c_0c_1 \in E(P)$ . The latter case is excluded by Lemma 1. If  $c_0b_0 \in E(P)$ , there is no poidge in  $P$  containing  $e_1$  and  $b_0$ , like in the symmetric subcase 2.1, a contradiction. Hence,  $c_0$  is an endvertex of  $P$ . From the analysis of the symmetric Subcase 1.2.1, it follows that  $e_1$  is an endvertex of  $P$ . It is easily seen that  $P = e_1a_0a_1a_2b_1c_0$  is a poidge-convex path (Type 5).

Subcase 2.2.2.  $b_1b_0 \in E(P)$ , see Fig. 15. Since  $a_0a_1a_2 \subset P$ , by Lemma 3,  $b_{-1} \notin V(P)$ . Assume  $c_{-1} \in V(P)$  or  $c_0 \in V(P)$ . Due to  $b_1c_0, a_1e_1 \notin E(P)$  and  $b_{-1} \notin V(P)$ , no poidge in  $P$  contains  $e_1, c_{-1}$  or  $e_1, c_0$ . Hence,  $c_0$  and  $c_{-1}$  do not belong to  $V(P)$ . Thus,  $b_0$  is an endvertex of  $P$ .

If  $e_1$  is also an endvertex of  $P$ , then  $P = e_1a_0a_1a_2b_1b_0$ , and it is easily checked that  $P$  is poidge-convex (Type 8), see Fig. 1. Suppose now that  $e_1$  is not an endvertex of  $P$ . Since  $b_1a_1 \notin E(P)$  and  $b_1 \in V(P)$ , by Lemma 1, we have  $f_1 \notin V(P)$ . On the other hand, due to  $a_0a_1a_2 \subset P$  and  $b_0 \in V(P)$ , according to Lemma 3, we have  $e_0 \notin V(P)$  and  $e_1f_2 \notin E(P)$ . Thus,  $e_1e_2 \in E(P)$ . Like in the case of  $b_0$ , we conclude that  $e_2$  is an endvertex of  $P$ . Then  $P = e_2e_1a_0a_1a_2b_1b_0$ , and this path is poidge-convex (Type 9).  $\square$

### 3. Poidge-convex grid graphs

We investigate here the grid graphs in  $\mathcal{T}$  with convex boundary, and characterize those which are poidge-convex.

**Proposition 1.** Let  $L_1$  and  $L_2$  be two parallel lattice lines, and  $a \in L_1$  be a lattice point. Then  $\pi(a, L_2)$  is a lattice point if and only if  $\rho(L_1, L_2)$  is even.

**Proposition 2.** Suppose that  $a, b$  are two vertices on the same lattice line. If  $L_a$  and  $L_b$  are two parallel distinct lattice lines containing  $a$  and  $b$ , respectively, then  $\rho(L_a, L_b) = \delta(a, b)$ .

In the following results,  $G$  is a grid graph and  $C$  its boundary cycle.

**Lemma 4.** Let  $G$  be a grid graph with  $\text{set}(C)$  convex. If  $G$  is poidge-convex and the angles at two vertices of  $\text{set}(C)$  are acute, then the vertices are neighbours (in  $\text{set}(C)$ ).

**Proof.** Suppose the lemma were false. Let  $a$  and  $b$  be two vertices of  $\text{set}(C)$  such that the angles at  $a$  and  $b$  are acute, and  $a$  and  $b$  are not neighbours. According to the convexity of  $\text{set}(C)$ , there exist two parallel lattice lines  $\overline{aa_1}$  and  $\overline{bb_1}$ , such that  $\text{set}(C) \subseteq \text{conv}\{a, a_1, b, b_1\}$  and  $a_1\overline{b}, \overline{ab_1}$  are also two parallel lattice lines, as shown in Fig. 16. There exist two parallel lattice lines  $L_a$  and  $L_b$ , such that  $L_a \cap \text{conv}\{a, a_1, b, b_1\} = \{a\}$ ,  $L_b \cap \text{conv}\{a, a_1, b, b_1\} = \{b\}$ ; then  $\text{conv}\{a, a_1, b, b_1\}$  is contained in the strip bounded by  $L_a$  and  $L_b$ . Thus,  $\pi(a, L_b)$  and  $\pi(b, L_a)$  are not vertices of  $G$ , or they equal  $b$  and  $a$ , respectively. On the other hand, since  $\widehat{aa_1b}$  and  $\widehat{ab_1b}$  are obtuse,  $\pi(a, a_1b), \pi(b, \overline{aa_1}), \pi(a, \overline{bb_1})$  and  $\pi(b, \overline{ab_1})$  do not belong to  $V(G)$ . Thus, there is no poidge in  $G$  containing  $a$  and  $b$ , contrary to the poidge-convexity of  $G$ .  $\square$

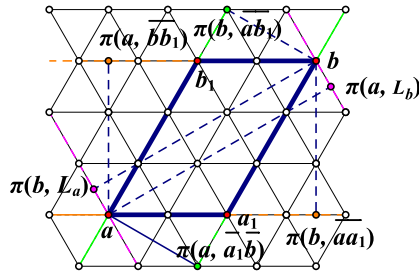


Fig. 16.  $\text{conv}\{a, a_1, b, b_1\}$ .

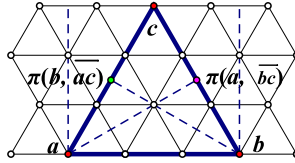


Fig. 17.  $\text{set}(C) = ab \cup bc \cup ca$ .

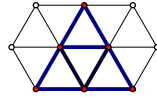


Fig. 18. The grid graph  $G$ .

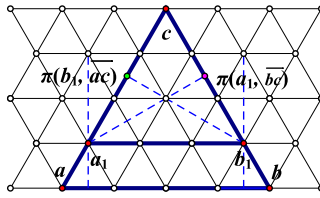


Fig. 19.  $\text{set}(C) \subset \text{conv}\{a, b, c\}$ .

**Lemma 5.** Let  $G$  be a grid graph, such that  $\text{set}(C)$  is a convex polygon having two neighbouring vertices  $a, b$  with acute angles. If  $G$  is poidge-convex, then  $\delta(a, b)$  is even.

**Proof.** On the contrary, suppose that  $\delta(a, b)$  is odd. By Propositions 1 and 2,  $\pi(a, \overline{bc})$ ,  $\pi(b, \overline{ac})$  are not in  $V(G)$ , see Fig. 17. On the other hand, since the angles at  $a, b$  are acute, there exists no vertex  $u \in V(G)$  such that  $\angle uab = \frac{\pi}{2}$  or  $\angle uba = \frac{\pi}{2}$ . Therefore, no poidge in  $G$  contains  $a$  and  $b$ , and this is not possible.  $\square$

**Theorem 2.** Suppose that  $G$  is a grid graph, such that  $\text{set}(C)$  is a convex polygon having at least two vertices  $a, b$  with acute angles. Then  $G$  is poidge-convex if and only if  $\text{set}(C)$  is an equilateral triangle with side length 2 (see Fig. 18).

**Proof.** Let  $a, b$  be two vertices of  $\text{set}(C)$ , such that the angles at  $a, b$  are acute. We now claim that  $\delta(a, b) = 2$ . Suppose, contrary to our claim, that  $\delta(a, b) \geq 3$ .

According to Lemma 4, the vertices  $a, b$  are neighbours in  $\text{set}(C)$ ; thus, there exists a lattice point  $c$  such that  $\text{set}(C) \subset \text{conv}\{a, b, c\}$ , see Fig. 19. Let  $a_1 \in ac \cap V(C)$  and  $b_1 \in bc \cap V(C)$  such that  $\delta(a, a_1) = \delta(b, b_1) = 1$ . Then  $\angle ca_1b_1 < \frac{\pi}{2}$  and  $\angle cb_1a_1 < \frac{\pi}{2}$ . Thus, there exists no vertex  $u \in V(G)$  such that  $\angle ua_1b_1 = \frac{\pi}{2}$  or  $\angle ub_1a_1 = \frac{\pi}{2}$ . Since, by Lemma 5,  $\delta(a, b)$  is even,  $\delta(a_1, b_1)$  is odd. Then, by Propositions 1 and 2,  $\pi(a_1, \overline{bc})$  and  $\pi(b_1, \overline{ac})$  are not in  $V(G)$ . Therefore, no poidge in  $G$  contains  $a_1$  and  $b_1$ , a contradiction.  $\square$

**Lemma 6.** Let  $G$  be a grid graph such that  $\text{set}(C)$  is a convex pentagon  $[abcde]$ ,  $\delta(a, e) \geq \delta(a, b)$  and the angle at  $a$  is acute. If  $\delta(a, e)$  is odd, then  $G$  is not poidge-convex.

**Proof.** We first suppose that  $\delta(a, e) > 1$ . Let  $f \in ed \cap V(G)$  such that  $\delta(e, f) = 2$ . Assume that  $\overline{aa'}$ ,  $\overline{ee'}$  and  $\overline{ff'}$  are three lattice lines parallel to  $cd$ , and  $\overline{ff''}$  is a lattice line parallel to  $bc$ , where  $e', f'$  and  $f''$  belong to  $V(C)$ , see Fig. 20.

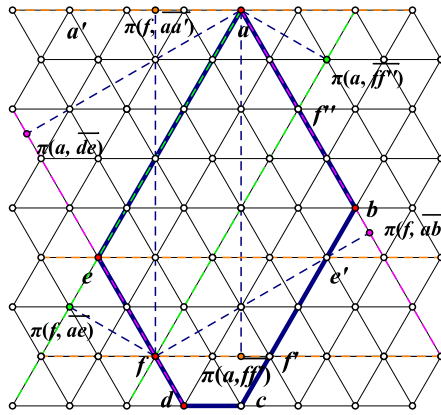


Fig. 20.  $\delta(d, e) > 1$ .

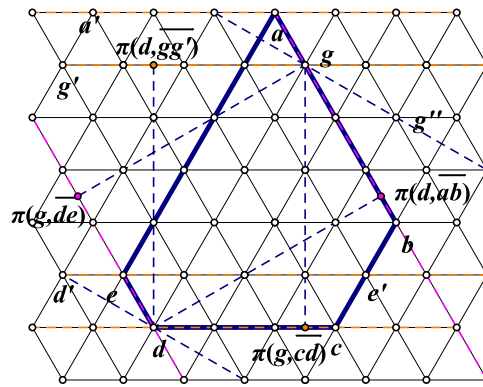


Fig. 21.  $g \in ab \cap V(G)$  and  $\delta(a, g) = 1$ .

Since  $\delta(a, e)$  is odd,  $\rho(\overline{ab}, \overline{de})$  is odd, according to Proposition 2. By Proposition 1,  $\pi(f, \overline{ab})$  and  $\pi(a, \overline{de})$  do not belong to  $V(G)$ . Because  $\rho(aa', \overline{ff'}) = \rho(aa', ee') + \rho(ee', \overline{ff'})$ ,  $\rho(aa', \overline{ff'})$  is odd. According to Proposition 1,  $\pi(f, \overline{aa'})$  and  $\pi(a, \overline{ff'})$  do not belong to  $V(G)$ . Since  $\widehat{ae'f}$  and  $\widehat{af''f}$  are obtuse,  $\pi(f, \overline{ae'})$  and  $\pi(a, \overline{ff''})$  do not belong to  $V(G)$ . Thus, no poidge in  $G$  contains  $f$  and  $a$ .

Now, assume  $\delta(d, e) = 1$ .

Let  $g \in ab \cap V(G)$  such that  $\delta(a, g) = 1$ . Assume that  $\overline{gg'}$  is a lattice line parallel to  $\overline{cd}$ , while  $\overline{dd'}$  and  $\overline{gg''}$  are orthogonal to  $\overline{dg}$ , see Fig. 21.

As before,  $\rho(\overline{ab}, \overline{de})$  is odd. By Proposition 1,  $\pi(g, \overline{de})$  and  $\pi(d, \overline{ab})$  do not belong to  $V(G)$ . Since  $\delta(g, d) = \delta(a, e)$  is odd,  $\rho(\overline{gg'}, \overline{cd})$  is odd. According to Proposition 1,  $\pi(d, \overline{gg'})$  and  $\pi(g, \overline{cd})$  do not belong to  $V(G)$ . Since  $\widehat{agd}$  is obtuse and  $d(a, g) = 1$ ,  $\overline{gg''} \cap (V(G) \setminus \{g\}) = \emptyset$ . Since  $\widehat{cde}$  is obtuse,  $\overline{dd'} \cap (V(G) \setminus \{d\}) = \emptyset$ . Thus, no poidge in  $G$  contains  $d$  and  $g$ .  $\square$

**Lemma 7.** Let  $G$  be a grid graph and  $\text{set}(C)$  a convex pentagon  $[abcde]$ , such that  $\delta(a, e) \geq \delta(a, b)$  and the angle at  $a$  is acute. If  $G$  is poidge-convex, then  $1 \leq \delta(d, e) \leq 2$ .

**Proof.** Set  $\delta(a, e) = n$ . Since  $\text{set}(C)$  is a convex pentagon  $abcde$ ,  $1 \leq \delta(d, e)$ . Now we claim that  $\delta(d, e) \leq 2$ . Suppose, contrary to our claim, that  $\delta(d, e) \geq 3$ . Let  $p \in ed \cap V(G)$  such that  $\delta(e, p) = 3$ . Suppose  $\overline{pq}$  is a lattice line parallel to  $\overline{cd}$  such that  $q \in V(G)$  and  $\delta(p, q) = 1$ , see Fig. 22. Since  $\delta(a, e) = n$  is even by Lemma 6, and  $\delta(e, p) = 3$  by Proposition 2,  $\rho(\overline{aa'}, \overline{pq})$  is odd. According to Proposition 1,  $\pi(q, \overline{aa'})$  and  $\pi(a, \overline{pq})$  are not vertices of  $(G)$ . Because  $\delta(p, q) = 1$ ,  $\rho(\overline{ab}, \overline{qq'}) = \rho(\overline{ab}, \overline{ed}) - \rho(\overline{ed}, \overline{qq'}) = n - 1$  is odd. According to Proposition 1,  $\pi(a, \overline{qq'})$  and  $\pi(q, \overline{ab})$  are not in  $V(G)$ . Since  $\widehat{eab}$  is acute,  $\pi(a, \overline{qq''}) \notin V(G)$ . Since  $\delta(e, q) = 3$  and  $\delta(p, q) = 1$ ,  $\rho(\overline{ae}, \overline{qq''}) = 4$ . Thus,  $\pi(q, \overline{ae}) \notin V(G)$ . Therefore, there is no poidge in  $G$  containing  $a$  and  $q$ , a contradiction. Hence,  $\delta(d, e) \leq 2$ .  $\square$

**Lemma 8.** Let  $G$  be a grid graph such that  $\text{set}(C)$  is a convex pentagon  $[abcde]$ ,  $\delta(a, e) \geq \delta(a, b)$ , and the angle at  $a$  is acute. If  $G$  is poidge-convex, then  $\delta(a, b) = \delta(a, e)$ .

**Proof.** By Lemma 6,  $\delta(a, e)$  is even. We first suppose that  $\delta(a, b)$  is odd. Then  $\delta(b, c) \geq 2$ . Let  $f \in bc \cap V(G)$  such that  $\delta(b, f) = 2$ . Assume that  $\overline{aa'}, \overline{ff'}$  and  $\overline{bb'}$  are three lattice lines parallel to  $\overline{cd}$ , and  $\overline{ff''}$  is a lattice line parallel to  $\overline{ab}$ , where  $b', f'$  and  $f''$  belong to  $V(C)$ , see Fig. 23.



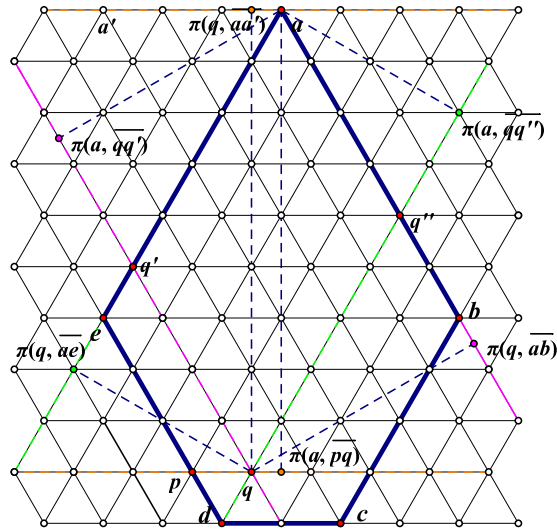


Fig. 22.  $\delta(d, e) \geq 3$ .

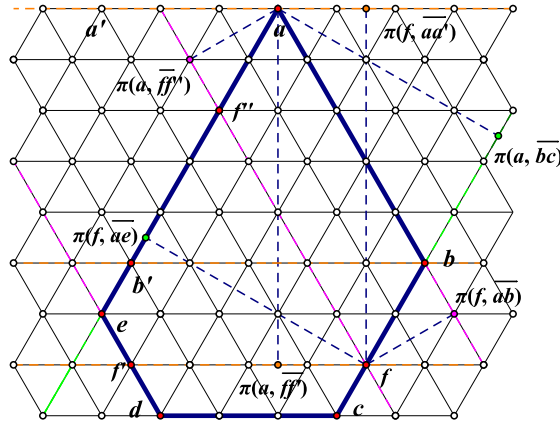


Fig. 23.  $f \in bc \cap V(G)$  and  $\delta(b, f) = 2$ .

By Proposition 1,  $\pi(a, \overline{bc})$  and  $\pi(f, \overline{ae})$  are not vertices of  $G$ . Moreover,  $\rho(\overline{aa'}, \overline{ff'}) = \rho(\overline{aa'}, \overline{bb'}) + \rho(\overline{bb'}, \overline{ff'}) = \delta(a, b) = 2$  is odd. Again by Proposition 1,  $\pi(a, \overline{ff'})$  and  $\pi(f, \overline{aa'})$  do not belong to  $V(G)$ . Since  $\overline{eab}$  is acute and  $\overline{abc}$  obtuse,  $\pi(f, \overline{ab})$ ,  $\pi(a, \overline{ff'}) \notin V(G)$ . Thus, there is no poidge in  $G$  containing  $f$  and  $a$ . Therefore,  $\delta(a, b)$  is even.

Next we suppose that  $\delta(a, b) < \delta(a, e) - 2$ . Since  $\delta(a, b)$  and  $\delta(a, e)$  are even,  $\delta(a, b) \leq \delta(a, e) - 4$ . Let  $g \in ed \cap V(G)$  and  $h \in ab \cap V(G)$  such that  $\delta(e, g) = 1$  and  $\delta(a, h) = 2$ . Assume that  $\overline{gm}$ ,  $\overline{ee'}$  and  $\overline{hk}$  are three lattice lines parallel to  $cd$ , where  $m \in V(C)$  and  $k \in gk \cap V(G)$ , see Fig. 24. Suppose that  $\overline{kk'}$  and  $\overline{mm'}$  are two lattice lines parallel to  $ab$ , where  $m'$  and  $k'$  belong to  $V(C)$ .

Since  $\delta(a, b)$  is even and  $\delta(e, g) = 1$ ,  $\rho(\overline{bc}, \overline{gk}) = \rho(\overline{bc}, \overline{ae}) - \rho(\overline{ae}, \overline{gk}) = \rho(\overline{bc}, \overline{ae}) - 1$  is odd. According to Proposition 1,  $\pi(m, \overline{gk})$  and  $\pi(k, \overline{bc})$  are not in  $V(G)$ . By Proposition 2, since  $\delta(a, e)$  is even,  $\rho(\overline{hk}, \overline{mg}) = \rho(\overline{aa'}, \overline{ee'}) - \rho(\overline{aa'}, \overline{hk}) + \rho(\overline{ee'}, \overline{mg}) = \rho(\overline{aa'}, \overline{ee'}) - 1$  is odd. Again by Proposition 1,  $\pi(m, \overline{hk})$  and  $\pi(k, \overline{mg})$  do not belong to  $V(G)$ . Due to  $\angle abc > \frac{\pi}{2}$  and  $k' \in ab$ , we have  $\angle k'k'm > \frac{\pi}{2}$  and  $\pi(m, \overline{kk'}) \notin V(G)$ . Since  $\delta(a, b) \leq \delta(a, e) - 4$ ,  $\delta(m, b) \geq 4 + 1 = 5$  and  $\rho(\overline{mm'}, \overline{kk'}) \geq 4$ . Thus,  $\pi(k, \overline{mm'}) \notin V(G)$ . Therefore, there is no poidge in  $G$  containing  $m$  and  $k$ , a contradiction. Hence,  $\delta(a, b) \geq \delta(a, e) - 2$ .

Now we claim that  $\delta(a, b) \neq \delta(a, e) - 2$ . Suppose, contrary to our claim, that  $\delta(a, b) = \delta(a, e) - 2$ . According to Lemma 7,  $\delta(d, e) = 1$  or  $\delta(d, e) = 2$ .

Case 1. If  $\delta(d, e) = 1$ , then  $\delta(b, c) = 3$ . Let  $p \in cd \cap V(G)$  such that  $\delta(c, p) = 1$ . Assume that  $\overline{pp'}$  is a lattice line parallel to  $ab$ , and  $\overline{pp''}$  is a lattice line parallel to  $ae$ , see Fig. 25.

Since  $\delta(a, e)$  is even and  $\delta(d, e) = 1$ ,  $\rho(\overline{aa'}, \overline{cd}) = \rho(\overline{aa'}, \overline{ee'}) + \rho(\overline{ee'}, \overline{cd}) = \rho(\overline{aa'}, \overline{ee'}) + 1$  is odd. According to Proposition 1,  $\pi(a, \overline{cd})$  and  $\pi(p, \overline{aa'})$  are not in  $V(G)$ . Because  $\delta(a, b)$  is even and  $\delta(c, p) = 1$ ,  $\rho(\overline{pp''}, \overline{ae}) = \rho(\overline{bc}, \overline{ae}) - \rho(\overline{bc}, \overline{pp''}) = \rho(\overline{bc}, \overline{ae}) - 1$  is odd. Again by Proposition 1,  $\pi(a, \overline{pp''})$  and  $\pi(p, \overline{ae})$  do not belong to  $V(G)$ . Due to  $\angle bae < \frac{\pi}{2}$ ,  $\pi(a, \overline{pp'})$  does not belong to  $V(G)$ . Because  $\angle bae > \frac{\pi}{2}$  and  $\delta(c, p) = 1$ ,  $\pi(p, \overline{ab})$  is not a vertex of  $G$ . Therefore, there is no poidge in  $G$  containing  $a$  and  $p$ , a contradiction.

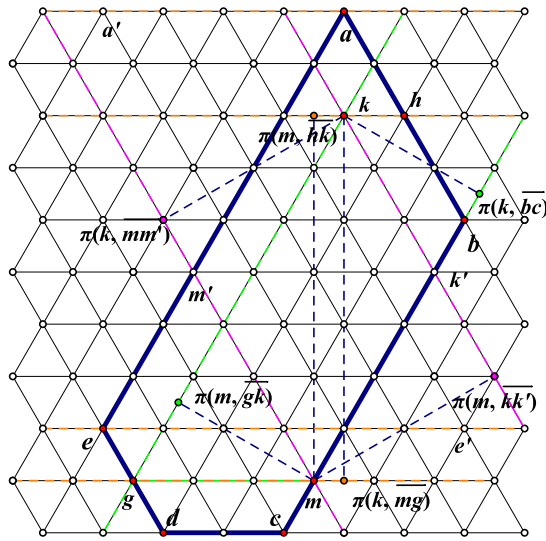


Fig. 24.  $\delta(a, b) \leq \delta(a, e) - 4$ .

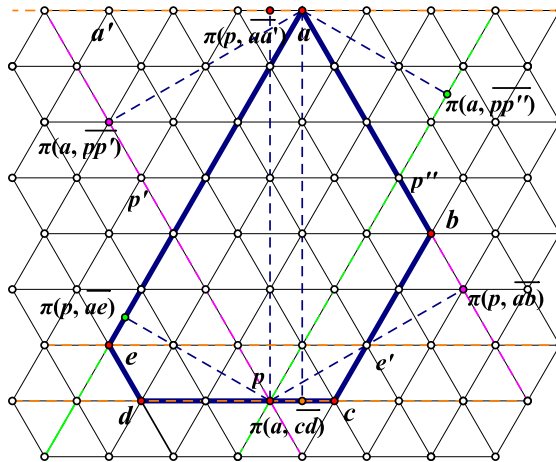


Fig. 25. Case 1.

Case 2. If  $\delta(d, e) = 2$ , then  $\delta(b, c) = 4$ . Since  $\delta(d, e) = 2$  and  $\delta(a, b) = \delta(a, e) - 2$ ,  $\delta(c, d) \geq 2$ . Let  $s \in ab \cap V(G)$  and  $t \in cd \cap V(G)$  such that  $\delta(a, s) = 1$  and  $\delta(c, t) = 2$ . Assume that  $ss'$  is a lattice line parallel to  $cd$ ,  $tt'$  is a lattice line parallel to  $ab$ , and  $ss''$  and  $tt''$  are lattice lines parallel to  $bc$ , see Fig. 26.

Since  $\delta(a, s) = 1$ ,  $\rho(\overline{aa'}, \overline{ss'}) = 1$  and  $\rho(\overline{ss''}, \overline{ae}) = 1$ . Since  $\delta(d, e) = 2$  and  $\delta(a, e)$  is even,  $\rho(\overline{ss'}, \overline{cd}) = \rho(\overline{aa'}, \overline{ee'}) + \rho(\overline{ee'}, \overline{cd}) - \rho(\overline{aa'}, \overline{ss'}) = \rho(\overline{aa'}, \overline{ee'}) + 1$  is odd. According to Proposition 1,  $\pi(s, \overline{cd})$  and  $\pi(t, \overline{ss'})$  are not in  $V(G)$ . Since  $\delta(c, t) = 2$  and  $\delta(a, b)$  is even,  $\rho(\overline{ss''}, \overline{tt''}) = \rho(\overline{bc}, \overline{ae}) - \rho(\overline{bc}, \overline{tt''}) - \rho(\overline{ss''}, \overline{ae}) = \rho(\overline{bc}, \overline{ae}) - 3$  is odd. Again by Proposition 1,  $\pi(s, \overline{tt''})$  and  $\pi(t, \overline{ss''})$  do not belong to  $V(G)$ . Because  $\widehat{bae}$  is acute and  $\delta(a, s) = 1$ ,  $\pi(s, \overline{tt'}) \notin V(G)$ . Since  $\widehat{bcd}$  is obtuse and  $\delta(c, t) = 2$ ,  $\pi(t, \overline{ab}) \notin V(G)$ . Thus, there is no poidge in  $G$  containing  $s$  and  $t$ , a contradiction.

Hence,  $\delta(a, b) = \delta(a, e)$ .  $\square$

**Theorem 3.** Let  $G$  be a grid graph such that set( $C$ ) is a convex pentagon  $[abcde]$  and the angle at  $a$  is acute. Then  $G$  is poidge-convex if and only if  $\delta(a, b) = \delta(a, e)$  is even and  $1 \leq \delta(d, e) \leq 2$ .

**Proof.** Suppose without loss of generality that  $\delta(a, e) \geq \delta(a, b)$ . According to Lemma 6, Lemma 7 and Lemma 8, the necessity of the condition in the statement is obvious. Now we prove its sufficiency.

Let  $x, y$  belong to  $V(G)$ . We prove the existence of a poidge in  $G$  containing both  $x$  and  $y$ .

Case 1. There exists a lattice line  $L$  containing  $x$  and  $y$ .

Subcase 1.1.  $L$  parallel to  $cd$ . Let  $L \cap (ab \cup bc) = \{u\}$  and  $L \cap (ae \cup ed) = \{v\}$ . If  $u = b$ , then  $\underline{L} \cap \text{set}(C) = \{b, e\}$  and  $be$  is a linear path in  $G$  containing  $x$  and  $y$ . Since  $\delta(a, e)$  is even and the angles  $\widehat{bae}$  and  $\widehat{abe}$  are acute,  $\pi(e, \overline{ab}) \in V(G)$ . Thus,  $\{\pi(e, \overline{ab})\} \cup be$  is a

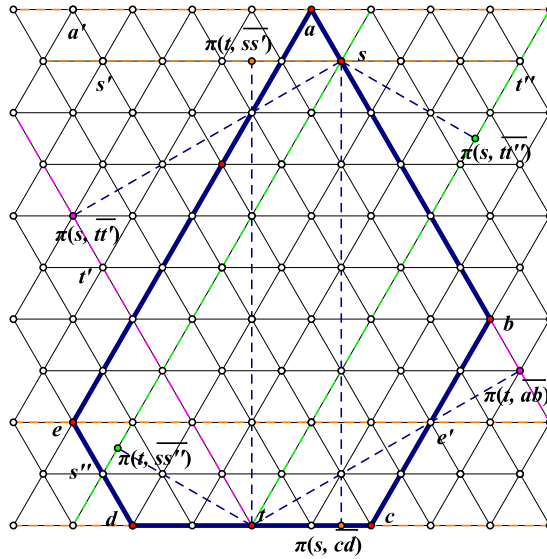


Fig. 26. Case 2.

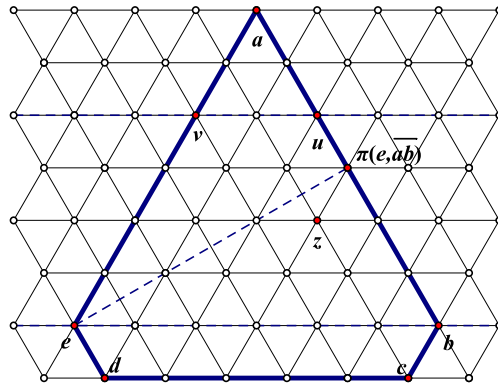


Fig. 27. Subcase 1.1.

poidge in  $G$  containing  $x$  and  $y$ , see Fig. 27. In case  $u \neq b$ , let  $P$  be the linear path between  $u$  and  $v$ . It is easily seen that there exists  $z \in V(G)$  such that  $\delta(u, z) = 2$  and  $\angle zuv = \frac{\pi}{2}$ . Thus,  $\{z\} \cup P$  is a poidge in  $G$  containing  $x$  and  $y$ , see Fig. 27.

Subcase 1.2.  $L$  parallel to  $ab$ . Let  $L \cap (bc \cup cd) = \{u\}$  and  $L \cap ae = \{v\}$ ; then,  $u, v \in V(G)$ . Assume that  $P$  is the linear path between  $u$  and  $v$ .

If  $u \in bc$ , then  $\{\pi(u, \overline{ae})\} \cup P$  is a poidge in  $G$  containing  $x$  and  $y$ , because  $\rho(\overline{bc}, \overline{ae})$  is even, which implies  $\pi(u, \overline{ae}) \in V(G)$ , see Fig. 28.

If  $u \in (cd \setminus \{c\})$ , then there exists a vertex  $z \in V(G)$  such that  $\delta(u, z) = 2$  and  $\angle zuv = \frac{\pi}{2}$ . Thus,  $\{z\} \cup P$  is a poidge in  $G$  containing  $x$  and  $y$ , see Fig. 29.

Subcase 1.3.  $L$  parallel to  $ae$ . This is symmetric to Subcase 1.2.

Case 2. There exist no lattice line containing  $x$  and  $y$ . Then there will exist two distinct parallel lattice lines  $L_1$  and  $L_2$  such that  $x \in L_1$ ,  $y \in L_2$  and  $\rho(L_1, L_2)$  is even. Let  $L_i \cap \text{set}(C) = \{u_i, v_i\}$ , then  $u_i, v_i \in V(G)$ , where  $i = 1, 2$ . Suppose without loss of generality that  $\delta(u_1, v_1) \leq \delta(u_2, v_2)$ .

Subcase 2.1. If  $\text{set}(C) \cap (L_1 \cup L_2)$  is a trapezoid or a hexagon, then  $\pi(u_1 v_1, L_2) \subset u_2 v_2$  (see Fig. 30). Since  $\rho(L_1, L_2)$  is even,  $\pi(x, L_2) \in V(G)$ , by Proposition 1. Let  $z = \pi(x, L_2)$ ,  $P$  be the linear path between  $z$  and  $u_2$ , and  $Q$  is the linear path between  $z$  and  $v_2$ , then  $y \in P$  or  $y \in Q$ . Therefore,  $\{x\} \cup P$  or  $\{x\} \cup Q$  is a poidge in  $G$  containing  $x$  and  $y$ .

Subcase 2.2. If  $\text{set}(C) \cap (L_1 \cup L_2)$  is a parallelogram, then  $\delta(b, c) = \delta(d, e) = 2$ , and  $ab \subset (L_1 \cup L_2)$  or  $cd \subset (L_1 \cup L_2)$ . Let  $\overline{aa'}$  be a lattice line parallel to  $cd$ .

Subcase 2.2.1. If  $ab \subset (L_1 \cup L_2)$ , then  $c \in (L_1 \cup L_2)$ . Assume without loss of generality that  $ab \subset L_1$ ,  $c \in L_2$  and  $L_2 \cap ae = c'$  (see Fig. 31).

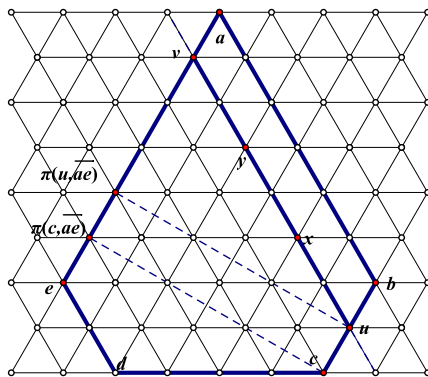


Fig. 28.  $u \in bc$ .

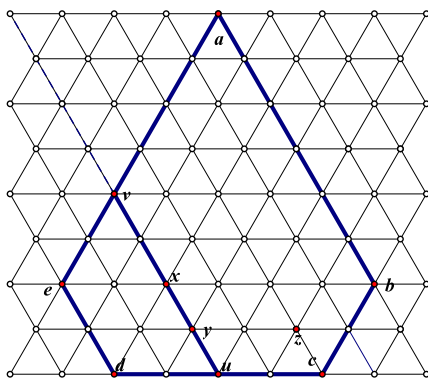


Fig. 29.  $u \in (cd \setminus \{c\})$ .

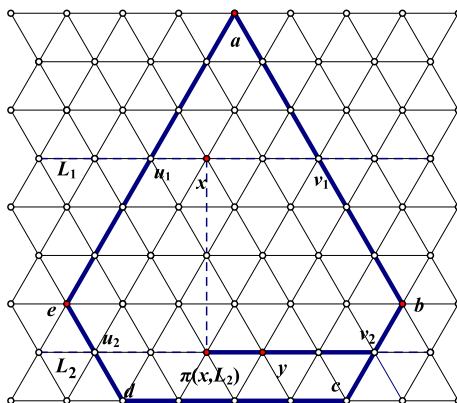


Fig. 30. Subcase 2.1.

If  $x = a$  and  $y = c$ , then  $\pi(x, L_1) \notin V(G)$  and  $\pi(y, L_2) \notin V(G)$ . Since  $\delta(a, b)$  and  $\delta(b, c)$  are even,  $\rho(\overline{aa'}, \overline{cd}) = \rho(\overline{aa'}, \overline{be}) + \rho(\overline{be}, \overline{cd})$  is even. According to Proposition 1,  $\pi(a, \overline{cd}) \in V(G)$ . Let  $P_1$  be the linear path between  $\pi(a, \overline{cd})$  and  $c$ , then  $\{a\} \cup P_1$  is a poidge in  $G$  containing  $x$  and  $y$ .

If  $x \neq a$ , then  $\pi(x, L_2) \in V(G)$ . Let  $P_2$  be the linear path between  $\pi(x, L_2)$  and  $c$ , and  $P_3$  be the linear path between  $\pi(x, L_2)$  and  $c'$ , then  $y \in P_3$  or  $y \in P_4$ . Therefore,  $\{x\} \cup P_3$  or  $\{x\} \cup P_4$  is a poidge in  $G$  containing  $x$  and  $y$ .

If  $y \neq c$ , then  $\pi(y, L_1) \in V(G)$ . Let  $P_5$  be the linear path between  $\pi(y, L_1)$  and  $a$ , and  $P_6$  be the linear path between  $\pi(y, L_1)$  and  $b$ , then  $x \in P_5$  or  $x \in P_6$ . Therefore,  $\{y\} \cup P_5$  or  $\{y\} \cup P_6$  is a poidge in  $G$  containing  $x$  and  $y$ .

Subcase 2.2.2. If  $ae \subset (L_1 \cup L_2)$ , then  $d \in (L_1 \cup L_2)$ . From the analysis of the symmetric Subcase 2.2.1, it follows that there exists a poidge in  $G$  containing  $x$  and  $y$ .  $\square$

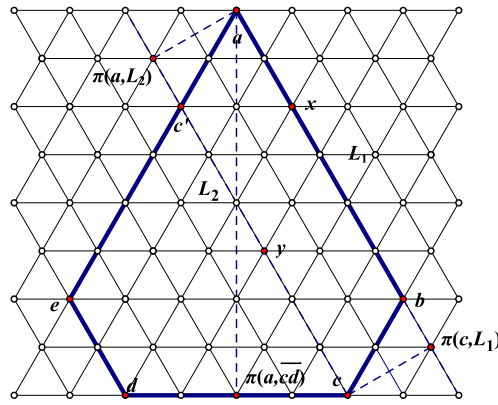


Fig. 31. Subcase 2.2.1.

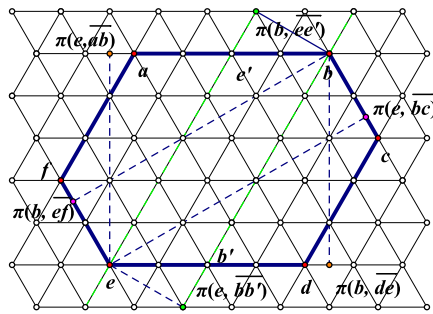


Fig. 32.  $\rho(\overline{ab}, \overline{de})$  and  $\rho(\overline{bc}, \overline{ef})$  are odd.

**Lemma 9.** Let  $G$  be a grid graph and  $\text{set}(C)$  be a convex hexagon  $[abcdef]$ . If  $G$  is poidge-convex, then at most one of the three numbers  $\rho(\overline{ab}, \overline{de})$ ,  $\rho(\overline{bc}, \overline{ef})$  and  $\rho(\overline{cd}, \overline{fa})$  is odd.

**Proof.** Suppose the assertion of the lemma is false. Assume without loss of generality that  $\rho(\overline{ab}, \overline{de})$  and  $\rho(\overline{bc}, \overline{ef})$  are odd and  $\rho(\overline{ab}, \overline{de}) \leq \rho(\overline{bc}, \overline{ef})$ . Since  $\rho(\overline{ab}, \overline{de})$  is odd,  $\pi(b, \overline{de})$  and  $\pi(e, \overline{ab})$  are not vertices of  $G$ , by Proposition 1. On the other hand,  $\rho(\overline{bc}, \overline{ef})$  is odd implies  $\pi(b, \overline{ef})$  and  $\pi(e, \overline{bc})$  do not belong to  $V(G)$ .

If  $\rho(\overline{ab}, \overline{de}) = \rho(\overline{bc}, \overline{ef})$ , then there exists a lattice line containing  $b$  and  $e$ . Because all of the four angles  $\widehat{abe}$ ,  $\widehat{cbe}$ ,  $\widehat{deb}$  and  $\widehat{feb}$  are acute, there exists no vertex of  $G$ , say  $u$ , such that  $\angle ube = \frac{\pi}{2}$  or  $\angle ueb = \frac{\pi}{2}$ . Thus, there is no poidge in  $G$  containing  $b$  and  $e$ , a contradiction.

If  $\rho(\overline{ab}, \overline{de}) < \rho(\overline{bc}, \overline{ef})$ , then there exist two distinct lattice lines parallel to  $cd$ , say  $\overline{bb'}$  and  $\overline{ee'}$ , such that  $b \in \overline{bb'}$  and  $e \in \overline{ee'}$ , where  $b'$  and  $e'$  belong to  $V(C)$ , see Fig. 32. Since the angles  $\widehat{baf}$  and  $\widehat{cde}$  are obtuse and the lattices lines  $\overline{bb'}$  and  $\overline{ee'}$  are parallel to  $cd$ ,  $\pi(b, \overline{ee'})$  and  $\pi(e, \overline{bb'})$  are not in  $V(G)$ . Thus, there is no poidge in  $G$  contains  $b$  and  $e$ , a contradiction.  $\square$

**Theorem 4.** Let  $G$  be a grid graph with  $\text{set}(C)$  a convex hexagon  $[abcdef]$  such that  $\rho(\overline{ab}, \overline{de})$  is odd. If  $G$  is poidge-convex, then  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{bc}, \overline{ef}) = \pm 1$  and  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{cd}, \overline{fa}) = \pm 1$ .

**Proof.** Since  $\rho(\overline{ab}, \overline{de})$  is odd, both  $\rho(\overline{bc}, \overline{ef})$  and  $\rho(\overline{cd}, \overline{fa})$  are even, according to Lemma 9. Let  $d_1, d_2 \in V(C)$  be the two neighbours of  $d$  and  $e_1, e_2 \in V(C)$  the two neighbours of  $e$ , such that  $d_1, e_1 \in \overline{de}$ . Suppose that  $\overline{d_1d'_1}$  is a lattice line parallel to  $cd$  and  $\overline{e_1e'_1}$  a lattice line parallel to  $bc$ . By Proposition 1, since  $\rho(\overline{ab}, \overline{de})$  is odd,  $\pi(a, \overline{de})$ ,  $\pi(b, \overline{de})$ ,  $\pi(d_1, \overline{ab})$  and  $\pi(e_1, \overline{ab})$  are not in  $V(G)$ , see Fig. 33.

Since  $\delta(d, d_1) = 1$ ,  $\rho(\overline{d_1d'_1}, \overline{fa}) = \rho(\overline{cd}, \overline{fa}) - 1$  is odd, by Proposition 2. According to Proposition 1,  $\pi(a, \overline{d_1d'_1})$  and  $\pi(d_1, \overline{fa})$  are not in  $V(G)$ . On the other hand,  $\overline{d_1d_2}$  is orthogonal to  $\overline{bc}$  and  $\overline{d_1d_2} \cap V(G) = \{d_1, d_2\}$ . By the poidge-convexity of  $G$ , there exists a poidge in  $G$  containing  $a$  and  $d_1$ , which means that  $\overline{ad_1}$  or  $\overline{ad_2}$  must be a lattice line. Similarly,  $\overline{be_1}$  or  $\overline{be_2}$  must be a lattice line.

If  $\overline{ad_1}$  is a lattice line, then  $\rho(\overline{d_1d'_1}, \overline{fa}) = \rho(\overline{ab}, \overline{de})$  and  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{cd}, \overline{fa}) = -1$ . Otherwise,  $\overline{ad_2}$  is a lattice line, which implies that  $\rho(\overline{cd}, \overline{fa}) = \rho(\overline{ab}, \overline{de_2e_1})$  and  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{cd}, \overline{fa}) = 1$ .

Symmetrically, if  $\overline{be_1}$  is a lattice line, then  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{bc}, \overline{ef}) = -1$ , and if  $\overline{be_2}$  is a lattice line, then  $\rho(\overline{ab}, \overline{de}) - \rho(\overline{bc}, \overline{ef}) = 1$ .  $\square$

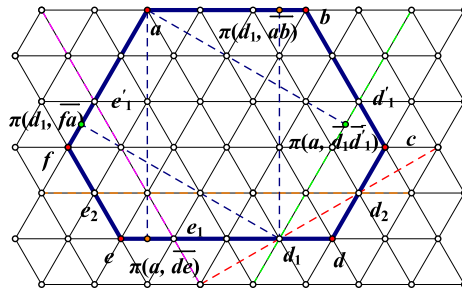


Fig. 33.  $\rho(\overline{ab}, \overline{de})$  is odd.

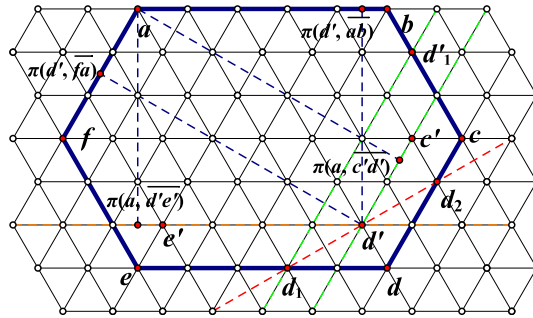


Fig. 34. Theorem 5.

**Theorem 5.** Let  $G$  be a grid graph such that  $\text{set}(C)$  is a convex hexagon  $[abcdef]$  and  $\rho(\overline{ab}, \overline{de})$ ,  $\rho(\overline{bc}, \overline{ef})$  and  $\rho(\overline{cd}, \overline{fa})$  are all even. If  $G$  is poidge-convex, then any pair of these numbers are either equal or consecutive even numbers.

**Proof.** Let  $c', d', e' \in V(G) \setminus V(C)$  such that  $\delta(c, c') = \delta(d, d') = \delta(e, e') = 1$ . Then  $\overline{c'd'}$  and  $\overline{d'e'}$  will be two lattice lines parallel to  $cd$  and  $de$ , respectively, see Fig. 34.

Let  $d_1 \in \overline{de}$  and  $d_2 \in \overline{cd}$  such that  $\delta(d, d_1) = \delta(d, d_2) = 2$ . Now we show that the poidge-convexity of  $G$  implies that one of the three lines  $\overline{ad}$ ,  $\overline{ad_1}$  and  $\overline{ad_2}$  is a lattice line. Since  $\delta(d, d') = 1$  and  $\rho(\overline{ab}, \overline{de})$  is even,  $\rho(\overline{ab}, \overline{d'e'}) = \rho(\overline{ab}, \overline{de}) - 1$  is odd. It follows that  $\pi(a, \overline{d'e'})$  and  $\pi(d', \overline{ab})$  do not belong to  $V(G)$ . By a similar argument,  $\delta(d, d') = 1$  and  $\rho(\overline{cd}, \overline{fa})$  even imply that  $\pi(a, \overline{c'd'})$  and  $\pi(d', \overline{fa})$  are not in  $V(G)$ . Obviously,  $\overline{d_1d_2}$  is orthogonal to  $bc$  and  $\overline{d_1d_2} \cap V(G) = \{d_1, d', d_2\}$ . By the poidge-convexity of  $G$ , there is a poidge in  $G$  containing  $a$  and  $d'$ , which means that one of the three lines  $\overline{ad}$ ,  $\overline{ad_1}$  and  $\overline{ad_2}$  must be a lattice line.

Suppose  $\overline{ad_1}$  is a lattice line parallel to  $cd$ . If  $\overline{ad_1}$  is a lattice line, then  $\rho(\overline{ab}, \overline{de}) = \rho(\overline{ad_1}, \overline{fa})$ . Because  $\delta(d, d_1) = 2$ ,  $\rho(\overline{cd}, \overline{fa}) = \rho(\overline{ab}, \overline{de}) + 2$ . If  $\overline{ad}$  is a lattice line, then  $\rho(\overline{ab}, \overline{de}) = \rho(\overline{cd}, \overline{fa})$ . Otherwise,  $\overline{ad_2}$  is a lattice line and  $\rho(\overline{cd}, \overline{fa}) = \rho(\overline{ab}, \overline{de}) - 2$ .

Similar arguments apply to  $b$  and  $e'$ ,  $f$  and  $c'$ , respectively; thus, if  $G$  is poidge-convex, then  $|\rho(\overline{ab}, \overline{de}) - \rho(\overline{bc}, \overline{ef})| \leq 2$  and  $|\rho(\overline{bc}, \overline{ef}) - \rho(\overline{cd}, \overline{fa})| \leq 2$ .  $\square$

We arrive at the following characterization.

**Theorem 6.** Let  $G$  be a grid graph such that  $\text{set}(C)$  is a convex hexagon, and  $p, q, r$  are the intrinsic distances of its opposite sides.  $G$  is poidge-convex if and only if at most one of the numbers  $p, q, r$  is odd, and the difference between any two of them is at most 2.

**Proof.** For the “only if” implication, combine Theorems 4 and 5. The verification of the “if” implication is a routine matter.  $\square$

To summarize the results of this section, a grid graph with boundary cycle  $C$  is poidge-convex, if and only if  $\text{set}(C)$  is an equilateral triangle with side length 2 or a pentagon, as described in Theorem 3, or a hexagon, as described in Theorem 6.

**Data availability**

No data was used for the research described in the article.

**Acknowledgements**

This work is supported by NSF of China (12271139, 11871192); the High-end Foreign Experts Recruitment Program of People’s Republic of China (G2023003003L); the Program for Foreign Experts of Hebei Province; the Natural Science Foundation of Hebei

Province (A2023205045); the Special Project on Science and Technology Research and Development Platforms, Hebei Province (22567610H) and the program for 100 Foreign Experts Plan of Hebei Province.

## References

- [1] R. Blind, G. Valette, T. Zamfirescu, Rectangular convexity, *Geom. Dedic.* 9 (1980) 317–327.
- [2] K. Böröczky Jr., Rectangular convexity of convex domains of constant width, *Geom. Dedic.* 34 (1990) 13–18.
- [3] C.C. Centeno, S. Dantas, M.C. Dourado, et al., Convex partitions of graphs induced by paths of order three, *Discret. Math. Theor. Comput. Sci.* 12 (5) (2010) 175–184.
- [4] M. Changat, J. Mathew, On triangle path convexity in graphs, *Discrete Math.* 206 (1999) 91–95.
- [5] P. Duchet, Convex sets in graphs II: minimal path convexity, *J. Comb. Theory, Ser. B* 44 (1988) 307–316.
- [6] M. Farber, R.E. Jamison, On local convexity in graphs, *Discrete Math.* 66 (1987) 231–247.
- [7] B. Wang, L. Yuan, T. Zamfirescu, Poidge-convexity, *J. Convex Anal.* 28 (2021) 1155–1170.
- [8] L. Yuan, T. Zamfirescu, Right triple convex completion, *J. Convex Anal.* 22 (2015) 291–301.
- [9] L. Yuan, T. Zamfirescu, Right triple convexity, *J. Convex Anal.* 23 (2016) 1219–1246.
- [10] T. Zamfirescu, Right convexity, *J. Convex Anal.* 21 (2014) 253–260.