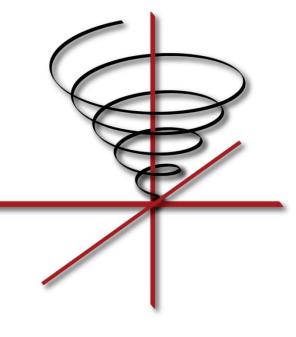
# SPIRAL, FFTX, and the Path to SpectralPACK



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In collaboration with the SPIRAL and FFTX team @ CMU and LBL

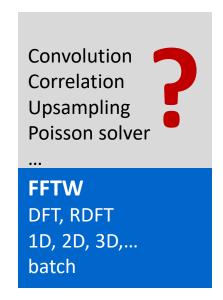


#### **Have You Ever Wondered About This?**

#### **Numerical Linear Algebra**

# LAPACK ScaLAPACK LU factorization Eigensolves SVD BLAS, BLACS BLAS-1 BLAS-2 BLAS-3

#### **Spectral Algorithms**



#### No LAPACK equivalent for spectral methods

- Medium size 1D FFT (1k—10k data points) is most common library call
  applications break down 3D problems themselves and then call the 1D FFT library
- Higher level FFT calls rarely used
   FFTW guru interface is powerful but hard to used, leading to performance loss
- Low arithmetic intensity and variation of FFT use make library approach hard
   Algorithm specific decompositions and FFT calls intertwined with non-FFT code

#### It Is Worse Than It Seems

#### FFTW is de-facto standard interface for FFT



- FFTW 3.X is the high performance reference implementation: supports multicore/SMP and MPI, and Cell processor
- Vendor libraries support the FFTW 3.X interface:
   Intel MKL, IBM ESSL, AMD ACML (end-of-life), Nvidia cuFFT, Cray LibSci/CRAFFT

#### **Issue 1: 1D FFTW call is standard kernel for many applications**

- Parallel libraries and applications reduce to 1D FFTW call P3DFFT, QBox, PS/DNS, CPMD, HACC,...
- Supported by modern languages and environments
   Python, Matlab,...

#### **Issue 2:** FFTW is slowly becoming obsolete

- FFTW 2.1.5 (still in use, 1997), FFTW 3 (2004) minor updates since then
- Development currently dormant, except for small bug fixes
- No native support for accelerators (GPUs, Xeon PHI, FPGAs) and SIMT
- Parallel/MPI version does not scale beyond 32 nodes

Risk: loss of high performance FFT standard library

## **FFTX:** The FFTW Revamp for ExaScale

#### **Modernized FFTW-style interface**



- Backwards compatible to FFTW 2.X and 3.X
   old code runs unmodified and gains substantially but not fully
- Small number of new features
   futures/delayed execution, offloading, data placement, callback kernels

#### Code generation backend using **SPIRAL**

- Library/application kernels are interpreted as specifications in DSL extract semantics from source code and known library semantics
- Compilation and advanced performance optimization cross-call and cross library optimization, accelerator off-loading,...
- Fine control over resource expenditure of optimization compile-time, initialization-time, invocation time, optimization resources
- Reference library implementation and bindings to vendor libraries
   library-only reference implementation for ease of development

# FFTX and SpectralPACK: Long Term Vision

#### **Numerical Linear Algebra**

# LAPACK LU factorization Eigensolves SVD ...

#### BLAS

BLAS-1 BLAS-2

BLAS-3



#### **Spectral Algorithms**

#### **SpectralPACK**

Convolution
Correlation
Upsampling
Poisson solver

#### **FFTX**

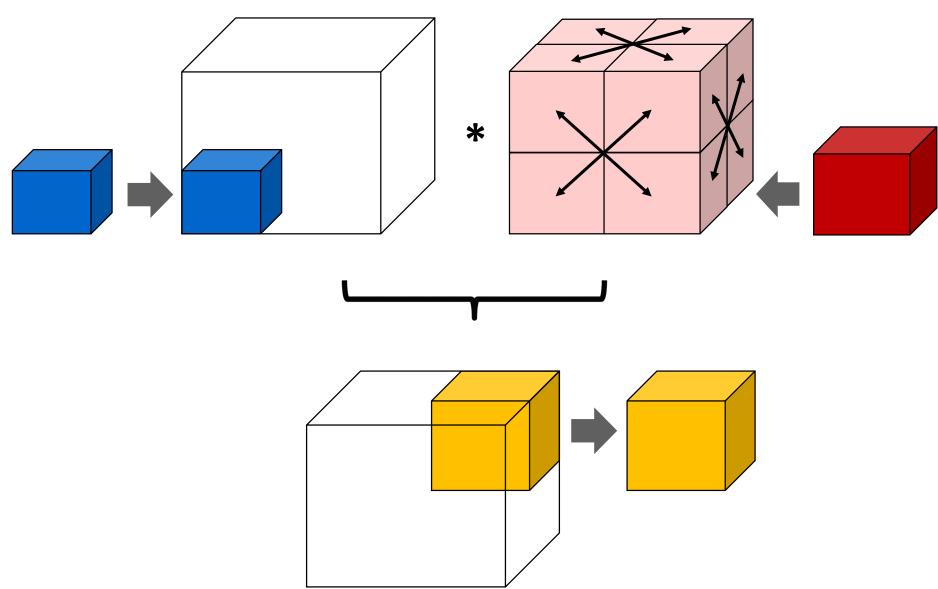
DFT, RDFT 1D, 2D, 3D,... batch

#### Define the LAPACK equivalent for spectral algorithms

- Define FFTX as the BLAS equivalent provide user FFT functionality as well as algorithm building blocks
- Define class of numerical algorithms to be supported by SpectralPACK
   PDE solver classes (Green's function, sparse in normal/k space,...), signal processing,...
- Define SpectralPACK functions
   circular convolutions, NUFFT, Poisson solvers, free space convolution,...

FFTX and SpectralPACK solve the "spectral dwarf" long term

# **Example: Hockney Free Space Convolution**



# **Example: Hockney Free Space Convolution**

```
fftx plan pruned real convolution plan(fftx real *in, fftx real *out, fftx complex *symbol,
        int n, int n in, int n out, int n freq) {
    int rank = 1,
    batch rank = 0,
    fftx plan plans[5];
    fftx plan p;
    tmp1 = fftx create zero temp real(rank, &padded dims);
   plans[0] = fftx plan guru copy real(rank, &in dimx, in, tmp1, MY FFTX MODE SUB);
    tmp2 = fftx create temp complex(rank, &freq dims);
   plans[1] = fftx plan guru dft r2c(rank, &padded dims, batch rank,
        &batch dims, tmp1, tmp2, MY FFTX MODE SUB);
    tmp3 = fftx create temp complex(rank, &freq dims);
    plans[2] = fftx plan guru pointwise c2c(rank, &freq dimx, batch rank, &batch dimx,
        tmp2, tmp3, symbol, (fftx callback) complex scaling,
        MY FFTX MODE SUB | FFTX PW POINTWISE);
    tmp4 = fftx create temp real(rank, &padded dims);
   plans[3] = fftx plan guru dft c2r(rank, &padded dims, batch rank,
        &batch dims, tmp3, tmp4, MY FFTX MODE SUB);
   plans[4] = fftx plan guru copy real(rank, &out dimx, tmp4, out, MY FFTX MODE SUB);
   p = fftx plan compose(numsubplans, plans, MY FFTX MODE TOP);
    return p;
```

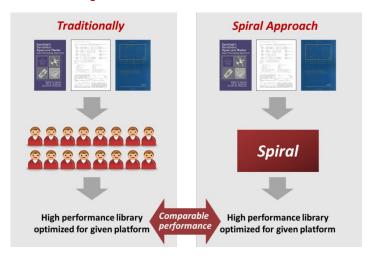
Looks like FFTW calls, but is a specification for SPIRAL



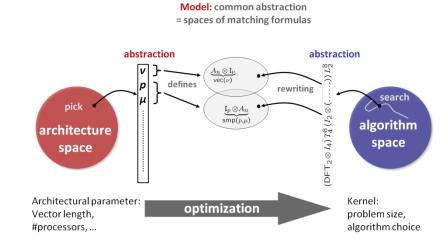


# Spiral Technology in a Nutshell

#### **Library Generator**



#### **Mathematical Foundation**



#### **Performance Portability**



Intel Xeon 8180M 2.25 Tflop/s, 205 W 28 cores, 2.5—3.8 GHz 2-way—16-way AVX-512



**IBM POWER9 768 Gflop/s, 300 W** 24 cores, 4 GHz 4-way VSX-3



Nvidia Tesla V100 7.8 Tflop/s, 300 W 5120 cores, 1.2 GHz 32-way SIMT



Intel Xeon Phi 7290F 1.7 Tflop/s, 260 W 72 cores, 1.5 GHz 8-way/16-way LRBni



**Snapdragon 835 15 Gflop/s, 2 W**8 cores, 2.3 GHz
A540 GPU, 682 DSP, NEON



Intel Atom C3858 32 Gflop/s, 25 W 16 cores, 2.0 GHz 2-way/4-way SSSE3

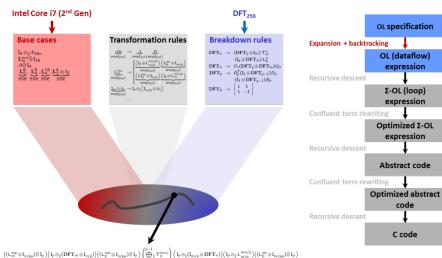


**Dell PowerEdge R940 3.2 Tflop/s, 6 TB, 850 W**4x 24 cores, 2.1 GHz
4-way/8-way AVX



**Summit** 187.7 Pflop/s, 13 MW 9,216 x 22 cores POWER9 + 27,648 V100 GPUs

#### **Code Synthesis and Autotuning**



# Algorithms: Rules in Domain Specific Language

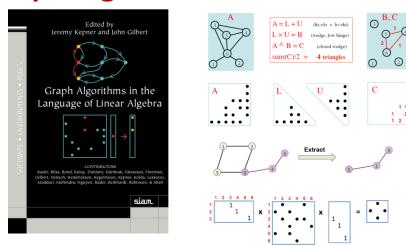
#### **Linear Transforms**

$$\begin{array}{lll} \operatorname{DFT}_n & \to & (\operatorname{DFT}_k \otimes \operatorname{I}_m) \operatorname{T}_m^n (\operatorname{I}_k \otimes \operatorname{DFT}_m) \operatorname{L}_k^n, & n = km \\ \operatorname{DFT}_n & \to & P_n (\operatorname{DFT}_k \otimes \operatorname{DFT}_m) Q_n, & n = km, \operatorname{gcd}(k,m) = 1 \\ \operatorname{DFT}_p & \to & R_p^T (\operatorname{I}_1 \oplus \operatorname{DFT}_{p-1}) D_p (\operatorname{I}_1 \oplus \operatorname{DFT}_{p-1}) R_p, & p \operatorname{ prime} \\ \operatorname{DCT-3}_n & \to & (\operatorname{I}_m \oplus \operatorname{J}_m) \operatorname{L}_m^n (\operatorname{DCT-3}_m (1/4) \oplus \operatorname{DCT-3}_m (3/4)) \\ & & \cdot (\operatorname{F}_2 \otimes \operatorname{I}_m) \begin{bmatrix} \operatorname{I}_m & \operatorname{O} \oplus - \operatorname{J}_{m-1} \\ \frac{1}{\sqrt{2}} (\operatorname{I}_1 \oplus \operatorname{2} \operatorname{I}_m) \end{bmatrix}, & n = 2m \\ \operatorname{DCT-4}_n & \to & S_n \operatorname{DCT-2}_n \operatorname{diag}_{0 \leq k < n} (1/(2 \operatorname{cos}((2k+1)\pi/4n))) \\ \operatorname{IMDCT}_{2m} & \to & (\operatorname{J}_m \oplus \operatorname{I}_m \oplus \operatorname{I}_m \oplus \operatorname{J}_m) \left( \left( \begin{bmatrix} 1 \\ -1 \end{bmatrix} \otimes \operatorname{I}_m \right) \oplus \left( \begin{bmatrix} -1 \\ -1 \end{bmatrix} \otimes \operatorname{I}_m \right) \right) \operatorname{J}_{2m} \operatorname{DCT-4}_{2m} \\ \operatorname{WHT}_{2^k} & \to & \prod_{i=1}^t (\operatorname{I}_{2^{k_1 + \dots + k_{i-1}}} \otimes \operatorname{WHT}_{2^{k_i}} \otimes \operatorname{I}_{2^{k_{i+1} + \dots + k_t}}), & k = k_1 + \dots + k_t \\ \operatorname{DFT}_2 & \to & \operatorname{F}_2 \\ \operatorname{DCT-2}_2 & \to & \operatorname{diag}(1, 1/\sqrt{2}) \operatorname{F}_2 \\ \operatorname{DCT-4}_2 & \to & \operatorname{J}_2 \operatorname{R}_{13\pi/8} \end{array}$$

#### **Numerical Linear Algebra**



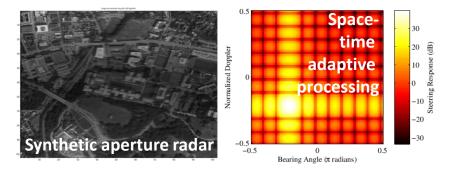
#### **Graph Algorithms**



In collaboration with CMU-SEI

#### **Spectral Domain Applications**





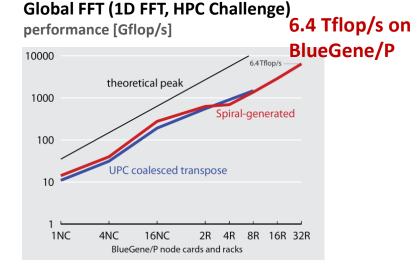
# **SPIRAL: Success in HPC/Supercomputing**

- NCSA Blue Waters
   PAID Program, FFTs for Blue Waters
- RIKEN K computer
  FFTs for the HPC-ACE ISA
- LANL RoadRunner
  FFTs for the Cell processor
- PSC/XSEDE BridgesLarge size FFTs
- LLNL BlueGene/L and P FFTW for BlueGene/L's Double FPU
- ANL BlueGene/Q Mira
   Early Science Program, FFTW for BGQ QPX









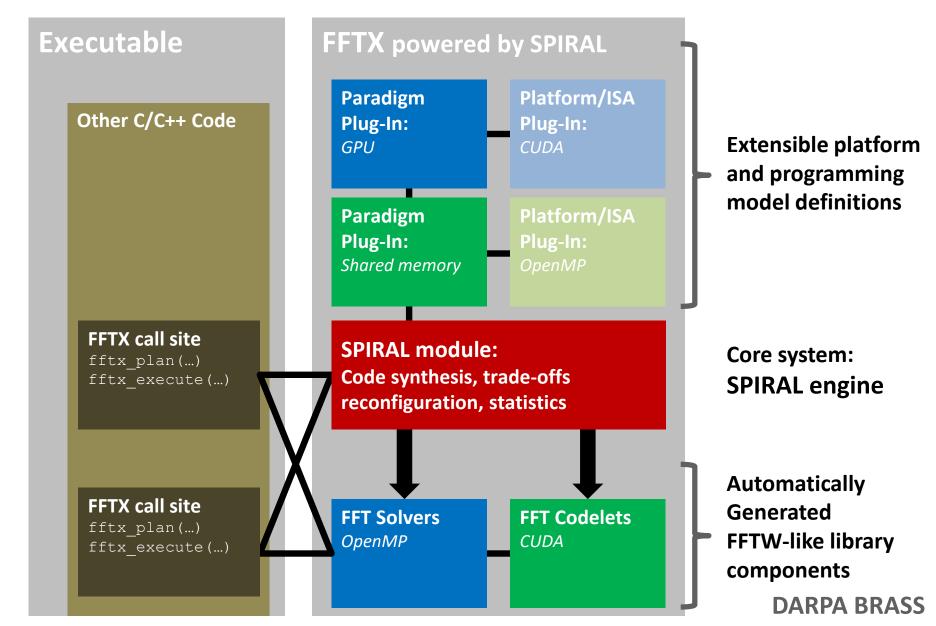
BlueGene/P at Argonne National Laboratory 128k cores (quad-core CPUs) at 850 MHz





2006 Gordon Bell Prize (Peak Performance Award) with LLNL and IBM 2010 HPC Challenge Class II Award (Most Productive System) with ANL and IBM

### **FFTX Backend: SPIRAL**



# FFTX: First Results for Hockney on Volta

FFTX and SpectralPack: A First Look

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Abstract-We propose FFTX, a new framework for building high-performance FFT-based applications on exascale machines. Complex node architectures lead to multiple levels of parallelism and demand efficient ways of data communication. The current FFTW interface falls short in maximizing performance in such scenarios. FFTX is designed to enable application developers to leverage expert-level, automatic optimizations while navigating a familiar interface. FFTX is backwards compatible to FFTW and extends the FFTW Interface into an embedded Domain Specific Language (DSL) expressed as a library interface. By means of a SPIRAL-based back end, this enables build-time source-to-source translation and advanced performance optimizations, such as cross-library calls optimizations, targeting of accelerators through offloading, and inlining of user-provided kernels. We demonstrate the use of FFTX with the prototypical example of 1D and 3D pruned convolutions and discuss future extensions.

Keywords-FFT; exasclale; code generator; high-performance;

#### I. INTRODUCTION

The Discrete Fourier Transform (DFT) — and in particular its implementation using Fast Fourier Transform (FFT) [1], [2] algorithms — is a fundamental component for the design of scientific applications suitable for the emerging exascale computing ecosystem. Application domains include material science, chemistry, molecular dynamics, and cosmology. Some examples are FFT-based differential equation solvers to compute properties of materials such as stress and strain [3], simulation of incompressible flows [4], plane wave based electronic structure methods [5] [6] and particlein-cell (PIC) codes [7].

Capturing application-specific properties of the FFTs is very important to enable high-performance execution on emerging platforms. For example, applications where a large subset of the inputs or outputs is either set to zero or not computed at all should exploit the zero structures to reduce data movements within the memory hierarchy on a node or across the network. The structure of multidimensional FFTs provide opportunities for parallelism at multiple levels — SIMD, threads, accelerators, and distributed systems. The best strategy for exploiting these opportunities strongly depends on the details of the use case and of the computing platform, as well as tradeoffs with other needs of the application in which the transforms are embedded.

Conventional FFT-based implementation approach and its limits. The implementation strategy for most of today's large science applications that depend on FFTs constists in transforming multidimensional problems into a sequence of 1D FFT calls, with the latter being performed by a library. Over the last decade or so the FFTW API became the de-facto standard FFT interface [8], [9], [10]. Vendors that provide FFT libraries like Intel, Cray, and IBM may still provide their own proprietary interface for backwards compatibility reasons, but all current vendor high-performance libraries, including Intel MKL [11], IBM ESSL [12], and Nvidia cuFFT [13], implement (at least a subset) of the FFTW interface. Thus, FFTW defines the standard FFT library interface and oftentimes is considered a key component of today's applications.

However, the approach of building up high-performance implementations out of calls to 1D FFTW kernels is breaking down on the current and emerging HPC platforms, for two reasons: First, node architectures have become more complex. Multiple cores and accelerators lead to multiple levels of parallelism, including threading and SIMD/SIMT. In addition, there are on-node complex memory hierarchies that are to varying extents user-controlled, and across which it is expensive to move data. This leads to more complex mappings of multidimensional FFT-based applications to the core 1D FFTs. Some of these are simply not expressible in the current FFTW interface; others can be expressed, but with significant programming effort, and often below the theoretically-predicted performance due to unexpected and opaque behavior of the FFT library software. Second, FFTW is no longer supported. The system comprises a high-level domain-specific language, a symbolic transformation/code

FFTX with SPIRAL and OpenACC: on par with cuFFT expert interface



FFTX with SPIRAL and OpenACC: 15 % faster than cuFFT expert interface



F. Franchetti, D. G. Spampinato, A. Kulkarni, D. T. Popovici, T. M. Low, M. Franusich, A. Canning, P. McCorquodale, B. Van Straalen, P. Colella:

**FFTX and SpectralPack: A First Look**, Workshop on Parallel Fast Fourier Transforms (PFFT), *to appear*.

http://www.spiral.net/doc/fftx

# **SPIRAL 8.0: Available Under Open Source**

- Open Source SPIRAL available
  - non-viral license (BSD)
  - Initial version, effort ongoing to open source whole system
  - Commercial support via SpiralGen, Inc.
- Developed over 20 years
  - Funding: DARPA (OPAL, DESA, HACMS, PERFECT, BRASS), NSF, ONR, DoD HPC, JPL, DOE, CMU SEI, Intel, Nvidia, Mercury
- Open sourced under DARPA PERFECT

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