Quotient inductive-inductive types and higher friends

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Motivation

Type theory in type theory:

- simple inductive types (ITs):
  - Abel–Öhman–Vezzosi, POPL 2018
- inductive-inductive types (IITs, Nordvall Forsberg PhD 2013):
  - Chapman: Type theory should eat itself, ENTCS 2009
- quotient inductive-inductive types (QIITs, this talk):
  - Altenkirch–Kaposi, POPL 2016

Other examples:

- real numbers (HoTT book)
- ordinal numbers (Lumsdaine–Shulman, 2019)
- partiality monad (Altenkirch–Danielsson–Kraus, FoSSaCS 2017)

Simple language of dependent types as a QIIT

$$
\begin{aligned}
\mathsf{Con} &: \mathsf{Set} \\
\mathsf{Ty} &: \mathsf{Con} \to \mathsf{Set} \\
\bullet &: \mathsf{Con} \\
- \triangleright - &: (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\,\Gamma \to \mathsf{Con} \\
\mathsf{U} &: (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\,\Gamma \\
\mathsf{El} &: (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\,(\Gamma \triangleright \mathsf{U}\,\Gamma) \\
\Sigma &: (A : \mathsf{Ty}\,\Gamma) \to \mathsf{Ty}\,(\Gamma \triangleright A) \to \mathsf{Ty}\,\Gamma \\
\Sigma\triangleright &: \Gamma \triangleright A \triangleright B = \Gamma \triangleright \Sigma\,A\,B
\end{aligned}
$$

Simple language of dependent types as IITs

$$
\begin{aligned}
&\mathsf{Con} && : \mathsf{Set} \\
&\mathsf{Ty} && : \mathsf{Con} \to \mathsf{Set} \\
&\mathsf{Con}_\sim && : \mathsf{Con} \to \mathsf{Con} \to \mathsf{Set} \\
&\mathsf{Ty}_\sim && : \mathsf{Con}_\sim\, \Gamma\, \Gamma' \to \mathsf{Ty}\, \Gamma \to \mathsf{Ty}\, \Gamma' \to \mathsf{Set} \\
&\bullet && : \mathsf{Con} \\
&- \rhd - && : (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\, \Gamma \to \mathsf{Con} \\
&\mathsf{U} && : (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\, \Gamma \\
&\mathsf{El} && : (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\, (\Gamma \rhd \mathsf{U}\, \Gamma) \\
&\Sigma && : (A : \mathsf{Ty}\, \Gamma) \to \mathsf{Ty}\, (\Gamma \rhd A) \to \mathsf{Ty}\, \Gamma \\
&\Sigma\rhd && : \mathsf{Con}_\sim\, (\Gamma \rhd A \rhd B)\, (\Gamma \rhd \Sigma\, A\, B) \\
&\bullet_\sim && : \mathsf{Con}_\sim\, \bullet\, \bullet \\
&\rhd_\sim && : (\overline{\Gamma} : \mathsf{Con}_\sim\, \Gamma\, \Gamma') \to \mathsf{Ty}_\sim\, \overline{\Gamma}\, A\, A' \to \mathsf{Con}_\sim\, (\Gamma \rhd A)\, (\Gamma' \rhd A') \\
&\ldots
\end{aligned}
$$

Simple language of dependent types as ITs

$$\Gamma ::= \bullet \mid \Gamma \rhd A \qquad\qquad A, B ::= \mathsf{U}\,\Gamma \mid \mathsf{El}\,\Gamma \mid \Sigma\,A\,B$$

$$\boxed{\vdash \Gamma} \qquad \boxed{\Gamma \vdash A} \qquad \boxed{\Gamma \sim \Gamma'} \qquad \boxed{\Gamma \vdash A \sim A'}$$

$$\frac{}{\vdash \bullet} \qquad \frac{\vdash \Gamma \qquad \Gamma \vdash A}{\vdash \Gamma \rhd A} \qquad \frac{\vdash \Gamma}{\Gamma \vdash \mathsf{U}\,\Gamma} \qquad \frac{\vdash \Gamma}{\Gamma \rhd \mathsf{U}\,\Gamma \vdash \mathsf{El}\,\Gamma}$$

$$\frac{\vdash \Gamma \qquad \Gamma \vdash A \qquad \Gamma \rhd A \vdash B}{\Gamma \vdash \Sigma\,A\,B}$$

$$\frac{\vdash \Gamma \qquad \Gamma \vdash A \qquad \Gamma \rhd A \vdash B}{\Gamma \rhd A \rhd B \sim \Gamma \rhd \Sigma\,A\,B} \qquad \frac{}{\bullet \sim \bullet} \qquad \frac{\Gamma \sim \Gamma' \qquad \Gamma \vdash A \sim A'}{\Gamma \rhd A \sim \Gamma' \rhd A'}$$

$$\frac{\Gamma \sim \Gamma'}{\Gamma \vdash \mathsf{U}\,\Gamma \sim \mathsf{U}\,\Gamma'} \qquad \ldots \qquad \frac{\Gamma \sim \Gamma' \qquad \Gamma \vdash A}{\Gamma' \vdash A} \qquad \ldots$$

Contents

- Formal specification of closed IITs
- Extension to QIITs
- Initial algebras
- HIITs
- Higher order abstract syntax (syntax with binding)

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How do we specify a QIIT in Agda?

```
data Nat  : Set where
     zero : Nat
     suc  : Nat → Nat

data Int  : Set where
     zero : Int
     suc  : Int → Int
     pred : Int → Int
     β    : ∀{n} → pred (suc  n) ≡ n
     η    : ∀{n} → suc  (pred n) ≡ n

data Con : Set
data Ty  : Con → Set

_▷'_ : (Γ : Con) → Ty Γ → Con
Σ'   : {Γ : Con}(A : Ty Γ) → Ty (Γ ▷' A) → Ty Γ

data Con   where
     •   : Con
     _▷_ : (Γ : Con) → Ty Γ → Con
     Σ▷  : ∀{Γ A B} → Γ ▷' A ▷' B ≡ Γ ▷' Σ' A B
data Ty    where
     U   : {Γ : Con} → Ty Γ
     El  : {Γ : Con} → Ty (Γ ▷ U)
     Σ   : {Γ : Con}(A : Ty Γ) → Ty (Γ ▷ A) → Ty Γ

_▷'_ = _▷_
Σ'   = Σ
```

Theory of closed IIT signatures

A signature is a context in a type theory (Carette–O'Connor, 2012).
Theory of signatures (ToS): category with families (CwF)

$$
\begin{aligned}
&\mathsf{Con} : \mathsf{Set} && \mathsf{Ty} : \mathsf{Con} \to \mathsf{Set} \\
&\mathsf{Sub} : \mathsf{Con} \to \mathsf{Con} \to \mathsf{Set} && \mathsf{Tm} : (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\,\Gamma \to \mathsf{Set} \\
&-[-] : \mathsf{Ty}\,\Delta \to \mathsf{Sub}\,\Gamma\,\Delta \to \mathsf{Ty}\,\Gamma && \ldots
\end{aligned}
$$

with a universe:

$$\mathsf{U} : \mathsf{Ty}\,\Gamma \qquad \mathsf{El} : \mathsf{Tm}\,\Gamma\,\mathsf{U} \to \mathsf{Ty}\,\Gamma,$$

Π types with small domain:

$$
\begin{aligned}
&\Pi : (a : \mathsf{Tm}\,\Gamma\,\mathsf{U}) \to \mathsf{Ty}\,(\Gamma \rhd \mathsf{El}\,a) \to \mathsf{Ty}\,\Gamma \\
&- \mathbin{@} - : \mathsf{Tm}\,\Gamma\,(\Pi\,a\,B) \to (u : \mathsf{Tm}\,\Gamma\,(\mathsf{El}\,a)) \to \mathsf{Tm}\,\Gamma\,(B[\mathsf{id}, u]),
\end{aligned}
$$

We will add more type formers for *open* and QIITs.

Closed IIT signatures: examples $((a \Rightarrow B) := \Pi\, a\, (B[\mathsf{p}]))$

- $\triangleright \quad \mathsf{U} \triangleright \quad \mathsf{El}\,\mathsf{q} \triangleright \quad \mathsf{q}[\mathsf{p}] \Rightarrow \mathsf{El}\,(\mathsf{q}[\mathsf{p}])$
- $\triangleright N : \mathsf{U} \triangleright \mathit{zero} : \mathsf{El}\, N \triangleright \mathit{suc} : N \Rightarrow \mathsf{El}\, N$
- $\triangleright$

$Con : \mathsf{U} \triangleright$

$Ty : Con \Rightarrow \mathsf{U} \triangleright$

$empty : \mathsf{El}\, Con \triangleright$

$ext : \Pi(\Gamma : Con). Ty \mathbin{@} \Gamma \Rightarrow \mathsf{El}\, Con \triangleright$

$U : \Pi(\Gamma : Con).\mathsf{El}\,(Ty \mathbin{@} \Gamma) \triangleright$

$El : \Pi(\Gamma : Con).\mathsf{El}\,(Ty \mathbin{@} (ext \mathbin{@} \Gamma \mathbin{@} (U \mathbin{@} \Gamma))) \triangleright$

$\Sigma : \Pi(\Gamma : Con).\Pi(A : Ty \mathbin{@} \Gamma). Ty \mathbin{@} (ext \mathbin{@} \Gamma \mathbin{@} A) \Rightarrow \mathsf{El}\,(Ty \mathbin{@} \Gamma)$

Strict positivity is enforced.

Isn't this circular?

(Q)IIT signatures are defined using a type theory, but this type theory is itself a QIIT.

We can bootstrap ToS using Church encoding (Awodey–Frey–Speight, LICS 2018).

Closed IIT signatures: semantics (i)

If $\mathcal{C}$ is a CwF, in $\hat{\mathcal{C}}$ we have (2-level type theory, Altenkirch–Capriotti–Kraus 2016, Annenkov–Capriotti–Kraus–Sattler, 2019):

$\mathsf{U}^\circ : \mathsf{Ty}_{\hat{\mathcal{C}}}\,\Gamma$ interpreted $|\mathsf{U}^\circ|_I\,\gamma := \mathsf{Ty}_{\mathcal{C}}\,I$

$\mathsf{El}^\circ : \mathsf{Tm}_{\hat{\mathcal{C}}}\,\Gamma\,\mathsf{U}^\circ \to \mathsf{Ty}_{\hat{\mathcal{C}}}\,\Gamma$ $|\mathsf{El}^\circ\,a|_I\,\gamma := \mathsf{Tm}_{\mathcal{C}}\,I\,(|a|_I\,\gamma)$

$\Pi^\circ : (a^\circ : \mathsf{Tm}_{\hat{\mathcal{C}}}\,\Gamma\,\mathsf{U}^\circ) \to \mathsf{Ty}_{\hat{\mathcal{C}}}\,(\Gamma \rhd \mathsf{El}^\circ\,a^\circ) \to \mathsf{Ty}_{\hat{\mathcal{C}}}\,\Gamma$

$|\Pi^\circ\,a^\circ\,B|_I\,\gamma := |B|_{I \rhd_{\mathcal{C}} |a|_I\,\gamma}(\gamma\mathsf{p}, \mathsf{q})$

If $\mathcal{C}$ has Id types, U° is closed under Id.

Π Π

Closed IIT signatures: semantics (ii)

We use Agda syntax to work in $\hat{\mathcal{C}}$.

$$
\begin{aligned}
&\mathsf{U}^\circ \ : \mathsf{Set} && (\mathsf{Ty}_\mathcal{C})\\
&\mathsf{El}^\circ : \mathsf{U}^\circ \to \mathsf{Set} && (\mathsf{Tm}_\mathcal{C})\\
&\Pi^\circ \ : (a^\circ : \mathsf{U}^\circ) \to (\mathsf{El}^\circ\, a^\circ \to \mathsf{Set}) \to \mathsf{Set} && (\triangleright_\mathcal{C})
\end{aligned}
$$

We define the standard model of ToS:

$$
\begin{aligned}
\mathsf{Con} &:= \mathsf{Set}\\
\mathsf{Ty}\,\Gamma &:= \Gamma \to \mathsf{Set}\\
\mathsf{Tm}\,\Gamma\,A &:= (\gamma : \Gamma) \to A\,\gamma\\
\mathsf{U}\,\gamma &:= \mathsf{U}^\circ\\
\mathsf{El}\,a\,\gamma &:= \mathsf{El}^\circ\,(a\,\gamma)\\
\Pi\,a\,B\,\gamma &:= \Pi^\circ\,(a\,\gamma)\,(B\,(\gamma, -))
\end{aligned}
$$

Example

Given the signature

$$\bullet \rhd \mathsf{U} \rhd \mathsf{El}\,\mathsf{q} \rhd \big(\mathsf{q}[\mathsf{p}] \Rightarrow \mathsf{El}\,(\mathsf{q}[\mathsf{p}])\big) : \mathsf{Con},$$

in the standard model this is

$$(N : \mathsf{U}^\circ) \times (\mathsf{El}^\circ\, N) \times (N \Rightarrow^\circ \mathsf{El}^\circ\, N) : \mathsf{Set}$$

which is a presheaf over $\mathcal{C}$, and interpreting it at the empty context of $\mathcal{C}$, we get

$$(N : \mathsf{Ty}_{\mathcal{C}}\, \bullet) \times \mathsf{Tm}_{\mathcal{C}}\ \bullet\ N \times \mathsf{Tm}_{\mathcal{C}}\,(\bullet \rhd N)\,(N[\mathsf{p}])$$

Closed IIT signatures: semantics (iii)

We use Agda syntax to work in $\hat{\mathcal{C}}$.

$$
\begin{aligned}
&\mathsf{U}^\circ \;:\mathsf{Set} && (\mathsf{Ty}_\mathcal{C})\\
&\mathsf{El}^\circ : \mathsf{U}^\circ \to \mathsf{Set} && (\mathsf{Tm}_\mathcal{C})\\
&\Pi^\circ \;: (a^\circ : \mathsf{U}^\circ) \to (\mathsf{El}^\circ\, a^\circ \to \mathsf{Set}) \to \mathsf{Set} && (\rhd_\mathcal{C})
\end{aligned}
$$

We can extend the standard model to the graph model:

$$
\begin{aligned}
&\mathsf{Con} &&:= (\Gamma^\mathsf{A} : \mathsf{Set}) && \times (\Gamma^\mathsf{M} : \Gamma^\mathsf{A} \to \Gamma^\mathsf{A} \to \mathsf{Set})\\
&\mathsf{U} &&:= \big(\lambda\gamma.\mathsf{U}^\circ && ,\;\; \lambda\_\; a^\circ\, a^{\circ\prime}.a^\circ \Rightarrow^\circ \mathsf{El}^\circ\, a^{\circ\prime}\big)\\
&\mathsf{El}\, a &&:= \big(\lambda\gamma.\mathsf{El}^\circ\,(a^\mathsf{A}\,\gamma) && ,\;\; \lambda\_\;\alpha\,\alpha'.(a^\mathsf{M}\,\_\;\alpha =_{\mathsf{El}^\circ\,(a\,\gamma')} \alpha')\big)\\
&\Pi\, a\, B &&:= \big(\lambda\gamma.\Pi^\circ\,(a^\mathsf{A}\,\gamma)\,(B^\mathsf{A}\,(\gamma, -)) && ,\;\; \lambda\_\; f\, f'.\Pi^\circ(x : a^\mathsf{A}\,\gamma).\\
& && && \qquad B^\mathsf{M}\,\_(f\,x)\,(f'\,(a^\mathsf{M}\,\_\; x'))\big)
\end{aligned}
$$

Example

Given the signature

$$\bullet \rhd \mathsf{U} \rhd \mathsf{El}\,\mathsf{q} \rhd \big(\mathsf{q}[\mathsf{p}] \Rightarrow \mathsf{El}\,(\mathsf{q}[\mathsf{p}])\big) : \mathsf{Con},$$

in the graph model this is

$$(N : \mathsf{U}^\circ) \times (\mathsf{El}^\circ\, N) \times (N \Rightarrow^\circ \mathsf{El}^\circ\, N)$$

and for any two (N, z, s), (N', z', s') a set

$$(\overline{N} : N \Rightarrow^\circ \mathsf{El}^\circ\, N') \times (\overline{N}\, z = z') \times (\Pi^\circ(n : N).\overline{N}\,(s\, n) = s'\,(\overline{N}\, n)),$$

and externally we obtain notions of $\mathbb{N}$-algebra

$$(N : \mathsf{Ty}_{\mathcal{C}}\,\bullet) \times \mathsf{Tm}_{\mathcal{C}}\ \bullet\ N \times \mathsf{Tm}_{\mathcal{C}}\,(\bullet \rhd N)\,(N[\mathsf{p}])$$

and homomorphism for any two algebras (N, z, s), (N', z', s'):

$$(\overline{N} : \mathsf{Tm}_{\mathcal{C}}\,(\bullet \rhd N)\, N') \times (\overline{N}[\epsilon, z] = z') \times (\overline{N}[\mathsf{p}, s] = s'[p, \overline{N}])$$

Closed IIT signatures: semantics (iv)

We use Agda syntax to work in $\hat{\mathcal{C}}$.

$$
\begin{aligned}
&\mathsf{U}^\circ \ : \mathsf{Set} && (\mathsf{Ty}_\mathcal{C})\\
&\mathsf{El}^\circ : \mathsf{U}^\circ \to \mathsf{Set} && (\mathsf{Tm}_\mathcal{C})\\
&\Pi^\circ \ : (a^\circ : \mathsf{U}^\circ) \to (\mathsf{El}^\circ\, a^\circ \to \mathsf{Set}) \to \mathsf{Set} && (\rhd_\mathcal{C})
\end{aligned}
$$

We can extend the graph model to the AMDS model

$$
\begin{aligned}
\mathsf{Con} := \ &(\Gamma^\mathsf{A} : \mathsf{Set})\times\\
&(\Gamma^\mathsf{M} : \Gamma^\mathsf{A} \to \Gamma^\mathsf{A} \to \mathsf{Set})\times\\
&(\Gamma^\mathsf{D} : \Gamma^\mathsf{A} \to \mathsf{Set})\times\\
&(\Gamma^\mathsf{S} : (\gamma : \Gamma^\mathsf{A}) \to \Gamma^\mathsf{D}\,\gamma \to \mathsf{Set})
\end{aligned}
$$

This is an inverse diagram model, see Shulman 2012, Lumsdaine 2018 HoTTEST talk, Lumsdaine–Kapulkin 2020.

Example

For natural numbers, the AMDS model gives notions of
Algebras:

$$(N : \mathrm{Set}) \times N \times (N \to N),$$

Morphisms between algebras (N, z, s), (N', z', s'):

$$(\overline{N} : N \to N') \times (\overline{N}\, z = z') \times (\overline{N}\,(s\, n) = s'\,(\overline{N}\, n)),$$

Displayed algebras over an algebra (N, z, s):

$$(\dot{N} : N \to \mathrm{Set}) \times (\dot{N}\, z) \times (\dot{N}\, n \to \dot{N}\,(s\, n)),$$

Sections of displayed algebras $(\dot{N}, \dot{z}, \dot{s})$:

$$(\overline{N} : (n : N) \to \dot{N}\, n) \times (\overline{N}\, z = z') \times (\overline{N}\,(s\, n) = s'\,(\overline{N}\, n)).$$

A CwF $\mathcal{C}$ supports a closed IIT

Externally, for a QIIT signature Ω, from the AMDS model we get:

$$\begin{aligned}
\Omega^{\mathsf{A}} &: \mathsf{Ty}_{\hat{\mathcal{C}}}\,\bullet \\
\Omega^{\mathsf{M}} &: \mathsf{Ty}_{\hat{\mathcal{C}}}\,(\bullet \rhd \Omega^{\mathsf{A}} \rhd \Omega^{\mathsf{A}}[\mathsf{p}]) \\
\Omega^{\mathsf{D}} &: \mathsf{Ty}_{\hat{\mathcal{C}}}\,(\bullet \rhd \Omega^{\mathsf{A}}) \\
\Omega^{\mathsf{S}} &: \mathsf{Ty}_{\hat{\mathcal{C}}}\,(\bullet \rhd \Omega^{\mathsf{A}} \rhd \Omega^{\mathsf{D}})
\end{aligned}$$

The CwF $\mathcal{C}$ supports a QIIT with signature Ω, if there is a

$$\mathsf{con} : \mathsf{Tm}_{\hat{\mathcal{C}}}\ \bullet\ \Omega^{\mathsf{A}}$$

and an

$$\mathsf{elim} : \mathsf{Tm}_{\hat{\mathcal{C}}}\,(\bullet \rhd \Omega^{\mathsf{D}}[\epsilon, \mathsf{con}])\,(\Omega^{\mathsf{S}}[\epsilon, \mathsf{con}[\mathsf{p}], \mathsf{q}]).$$

(This specifies definitional computation rules.)

Summary up to now

We showed what it means that a CwF $\mathcal{C}$ has closed IITs.

- A signature is a context in ToS.
- The AMDS model of ToS internal to $\hat{\mathcal{C}}$ uses U°, El°, Π°.
- Externally we get notions of constructors, eliminator.

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External parameters

New type former in ToS (internal to $\hat{\mathcal{C}}$):

$$\begin{aligned} \hat{\Pi} \quad &: (a^\circ : \mathsf{U}^\circ) \to (a^\circ \Rightarrow^\circ \mathsf{Ty}\,\mathit{\Gamma}) \to \mathsf{Ty}\,\mathit{\Gamma} \\ -\hat{@}- &: \mathsf{Tm}\,\mathit{\Gamma}\,(\hat{\Pi}\,a^\circ\,B) \to \Pi^\circ(x : a^\circ).\mathsf{Tm}\,\mathit{\Gamma}\,(B\,x) \end{aligned}$$

In the standard model,

$$\hat{\Pi}\,a^\circ\,B\,\gamma := \Pi^\circ(x : a^\circ).(B\,x\,\gamma)$$

If $\mathcal{C}$ has $\mathbb{N}$, then we have $\mathbb{N}^\circ : \mathsf{U}^\circ$ and we can specify vectors:

- $\begin{aligned} \triangleright V \quad &: \mathbb{N}^\circ \hat{\Rightarrow} \mathsf{U} \triangleright \\ \mathit{nil} \quad &: \mathsf{El}\,(V\,\hat{@}\,0) \triangleright \\ \mathit{cons} &: a^\circ \hat{\Rightarrow} \hat{\Pi}(n : \mathbb{N}^\circ).V\,\hat{@}\,n \Rightarrow \mathsf{El}\,(V\,\hat{@}\,(1+n)) \end{aligned}$

and the Chapman-style syntax of type theory with an infinite hierarchy of universes.

Equations (identity type with reflection)

New type former in ToS:

$$\begin{aligned} \mathsf{Eq} &: (a : \mathsf{Tm}\,\Gamma\,\mathsf{U}) \to \mathsf{Tm}\,\Gamma\,(\mathsf{El}\,a) \to \mathsf{Tm}\,\Gamma\,(\mathsf{El}\,a) \to \mathsf{Ty}\,\Gamma \\ \mathsf{reflect} &: \mathsf{Tm}\,\Gamma\,(\mathsf{Eq}\,a\,u\,v) \to u = v \end{aligned}$$

In the standard model:

$$\mathsf{Eq}_a\,u\,v\,\gamma := (u\,\gamma =_{\mathsf{El}^\circ\,a\,\gamma} v\,\gamma)$$

Now we can specify all strict QIITs (where the equations are definitional equalities). E.g. integers:

- $\triangleright Z : \mathsf{U} \triangleright \mathit{zero} : \mathsf{El}\,Z \triangleright \mathit{suc} : Z \Rightarrow \mathsf{El}\,Z \triangleright \mathit{pred} : Z \Rightarrow \mathsf{El}\,Z \triangleright$

$\beta : \Pi(i : Z).\mathsf{Eq}\,Z\,(\mathit{pred} \mathbin{@} (\mathit{suc} \mathbin{@} i))\,i \triangleright$

$\eta : \Pi(i : Z).\mathsf{Eq}\,Z\,(\mathit{suc} \mathbin{@} (\mathit{pred} \mathbin{@} i))\,i$

or type theory as a QIIT.

Equations (U is closed under identity with J)

New type former in ToS:

$$\mathsf{Id} : (a : \mathsf{Tm}\,\Gamma\,\mathsf{U}) \to \mathsf{Tm}\,\Gamma\,(\mathsf{El}\,a) \to \mathsf{Tm}\,\Gamma\,(\mathsf{El}\,a) \to \mathsf{Tm}\,\Gamma\,\mathsf{U}$$

with the usual J elimination rule.
If $\mathcal{C}$ has identity types with J, in $\hat{\mathcal{C}}$ we have
$\mathsf{id}^\circ : (a^\circ : \mathsf{U}^\circ) \to \mathsf{El}^\circ\,a^\circ \to \mathsf{El}^\circ\,a^\circ \to \mathsf{U}^\circ$. In the standard model:

$$\mathsf{Id}_a\,u\,v\,\gamma := \mathsf{El}^\circ\,(\mathsf{id}^\circ_{a\,\gamma}\,(u\,\gamma)\,(v\,\gamma))$$

Now we can specify all HIITs (Kaposi-Kovács 2020). E.g. the torus:

- $\bullet \rhd T : \mathsf{U} \rhd b : \mathsf{El}\,T \rhd p : \mathsf{El}\,(\mathsf{Id}_T\,b\,b) \rhd q : \mathsf{El}\,(\mathsf{Id}_T\,b\,b) \rhd t : \mathsf{Id}_{\mathsf{Id}_T\,b\,b}\,(p \cdot q)\,(q \cdot p)$

where $\cdot$ is defined using J.

Infinitary operators

New type former in ToS (internal to $\hat{\mathcal{C}}$):

$$\tilde{\Pi} \quad : (a^\circ : \mathsf{U}^\circ) \to (a^\circ \Rightarrow^\circ \mathsf{Tm}\,\Gamma\,\mathsf{U}) \to \mathsf{Tm}\,\Gamma\,\mathsf{U}$$
$$-\tilde{@}- : \mathsf{Tm}\,\Gamma\,(\tilde{\Pi}\,a^\circ\,b) \to \Pi^\circ(x : a^\circ).\mathsf{Tm}\,\Gamma\,(\mathsf{El}\,(b\,x))$$

If $\mathcal{C}$ has function space, in $\hat{\mathcal{C}}$ we have
$\pi^\circ : (a^\circ : \mathsf{U}^\circ) \to (a^\circ \Rightarrow^\circ \mathsf{U}^\circ) \to \mathsf{U}^\circ$.
In the standard model,

$$\tilde{\Pi}\,a^\circ\,b\,\gamma := \pi^\circ(x : a^\circ).(b\,x\,\gamma)$$

If $\mathcal{C}$ has $\mathbb{N}$, then we have $\mathbb{N}^\circ : \mathsf{U}^\circ$ and we can specify infinitely branching trees:

$$\bullet \rhd T : \mathsf{U} \rhd \mathit{leaf} : \mathsf{El}\;T \rhd \mathit{node} : (\mathbb{N}^\circ \tilde{\Rightarrow} T) \Rightarrow \mathsf{El}\;T$$

Now we can specify ToS itself, real numbers, the partiality monad.

Summary of operators

- U, El,
- Π with domain in U,
- $\hat{\Pi}$ with domain in U°,
- Eq: extensional identity,
- Id: intensional identity,
- $\tilde{\Pi}$ in U, with domain in U°.

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flCwF model (i)

If $\mathcal{C}$ is a model of ETT, the AMDS model can be extended to a finite limit CwF model: CwF + Σ + Eq + K (Nordvall Forsberg PhD 2013, c.f. democracy, Dybjer–Clairambault 2014):

$$\mathsf{K} : \mathsf{Con} \to \mathsf{Ty}\,\Gamma \qquad\qquad \mathsf{mkK} : \mathsf{Sub}\,\Gamma\,\Delta \cong \mathsf{Tm}\,\Gamma\,(\mathsf{K}\,\Delta) : \mathsf{unK}$$

The model (AMDS is the Con,Sub,Ty,Tm components):

- Contexts are flCwFs
- Substitutions strict flCwF morphisms
- Types are displayed flCwFs (c.f. Ahrens–Lumsdaine 2019)
- Terms are strict flCwF sections

this supports U, El, Π, $\hat{\Pi}$, Eq, but not $\tilde{\Pi}$, Id.
See Altenkirch–Kaposi-Kovács POPL 2019.

flCwF model (ii)

If $\mathcal{C}$ is a model of ETT, the AMDS model can be extended to a finite limit CwF model: CwF + Σ + Eq + K (Nordvall Forsberg PhD 2013, c.f. democracy, Dybjer–Clairambault 2014):

$$\mathsf{K} : \mathsf{Con} \to \mathsf{Ty}\,\Gamma \qquad\qquad \mathsf{mkK} : \mathsf{Sub}\,\Gamma\,\Delta \cong \mathsf{Tm}\,\Gamma\,(\mathsf{K}\,\Delta) : \mathsf{unK}$$

The model (AMDS is the Con,Sub,Ty,Tm components):

- Contexts are flCwFs
- Substitutions weak flCwF morphisms (pseudomorphisms)
- Types are split flCwF isofibrations
- Terms are weak flCwF sections

this supports U, El, Π, $\hat{\Pi}$, Eq, $\tilde{\Pi}$, Id.
See Kovács–Kaposi LICS 2020.

Initiality ↔ induction

For each signature, we obtain a CwF + Σ + Eq + K. We prove that initiality is equivalent to induction in the internal language. Assume a Θ : Con.

$$\begin{aligned}
&\mathsf{rec} &&: (\Gamma : \mathsf{Con}) \to \mathsf{Sub}\,\Theta\,\Gamma \\
&\mathsf{uni} &&: (\sigma\,\delta : \mathsf{Sub}\,\Theta\,\Gamma) \to \sigma = \delta \\
&\mathsf{elim} &&: (A : \mathsf{Ty}\,\Theta) \to \mathsf{Tm}\,\Theta\,A \\
&\mathsf{elim}\,A &&:= \mathsf{q}[\mathsf{rec}\,(\Theta \rhd A)] : \mathsf{Tm}\,\Theta\,(A[\underbrace{\mathsf{p} \circ \mathsf{rec}\,(\Theta \rhd A)}_{=\mathsf{id}\text{ by }\mathsf{uni}\,\mathsf{id}\,_}]) \\
&\mathsf{rec}\,\Gamma &&:= \mathsf{unK}\,(\mathsf{elim}\,(\mathsf{K}\,\Gamma)) \\
&\mathsf{uni}\,\sigma\,\delta &&:= \mathsf{ap}\,\mathsf{unK}\,\underbrace{\Big(\mathsf{reflect}\,\big(\mathsf{elim}\,(\mathsf{Eq}\,(\mathsf{mkK}\,\sigma)\,(\mathsf{mkK}\,\delta))\big)\Big)}_{:\,\mathsf{mkK}\,\sigma = \mathsf{mkK}\,\delta}
\end{aligned}$$

Initial algebras

If a model of ETT supports the ToS, then it supports all (Q)IITs specified by the ToS (for all combinations of ToS type formers).

Idea: natural numbers can be defined:

$$\begin{aligned}
\mathbb{N} &:= \mathsf{Tm}_{\mathsf{ToS}}\,(\bullet \triangleright N : \mathsf{U} \triangleright z : \mathsf{El}\,N \triangleright s : N \Rightarrow \mathsf{El}\,N)\,(\mathsf{El}\,N) \\
\mathsf{zero} &:= z \\
\mathsf{suc}\,t &:= s\,@\,t
\end{aligned}$$

If we interpret the term in the standard model A, we get Church encoding (implementing the recursor):

$$\mathsf{Tm}_{\mathsf{A}}\,(\bullet \triangleright N : \mathsf{U} \triangleright z : \mathsf{El}\,N \triangleright s : N \Rightarrow \mathsf{El}\,N)\,(\mathsf{El}\,N) = \big((N : \mathsf{Set}) \times N \times (N \to N)\big) \to N$$

If interpret in the graph model AM, we get the Awodey-Frey-Speight encoding (LICS 2018).

Results on existence of initial algebras

If a model of ETT supports the ToS, then it supports all (Q)IITs specified by the ToS (for all combinations of ToS type formers).

- In ETT with indexed W types, we can define the ToS with $\mathsf{U}, \mathsf{El}, \Pi, \hat{\Pi}$ (Kaposi–Lafont–Kovács, TYPES 2019 post-proc)
- WIP: show that the setoid model supports ToS with $\mathsf{U}, \mathsf{El}, \Pi, \hat{\Pi}, \mathsf{Id}, \tilde{\Pi}$ (Kaposi–Zongpu TYPES 2020)
- stealing from Brunerie–Menno de Boer's (HoTTEST talk) formalisation: they have U, El, Π, Id: in ETT + quotients + propext, we can derive all closed QIITs

Negative result: certain infinitary QIITs cannot be defined in ETT + quotients (Lumsdaine–Shulman 2019).

A direct reduction (see Altenkirch–Kaposi–Kovács–Von Raumer, TYPES 2019) might work in intensional models and would give definitional computation rules.

Contents

- Formal specification of closed IITs
- Extension to QIITs
- Initial algebras
- HIITs
- Higher order abstract syntax (syntax with binding)

Categorical semantics of HIITs

Capriotti and Sattler (see abstract at TYPES 2020):

- construct a higher category of algebras from a signature
- support $\mathsf{U}, \mathsf{El}, \Pi, \hat{\Pi}, \tilde{\Pi}, \mathsf{Id}$
- define displayed algebras and sections
- show the equivalence of initiality and induction
- work in $\hat{\mathcal{C}}$ for a model of HoTT $\mathcal{C}$

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Signatures for type theories (WIP) (i)

We know how to say that a CwF $\mathcal{C}$ supports a QIIT.

How do we say that a CwF supports Π types, Σ types, coinductive types etc.? We could define CwF with Π and Σ as a QIIT, but that has two problems:

- overhead: then our semantics says what it means that another CwF supports an (internal) CwF
- we would need to write substitution rules such as $\Pi\, A\, B[\sigma] = \Pi\,(A[\sigma])\,(B[\sigma \circ \mathsf{p}, \mathsf{q}])$ by hand.

A possible solution, based on Capriotti's Rule Framework (TYPES 2017):

- the QIIT-ToS has Ty which we call Ty^0 from now on,
- new sort for Ty^1 types with, $\uparrow$: $\mathsf{Ty}^0\,\Gamma \to \mathsf{Ty}^1\,\Gamma$
- Ty^1 has a function space with domain in Ty^0 and Eq of Ty^0
- a signature is a context in this general ToS

Signatures for type theories (WIP) (ii)

Signature for Π with β:

- $\triangleright \mathit{pi} \quad : \Pi^1(a : \mathsf{U}).(a \Rightarrow \mathsf{U}) \Rightarrow^1 \uparrow \mathsf{U} \triangleright$

  $\mathit{lam} : \Pi^1(a : \mathsf{U}).\Pi^1(b : a \Rightarrow \mathsf{U}).$

  $\qquad ((x : a) \Rightarrow \mathsf{El}\,(b \mathbin{@} x)) \Rightarrow^1 \uparrow (\mathsf{El}\,(\mathit{pi} \mathbin{@^1} a \mathbin{@^1} b)) \triangleright$

  $\mathit{app} : \Pi^1(a : \mathsf{U}).\Pi^1(b : a \Rightarrow \mathsf{U}).$

  $\qquad \mathsf{El}\,(\mathit{pi} \mathbin{@^1} a \mathbin{@^1} b) \Rightarrow^1 \uparrow ((x : a) \Rightarrow \mathsf{El}\,(b \mathbin{@} x)) \triangleright$

  $\beta \quad : \Pi^1(a : \mathsf{U}).\Pi^1(b : a \Rightarrow \mathsf{U}).\Pi^1(t : (x : a) \Rightarrow \mathsf{El}\,(b \mathbin{@} x)).$

  $\qquad \mathsf{Eq}_{(x:a) \Rightarrow \mathsf{El}\,(b @ x)} \big(\mathit{app} \mathbin{@^1} a \mathbin{@^1} b \mathbin{@^1} (\mathit{lam} \mathbin{@^1} a \mathbin{@^1} b \mathbin{@^1} t)\big)\, t$

Signatures for type theories (WIP) (iii)

Conversions:

- TT signature $\to$ QIIT signature:
  - adds substitution laws
  - obtain category of models, initiality
- QIIT signature $\to$ TT signature:
  - adds elimination principles
  - obtain syntactic description

We can generalise type theory signatures to arbitrary signatures with binding. In a CwF $\mathcal{C}$, $\mathsf{Ty}_{\mathcal{C}} : \mathsf{Ty}_{\hat{\mathcal{C}}}\,\bullet$, but $\mathsf{Tm}_{\mathcal{C}} : \overline{\mathsf{Ty}}_{\hat{\mathcal{C}}}\,(\bullet \rhd \mathsf{Ty}_{\mathcal{C}})$.

$$\overline{\mathsf{Ty}}_{\hat{\mathcal{C}}}\,\Gamma = (A : \mathsf{Ty}_{\hat{\mathcal{C}}}\,\Gamma) \times (- \rhd_A - : (I : |\mathcal{C}|) \to |\Gamma|_I \to |\mathcal{C}|) \times \mathcal{C}(J, I \rhd_A \gamma) \cong (f : \mathcal{C}(J, I)) \times |A|_J\,(\gamma f)$$

See also: Bocquet–Kaposi–Sattler TYPES 2020, Awodey's natural models 2014, Uemura 2019, HoTTEST talks: Sterling, Bauer, Altenkirch.

Summary

- A QIIT/HIIT can be described as a context in a well chosen type theory of signatures.
- Models of the type theory of signatures provide semantics for QIITs/HIITs.
- In ETT, if we have the ToS, we get all QIITs.
- We can extend the theory of QIIT signatures to the theory of type theory signatures.