



Privacy and Economics

The Swiss Blockchain Winter School, Interlaken, 2019

Rainer Böhme



This is **not** Interlaken. (Picture taken last week.)

Privacy in Blockchain-based Systems

Fundamental conflict between:



- public blockchain data, required for distributed verification,
- and
- users' right of personality, because the information contained in (financial) transactions reveals personal preferences and circumstances of life.

Hope of mitigation:

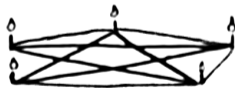
- Pseudonyms are not directly linkable to natural persons.

Nevertheless, EU lawyers consider blockchain data as **personal data** because:

1. the link to natural persons can often be established with **additional information** (e. g. known by intermediaries, such as exchanges or network relays);
2. the persistence of blockchain data increases the probability that this will happen some time in the future.

WHATCHYA DOING?

ECONOMICS.



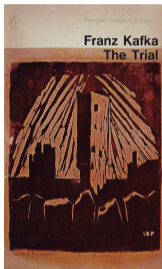
Outline

- 1. How to Plug Privacy into Economic Equations**
2. Observing a Market for Anonymity
3. How to Price Anonymity

Acknowledgement: Part of the material in this talk is joint work with Malte Möser and Daniel G. Arce.

Challenge

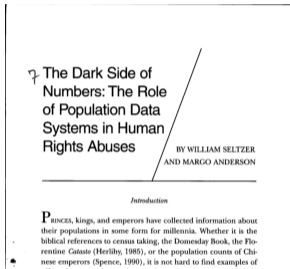
Worst case: the system operator is the attacker.



fiction



reality



Seltzer, W. & Anderson, M. in *Social Research* 68 (2), 2001

Bad, but almost unavoidable case: the system operator makes mistakes.

→ **Dead end: Quantifying the disutility of personal data abuse**

Schools of Thought

Classical economic theory

Efficient markets, perfect information, ...

Why does advertising exist?

New institutional economics

Information matters, asymmetry causes misallocation, ...

Adverse selection: privacy protects bad risks

Asset pricing

Present value of expected future benefit of (re)identification

Berthold & Böhme 2009

Behavioral economics

Bounded rationality, human-subject experiments

Find price tags for personal data

→ **No general theory of information distribution between economic agents**

Economics Value of Personal Data

Price discrimination by HTTP User-Agent string:

Room Rates
Orbitz is starting to show different results to users of Macs and PCs after finding Mac users spend more freely. In a recent search for hotels in Miami, Mac users saw more options over \$200. For El Paso, Texas, they saw more expensive options in the place of two cheaper ones.

Mac	Miami	PC	Mac	El Paso	PC
1. Hyatt House \$118	1. Hyatt House \$118		1-5. Same for both	1-5. Same for both	
2. Design Suites \$124	2. Catalina Hotel \$209		6. Wyndham El Paso \$76	6. Travelodge \$40	
3. Catalina Hotel \$209	3. Design Suites \$124		7. Studio Plus Deluxe \$54	7. Wyndham El Paso \$76	
4. Churchill Suites \$189	4. The Richmond Hotel \$156		8. Hyatt Place El Paso \$76	8. Studio Plus Deluxe \$54	
5. The Richmond Hotel \$156	5. Churchill Suites \$189		9. El Paso Marriott \$89	9. Days Inn \$55	
6. Eden Roc Renaissance \$212	6. Ocean Spray \$95		10. Radisson Hotel \$98	10. Hyatt Place El Paso \$76	
7. The Palms Hotel & Spa \$224	7. South Seas Hotel \$175				

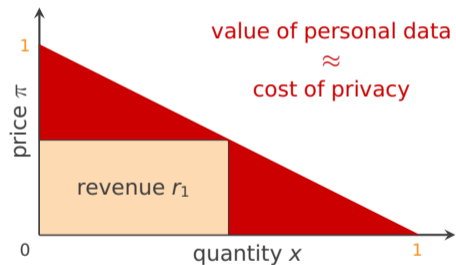
Source: WSJ searches of Orbitz that were performed at the same time for the same dates using a Mac with a Safari browser and a PC with Internet Explorer The Wall Street Journal

Source: The Wall Street Journal 2012

Privacy and Price Discrimination

Demand function $D : \pi \mapsto x$

Case 2: Seller knows each buyers' willingness to pay



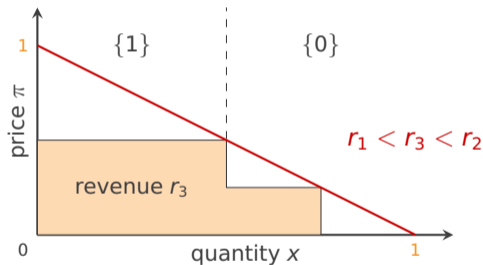
Odlyzko 2003

Assumptions: monopolistic seller, no arbitrage, zero marginal cost

Privacy and Price Discrimination

Demand function $D : \pi \mapsto x$

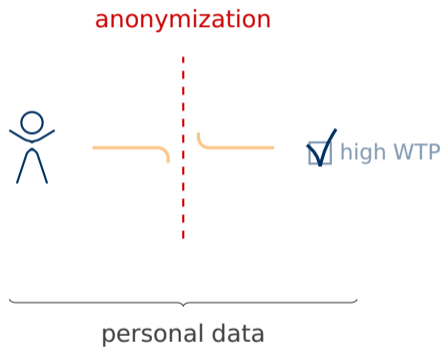
Case 3: Seller knows **one bit** about each buyers' WTP



Odlyzko 2003

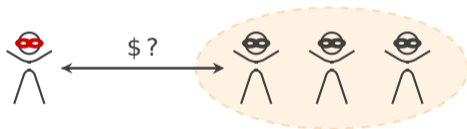
Assumptions: monopolistic seller, no arbitrage, zero marginal cost

Technical Approach



Economics of Anonymity

*“Anonymity is the state of not being identifiable within a set of subjects, the **anonymity set**.”*



“anonymity loves company”

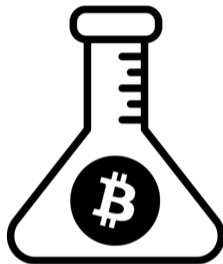
Dingledine & Mathewson 2006

Pfitzmann & Köhntopp 2001

Opportunity



The price of anonymity



Bitcoin as a
social science lab

Icons based on
Flask by Andrew Was from the Noun Project

Tag and Eye Mask by Creative Stall from the Noun Project

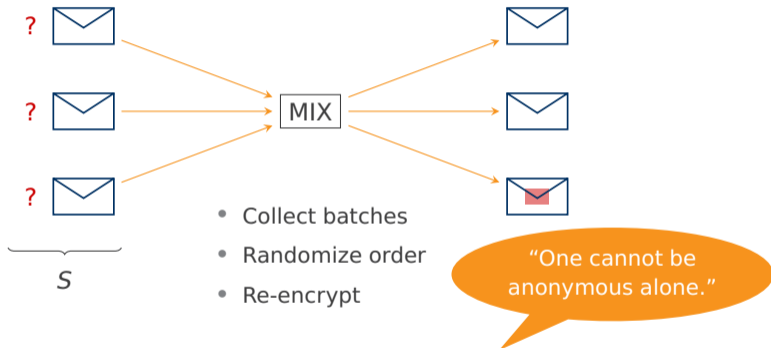
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The Mixing Principle

Establish **unlinkability** of messages in communication systems.

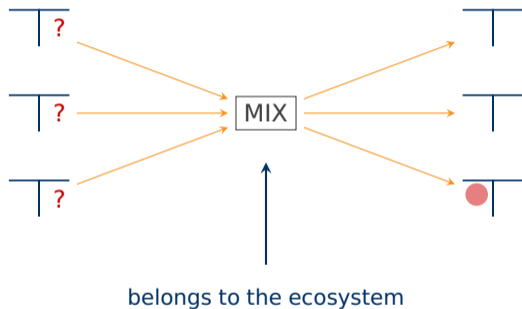


The size of the **anonymity set** $|S|$ is a measure of privacy.

Chaum, D. Untraceable Electronic Mail, Return Addresses, and Digital Pseudonyms. *CACM*, 24 (2), 1981, pp. 84–88.

Application of the Mixing Principle to Bitcoin

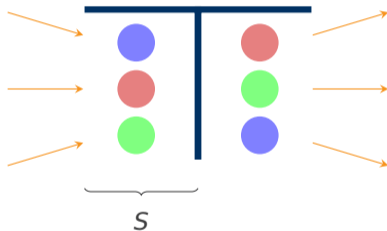
Establish unlinkability of flows in **transaction systems**.



Requires substantial **trust** in mix operator.

“CoinJoin” Transactions

Bitcoin’s **transaction logic** allows multiple inputs and outputs.



More secure alternative: all participants must sign the transaction.

Maxwell 2013

Matchmaking for CoinJoins

Join Market **Orders** Size Distribution Depth Export orders GitHub Getting Started

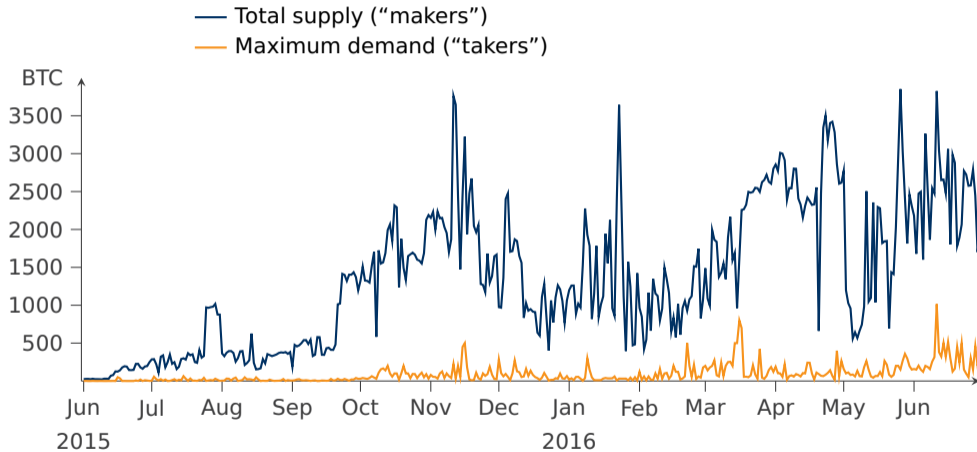
JoinMarket Orderbook

364 orders found by 91 counterparties

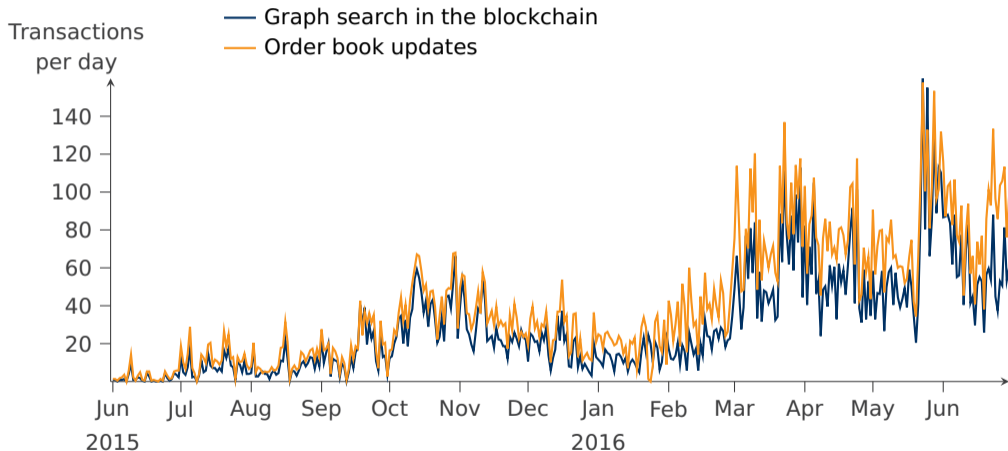
Type	Counterparty	Order ID	Fee	Miner Fee Contribution / BTC	Minimum Size / BTC	Maximum Size / BTC
Absolute Fee	J5BNWo4MhLbtAej1	0	0.00000400	0.00000002	0.00002730	0.00863625
Absolute Fee	J58HmZ2eFvZqELwx	5	0.00000400	0.00000002	0.00406800	0.00702015
Absolute Fee	J5E9nx6U7k976mCB	2	0.00000400	0.00000002	0.00644100	0.00863625
Absolute Fee	J5CUymfJ9hYrVWAc	0	0.00000500	0.00000000	0.00100000	13.58535521
Absolute Fee	J58HmZ2eFvZqELwx	22	0.00000539	0.00000060	0.00406800	0.00702015
Absolute Fee	J58HmZ2eFvZqELwx	3	0.00000800	0.00000000	0.00406800	0.00702015
Absolute Fee	J5E9nx6U7k976mCB	20	0.00000800	0.00000000	0.00644100	0.24100000
Absolute Fee	J59Z6KFWtWk4wcjM	6	0.00000800	0.00000000	0.00400000	0.24100000
Absolute Fee	J5Bmy7oTZ3lrpdVv	0	0.00000889	0.00000065	0.06886283	1.82194683
Absolute Fee	J59pheQXDj7MZzFp	0	0.00000950	0.00000150	0.00100000	0.71455724
Absolute Fee	J58HmZ2eFvZqELwx	1	0.00000950	0.00000150	0.00406800	0.00702015
Absolute Fee	J5E9nx6U7k976mCB	0	0.00000950	0.00000150	0.00644100	0.71455724
Absolute Fee	J58HmZ2eFvZqELwx	21	0.00000997	0.00000000	0.00406800	0.00702015

<http://joinmarket.io>, last access: February 25th, 2018

Supply and Demand



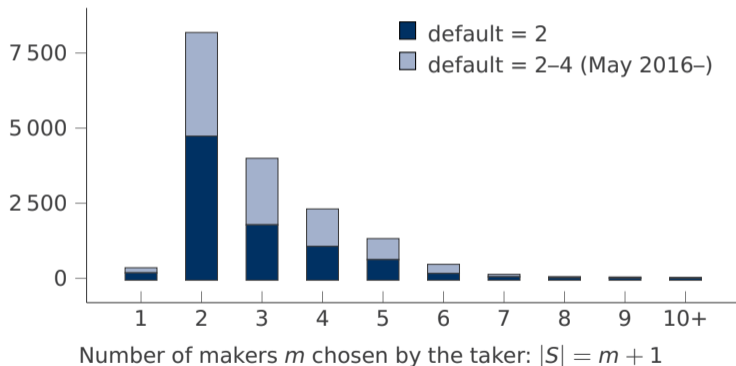
Identifying JoinMarket Transactions



Size of the Anonymity Set

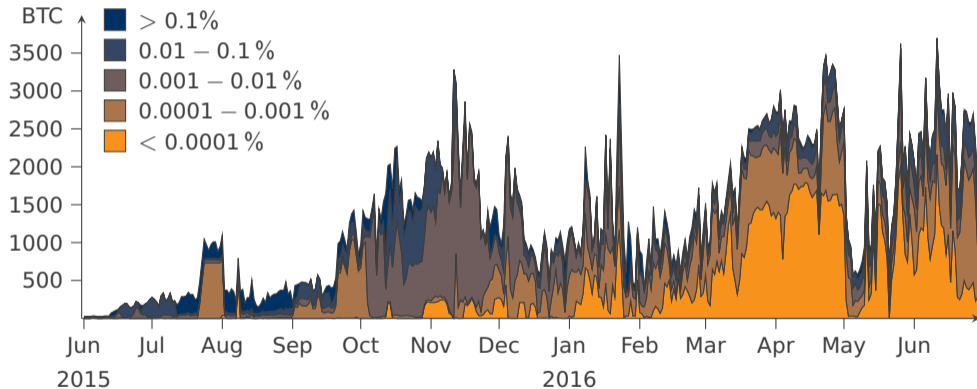
S is composed of exactly one **“taker”** and $m \geq 1$ **“makers”**.

Histogram of 16 K JoinMarket transactions



Empirical Prices of Anonymity

Order book analysis: fee per maker



For comparison: mix operators charge 1-3%.

Möser et al. 2013

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Anonymity Market

“One cannot be anonymous alone.”

Cooperative game theory to model the co-creation of anonymity.

Model

- 1 “taker” and $m \geq 1$ “makers”
- **Only the taker pays** for anonymity: fee f to each maker.
- The taker **and all makers benefit** from anonymity set, $|S| = m + 1$.
- Taker has an outside option, e. g., a mix charging fee $F \gg f$.
- Solve for f endogenously.

→ **Shapley value as solution concept.**

Utility of Anonymity

Assumption: the attacker guesses within the anonymity set (i. e., GPA)

Case 1: coalition with makers

Taker expects: $D \cdot \frac{m}{m+1}$

Maker expects: $d \cdot \frac{m}{m+1}$

$$D \gg d$$

Case 2: outside option

Taker expects: $\delta \cdot D$

Maker receives: 0

$$\delta \in (0, 1)$$

Solution for $m = 2$ Makers (sketch)

Characteristic function V

$$V(\{t\}) = \{x_t \mid x_t \leq \delta D - F\} \quad (1)$$

$$V(\{i\}) = \{x_i \mid x_i \leq 0 : i = 1, 2\} \quad (2)$$

$$V(\{t, 1\}) = \{(x_t, x_1) \mid x_t \leq D/2 - f, x_1 \leq d/2 + f\} \quad (3)$$

$$V(\{t, 2\}) = \{(x_t, x_2) \mid x_t \leq D/2 - f, x_2 \leq d/2 + f\} \quad (4)$$

$$V(\{1, 2\}) = \{(x_1, x_2) \mid x_1 \leq 0, x_2 \leq 0\} \quad (5)$$

$$V(\{t, 1, 2\}) = \{(x_t, x_1, x_2) \mid x_t \leq 2/3D - 2f, x_{1,2} \leq 2/3d + f\} \quad (6)$$

D, d : value of anonymity for taker/maker; f : fee per maker; δ, F : quality/fee of outside option

Solution for $m = 2$ Makers (sketch, cont'd)

Worth function ω

$$\omega(\{t\}) = \delta D - F \quad (7)$$

$$\omega(\{1\}) = \omega(\{2\}) = 0 \quad (8)$$

$$\omega(\{t, 1\}) = D/2 + d/2 \quad (9)$$

$$\omega(\{t, 2\}) = D/2 + d/2 \quad (10)$$

$$\omega(\{1, 2\}) = 0 \quad (11)$$

$$\omega(\{t, 1, 2\}) = 2/3D + 4/3d \quad (12)$$

D, d : value of anonymity for taker/maker; δ, F : quality/fee of outside option

Shapley Value and Associated Fee f

Shapley value φ

$$\varphi_t = \frac{14}{36}D + \frac{22}{36}d + \frac{1}{3}(\delta D - F) \quad (13)$$

$$\varphi_1 = \varphi_2 = \frac{5}{36}D + \frac{13}{36}d - \frac{1}{6}(\delta D - F) \quad (14)$$

“The Price of Anonymity”

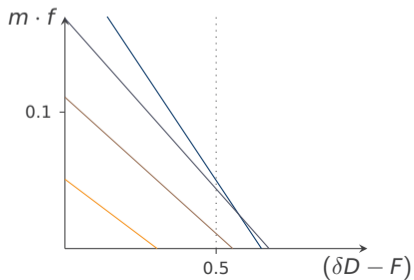
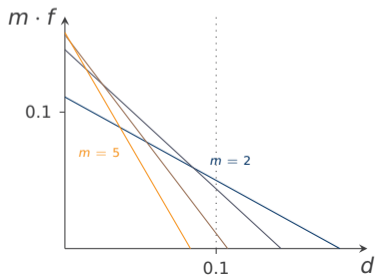
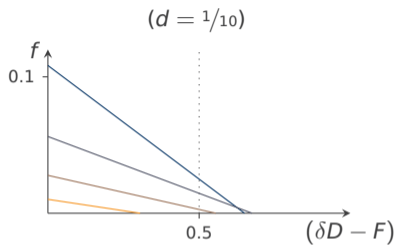
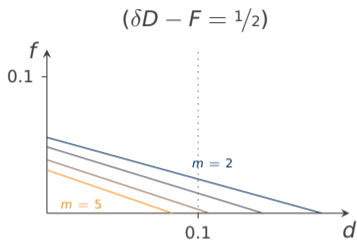
$$f = \frac{5}{36}D - \frac{11}{36}d - \frac{1}{6}(\delta D - F) \quad (15)$$

→ **General solution** for $m > 2$ is efficiently computable.

D, d : value of anonymity for taker/maker; f : fee per maker; δ, F : quality/fee of outside option

Visualization

Normalization: $D = 1$





We have introduced an NTU cooperative game-theoretic model of a CoinJoin anonymity market that is consistent with:

1. The prevalent measure of anonymity: the **anonymity set**.
2. The **peculiar nature** of anonymity markets:

One demand-side participant (the 'taker') pays for anonymity but all suppliers (the 'makers') **also** receive the good in demand (anonymity).

Take-Home Messages

1. Blockchain data is (almost always) personal data.
2. It is possible (but not trivial) to plug privacy into a utility function.
3. Forgotten objective in protocol design: incentivize privacy-enhancing behavior.
4. Blockchain systems are crystal balls for studying the economics of privacy.





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Thank you for your attention.



Talk, research visit, post-doc? rainer.boehme@uibk.ac.at

Further Reading

1. Möser, M. and Böhme, R. The Price of Anonymity: Empirical Evidence from a Market for Bitcoin Anonymization. *Journal of Cybersecurity*, 3, 2 (2017), 127–135.
2. Arce, D. G. and Böhme, R. Pricing Anonymity. In S. Meiklejohn and K. Sako, eds., *Financial Cryptography and Data Security*. 2018.
3. Abramova, S., Schöttle, P., and Böhme, R. Mixing Coins of Different Quality: A Game-Theoretic Approach. In *Financial Cryptography (4th Workshop on Bitcoin and Blockchain Research)*. Malta, 2017.
4. Acquisti, A., Taylor, C., and Wagman, L. The Economics of Privacy. *Journal of Economics Literature*, 54, 6 (2016), 442–292.
5. Böhme, R., Christin, N., Edelman, B., and Moore, T. Bitcoin: Economics, Technology, and Governance. *Journal of Economic Perspectives*, 29, 2 (2015), 213–238.

The Shapley Value Solution for m Makers

After overcoming some technical hurdles, e. g., specifying $\omega : V(\hat{S}) \rightarrow \mathbb{R}$, the **fees** of the Shapley value solution for m makers are:

$$f = \frac{1}{(m+1)}D - \frac{D}{m(m+1)} \sum_{n=1}^m \frac{n}{n+1} - \frac{d}{m(m+1)} \sum_{n=1}^m \frac{n^2}{n+1} + \frac{F - \delta D}{m(m+1)}$$

(A formula replacing the finite sums with harmonic numbers is given in the paper.)

- Increasing in D .
- Increasing in F .
- Decreasing in d .
- Theoretical lower bound for identifying dishonest makers.
- Experimentally/behaviorally testable by endowing subjects with D , d , δ and F values.