Retrieval-Technologien für die Plagiaterkennung in Programmen

Fabian Loose, Steffen Becker, Martin Potthast, Benno Stein @ webis.de Bauhaus University Weimar

Outline

- Overview
- Retrieval Models for Source Code
- · Hash-based Search



Fabian Loose



Steffen Becker



Martin Potthast

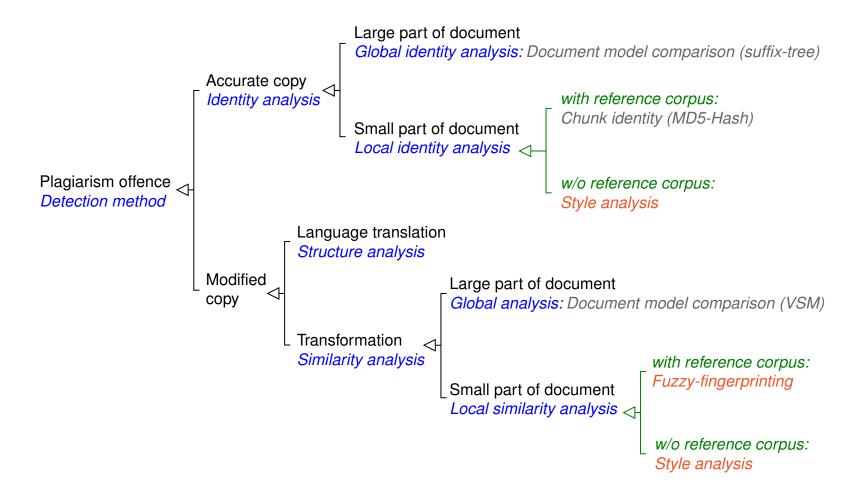
Plagiarism is the practice of claiming, or implying, original authorship of someone else's written or creative work, in whole or in part, into one's own without adequate acknowledgment.

[Wikipedia: Plagiarism]

- □ Plagiarism is observed in literature, music, software, scientific articles, newspaper, advertisement, Web sites, etc.
- □ A study among 18 000 university students in the United States shows that almost 40% of them have plagiarized at least once. [1]

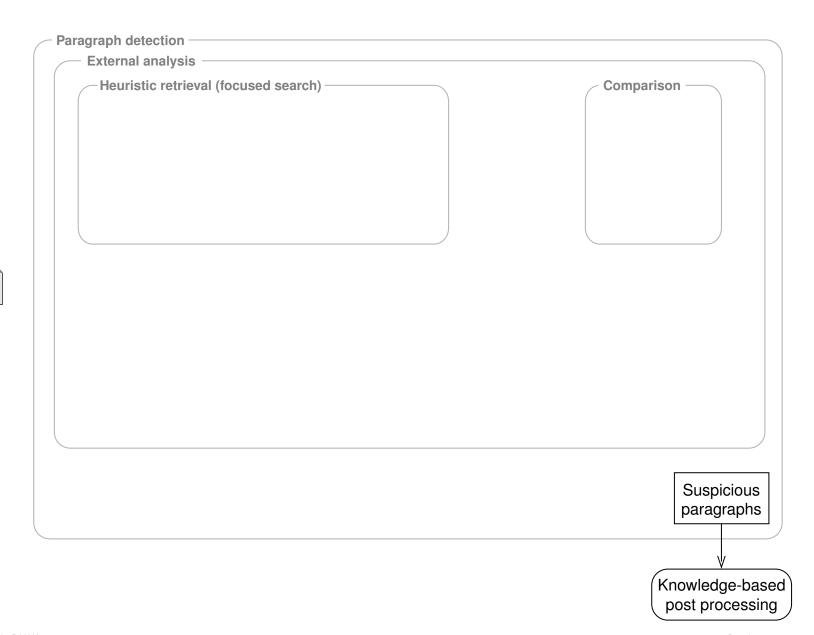
[1] D. McCabe. Research Report of the Center for Academic Integrity. http://www.academicintegrity.org, 2005.

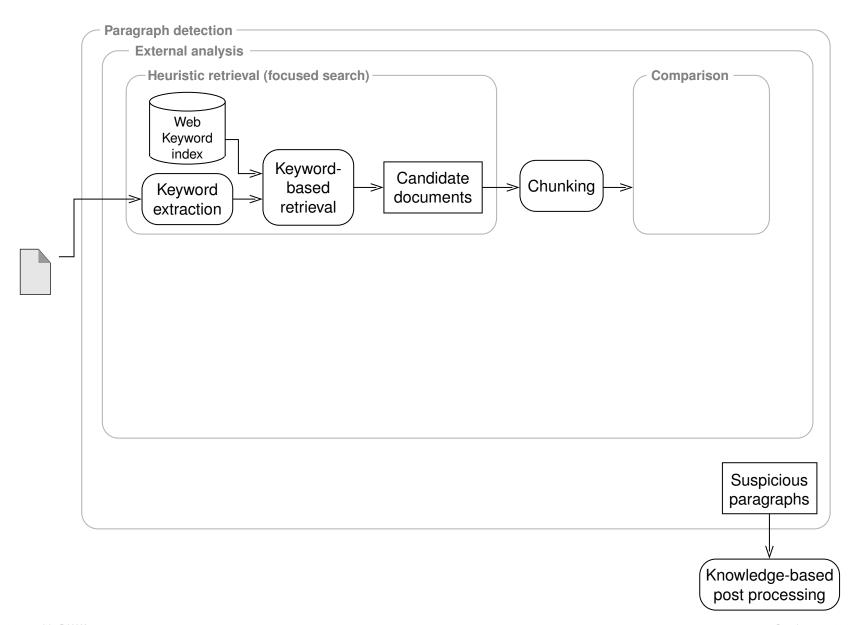
Taxonomy of Plagiarism Offenses

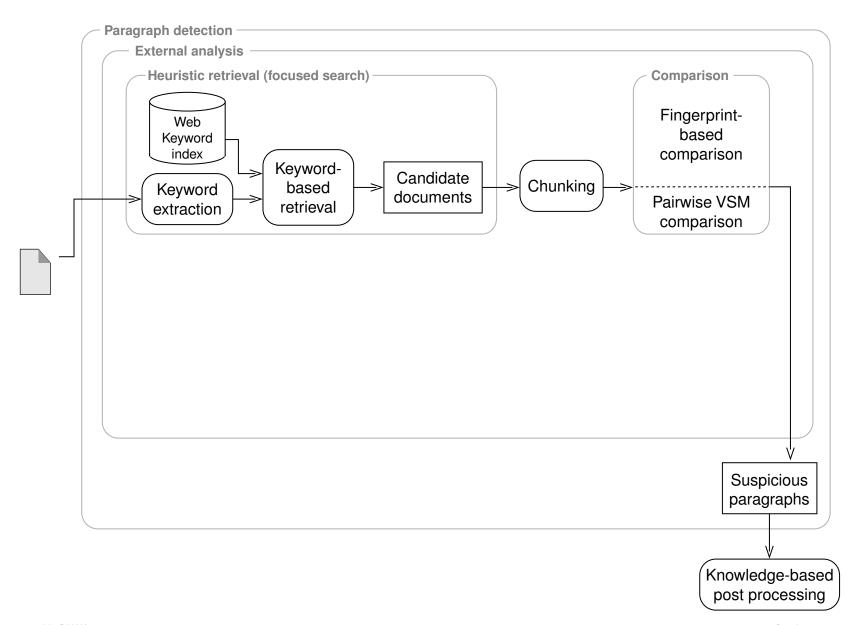


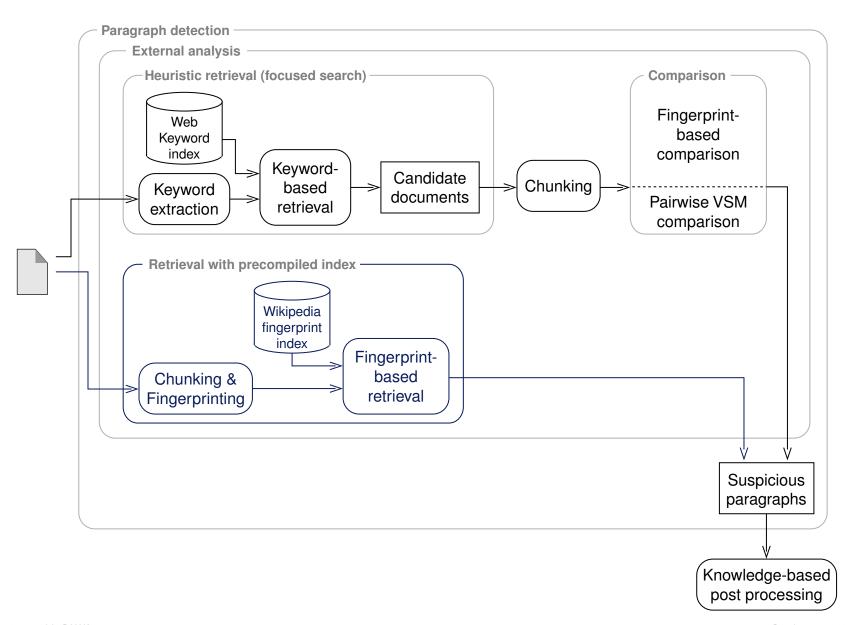
Paragraph detection —		

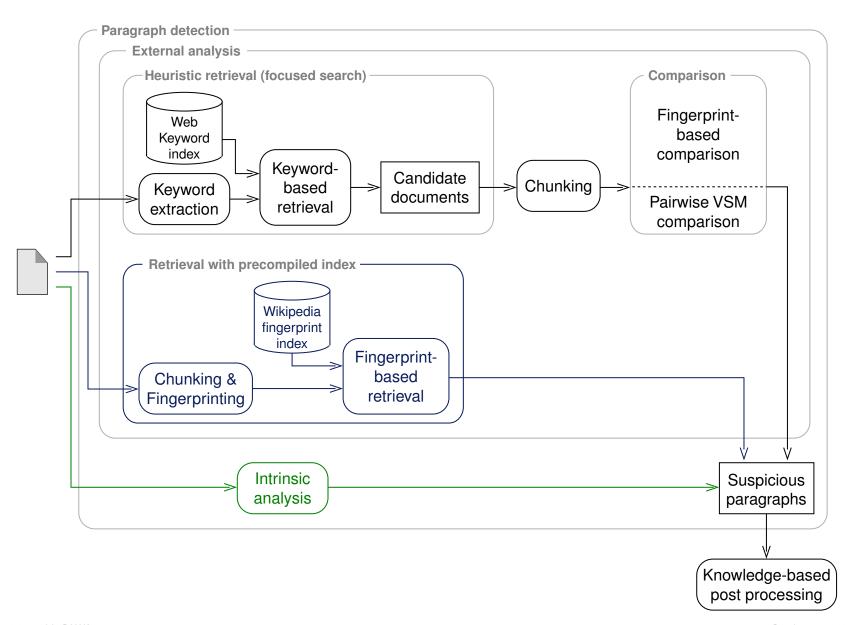
Knowledge-based post processing

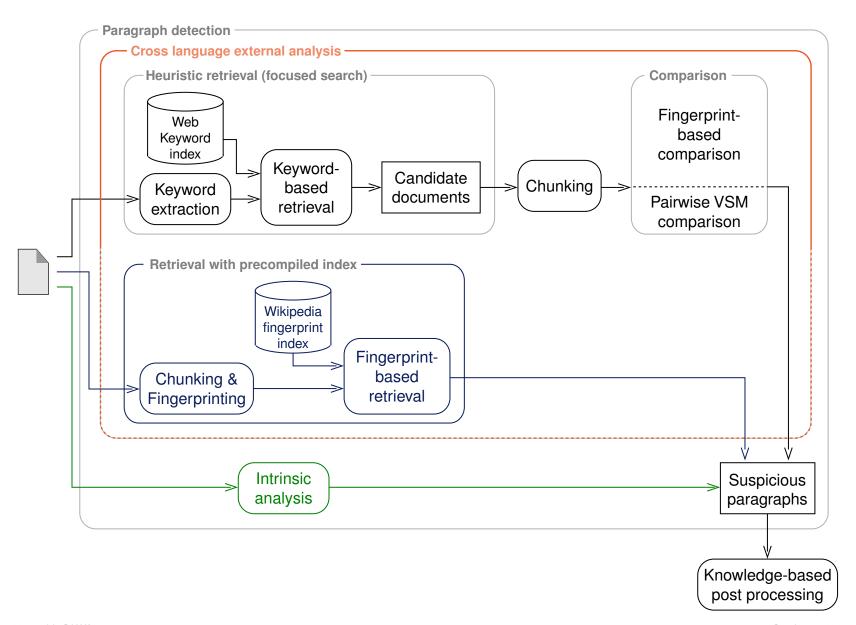












Examples for Identification Technology

□ Level 1. Identity analysis for paragraphs.

MD5 hashing

□ Level 2. Synchronized identity analysis for paragraphs.

hashed breakpoint chunking

□ Level 3. Tolerant similarity analysis for paragraphs.

Fuzzy-fingerprinting

□ Level 4. Intrinsic (style) analysis without a reference corpus.

statistical outlier analysis with Bayes, meta learning with logistic regression

□ Level 5. Correct citation.

knowledge-based analysis

Current research is corpus-centered, "external plagiarism analysis".

[Brin et al. 1995, Monostori et al. 2001-2004, Stein et al. 2004-2006, etc.]

External plagiarism analysis formulated as decision problem:

Problem. AVEXTERN (AV stands for Authorship Verification)

Given. A text d, allegedly written by author A, and set of texts D,

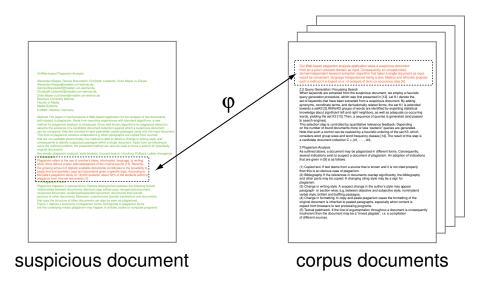
 $D = \{d_1, \ldots, d_n\}$, written by an arbitrary number of authors.

Question. Does d contain sections whose similarity to sections in D is above

a threshold θ ?

Basic Principle

- □ Partition each document in meaningful sections, also called "chunks".
- \Box Do a pairwise comparison using a similarity function φ .



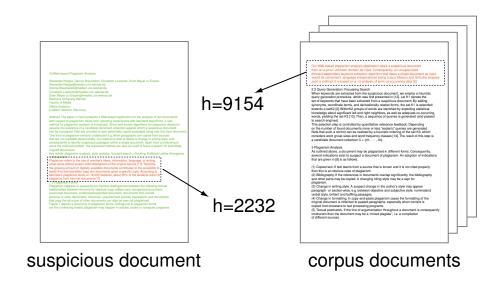
Complexity:

n documents in corpus, c chunks per document on average

 \rightarrow $O(n \cdot c^2)$ comparisons

Comparison with Fingerprints (Level 1)

- Partition each document into equidistant sections.
- \Box Compute fingerprints of the chunks using a hash function h.
- □ Put all hashes into a hash table. A collision indicates matching chunks.



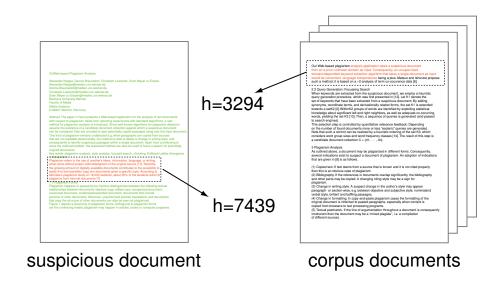
Complexity:

n documents in corpus, c chunks per document on average

 \rightarrow $O(n \cdot c)$ operations (fingerprint generation, hash table operations)

Comparison with Fingerprints (Level 2)

- □ Partition each document into *synchronized* sections.
- \Box Compute fingerprints of the chunks using a hash function h.
- □ Put all hashes into a hash table. A collision indicates matching chunks.



Complexity:

n documents in corpus, c chunks per document on average

 \rightarrow $O(n \cdot c)$ operations (fingerprint generation, hash table operations)

Comparison with Fingerprints (Level 3)

Discussion:

□ Hashing is fast, but sensitive to smallest changes:

$$h(c_1) = h(c_2) \Rightarrow c_1 = c_2$$
 (with very high probability)

Current research:

 \Box Focus on *fuzzy* hash functions h_{φ} :

$$h_{\varphi}(c_1) = h_{\varphi}(c_2) \quad \Rightarrow \quad P(\varphi(c_1, c_2) > \theta) \geq 1 - \varepsilon$$
 [Stein 2005-07]

- □ Fuzzy hash functions allow for large chunk sizes (speed-up)
- Fuzzy hash functions are not sensitive to small changes

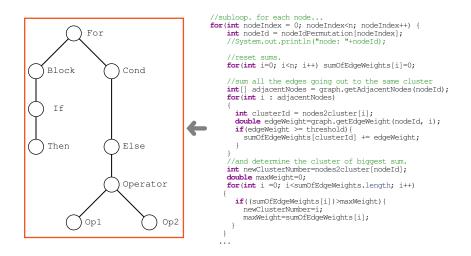
```
//subloop. for each node...
for (int nodeIndex = 0; nodeIndex<n; nodeIndex++) {</pre>
   int nodeId = nodeIdPermutation[nodeIndex];
   //System.out.println("node: "+nodeId);
    //reset sums.
   for(int i=0; i<n; i++) sumOfEdgeWeights[i]=0;</pre>
    //sum all the edges going out to the same cluster
    int[] adjacentNodes = graph.getAdjacentNodes(nodeId);
    for (int i : adjacentNodes)
      int clusterId = nodes2cluster[i];
      double edgeWeight=graph.getEdgeWeight(nodeId, i);
      if (edgeWeight >= threshold) {
        sumOfEdgeWeights[clusterId] += edgeWeight;
    //and determine the cluster of biggest sum.
   int newClusterNumber=nodes2cluster[nodeId];
   double maxWeight=0;
   for(int i =0; i < sumOfEdgeWeights.length; i++)</pre>
     if((sumOfEdgeWeights[i])>maxWeight){
       newClusterNumber=i;
       maxWeight=sumOfEdgeWeights[i];
```

Representation d

Sim. measure φ

 $\begin{array}{ccc} \textbf{Compilation} & \textbf{Runtime} \\ \textbf{level for } d & \textbf{for } \varphi \end{array}$

Structure-based Graph Models



Representation d	Sim. measure φ	Compilation level for \boldsymbol{d}	Runtime for φ	
abstract syntax trees	hash-based subtree search	syntactical	$O(\mathbf{d})$	[Baxter et al. 1998]
conceptual graphs	heuristically focused isomorphic graph search	semantic	$O(\mathbf{d} ^3)$	[Mishne et al. 2004]
program dep. graphs	isomorphic graph search	semantic	NP-complete	[Liu et al. 2006]

Attribute-based Vector Models

```
//subloop. for each node...
                                       for (int nodeIndex = 0; nodeIndex<n; nodeIndex++) {</pre>
         For
                                            int nodeId = nodeIdPermutation[nodeIndex];
                                            //System.out.println("node: "+nodeId);
                                            for(int i=0; i<n; i++) sumOfEdgeWeights[i]=0;</pre>
Block
                  Cond
                                            //sum all the edges going out to the same cluster
                                            int[] adjacentNodes = graph.getAdjacentNodes(nodeId);
                                            for (int i : adjacentNodes)
 Ιf
                                             int clusterId = nodes2cluster[i];
                                             double edgeWeight=graph.getEdgeWeight(nodeId, i);
                                             if (edgeWeight >= threshold) {
                                                sumOfEdgeWeights[clusterId] += edgeWeight;
                  Else
Then
                                            //and determine the cluster of biggest sum.
                                            int newClusterNumber=nodes2cluster[nodeId];
                                            double maxWeight=0;
                  Operator
                                            for (int i =0; i < sumOfEdgeWeights.length; i++)</pre>
                                             if((sumOfEdgeWeights[i])>maxWeight){
                                               newClusterNumber=i;
                                               maxWeight=sumOfEdgeWeights[i];
```

```
for ( int nodeIndex
= 0 : nodeIndex
< n : nodeIndex++
) { int nodeId
= nodeIdPer [ nodeIndex
] ; for (
int i = 0
; i < n
; i ++ )
sumOfEdgeWeights [ i ]
= 0 ; int
[ ] adjacentNodes =
graph . getAdNodes (
nodeId ) ; for
(int i:
adjacentNodes ) { int
clusterId = nodes2clu [
i ] ; double
edgeWeight = graph .
getEdgeWeight ( nodeId ,
i ) ; if
```

Representation d	Sim. measure φ	Compilation level for \boldsymbol{d}	Runtime for φ	
software metric features all n grams subset of all n grams	cosine Jaccard Jaccard	none lexical lexical	$O(\mathbf{d})$ $O(\mathbf{d})$ $O(\mathbf{d})$	[Ottenstein 1976] [Clough et al. 2002] [Schleimer et al. 2003]
n < 5				

Structure-based String Models

```
for ( int nodeIndex
                                   //subloop. for each node...
                                   for (int nodeIndex = 0; nodeIndex<n; nodeIndex++) {</pre>
                                                                                              = 0 : nodeIndex
        For
                                       int nodeId = nodeIdPermutation[nodeIndex];
                                                                                              < n : nodeIndex++
                                       //System.out.println("node: "+nodeId);
                                                                                              ) { int nodeId
                                                                                              = nodeIdPer [ nodeIndex
                                       for(int i=0; i<n; i++) sumOfEdgeWeights[i]=0;</pre>
                                                                                              ] ; for (
Block
                Cond
                                                                                              int i = 0
                                       //sum all the edges going out to the same cluster
                                       int[] adjacentNodes = graph.getAdjacentNodes(nodeId);
                                                                                              ; i < n
                                       for (int i : adjacentNodes)
                                                                                              ; i ++ )
 Ιf
                                                                                              sumOfEdgeWeights [ i ]
                                         int clusterId = nodes2cluster[i];
                                                                                                                                     BEGINFOR VARDEF BEGINFOR ASSIGN
                                         double edgeWeight=graph.getEdgeWeight(nodeId, i);
                                                                                              = 0 ; int
                                                                                                                                     VARDEF ASSIGN BEGINFOR ASSIGN
                                         if (edgeWeight >= threshold) {
                                                                                               [ ] adjacentNodes =
                                           sumOfEdgeWeights[clusterId] += edgeWeight;
                                                                                                                                     ENDFOR ASSIGN ENDFOR ...
                                                                                              graph . getAdNodes (
                Else
Then
                                                                                              nodeId ) ; for
                                       //and determine the cluster of biggest sum.
                                                                                               (int i:
                                       int newClusterNumber=nodes2cluster[nodeId];
                                                                                              adjacentNodes ) { int
                                       double maxWeight=0;
                Operator
                                       for (int i =0; i < sumOfEdgeWeights.length; i++)</pre>
                                                                                              clusterId = nodes2clu [
                                                                                              i ] ; double
                                        if((sumOfEdgeWeights[i])>maxWeight){
                                                                                              edgeWeight = graph .
                                          newClusterNumber=i;
                                          maxWeight=sumOfEdgeWeights[i];
                                                                                              getEdgeWeight ( nodeId ,
                                                                                              i ) ; if
```

Representation d	Sim. measure φ	Compilation level for d	Runtime for φ	
string of token types string of token types string of token types	compression ratio greedy string tiling longest common substring	lexical lexical lexical	$O(\mathbf{d} ^2)$ $O(\mathbf{d} ^3)$ $O(\mathbf{d} ^2)$	[Chen et al. 2004] [Prechelt et al. 2000] [Burrows et al. 2000]
string of token types	longest common subseq.	lexical	$O(\mathbf{d} ^2)$	[new]

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF VARDEF ASSIGN CASE BEGINSWITCH BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF VARDEF ASSIGN CASE BEGINSWITCH BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF VARDEF ASSIGN CASE BEGINSWITCH BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF VARDEF ASSIGN CASE BEGINSWITCH BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF VARDEF ASSIGN CASE BEGINSWITCH BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF | VARDEF ASSIGN | CASE BEGINSWITCH | BEGINFOR ASSIGN

Comparison of Structure-based String Models

For "compression ration", "greedy string tiling", and "longest common substring" the heart of φ is substring maximization.

BEGINFOR VARDEF BEGINFOR ASSIGN VARDEF ASSIGN BEGINFOR ASSIGN

BEGINFOR VARDEF | VARDEF ASSIGN | CASE BEGINSWITCH | BEGINFOR ASSIGN

Longest common subsequence:

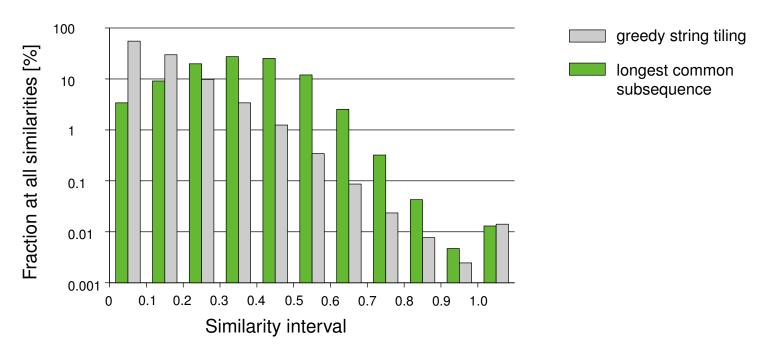
$$\varphi(\mathbf{s}_q, \mathbf{s}_x) = \frac{2 \cdot |\mathsf{lcs}(\mathbf{s}_q, \mathbf{s}_x)|}{|\mathbf{s}_q| + |\mathbf{s}_x|}$$

Comparison of Structure-based String Models

Corpus:

- open source project JNode, (Java New Operating System Design Effort)
- □ 18 subsequent release versions, 80 091 documents
- 121 215 methods

Experiment (plot below): sample of 50 000 method pairs, drawn i.i.d.



Fingerprint-based Models

```
//subloop. for each node...
for(int nodeIndex = 0; nodeIndex<n; nodeIndex++) {</pre>
    int nodeId = nodeIdPermutation[nodeIndex];
    //System.out.println("node: "+nodeId);
    //reset sums.
    for(int i=0; i<n; i++) sumOfEdgeWeights[i]=0;</pre>
    //sum all the edges going out to the same cluster
    int[] adjacentNodes = graph.getAdjacentNodes(nodeId);
    for (int i : adjacentNodes)
      int clusterId = nodes2cluster[i];
      double edgeWeight=graph.getEdgeWeight(nodeId, i);
      if (edgeWeight >= threshold) {
        sumOfEdgeWeights[clusterId] += edgeWeight;
    //and determine the cluster of biggest sum.
    int newClusterNumber=nodes2cluster[nodeId];
    double maxWeight=0;
   for(int i =0; i<sumOfEdgeWeights.length; i++)</pre>
     if((sumOfEdgeWeights[i])>maxWeight){
       newClusterNumber=i;
       maxWeight=sumOfEdgeWeights[i];
```

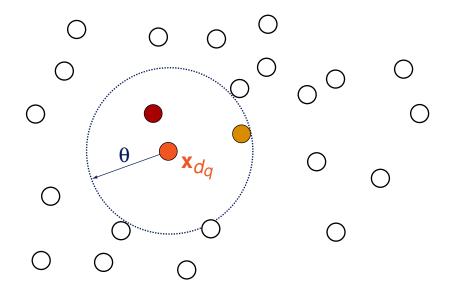
```
for ( int nodeIndex
= 0 : nodeIndex
< n : nodeIndex++
) { int nodeId
= nodeIdPer [ nodeIndex
] ; for (
: i < n
; i ++ )
sumOfEdgeWeights [ i ]
= 0 ; int
[ ] adjacentNodes =
graph . getAdNodes (
nodeId ) ; for
(int i:
adjacentNodes ) { int
clusterId = nodes2clu
i ] ; double
edgeWeight = graph .
getEdgeWeight ( nodeId ,
i ) ; if
```

{2323753332,345256745}

Rationale:

- the inherent quadratic situation becomes linear
- code repositories become extremely large
- because of the problem structure we are interested in plagiarism candidates; a human inspection is always necessary

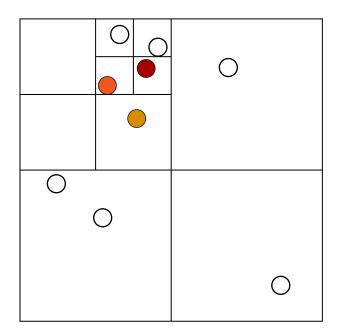
Nearest Neighbor Search



Applications:

- elimination of duplicates / near duplicates
- identification of versioned and plagiarized documents
- retrieval of similar documents
- identification of source code plagiarism

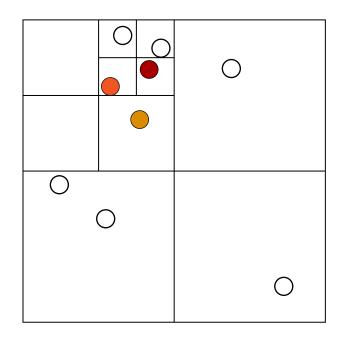
Nearest Neighbor Search

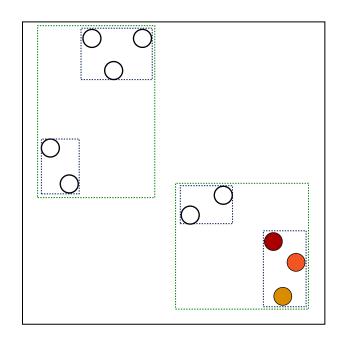


Indexing with space partitioning methods:

- Quad-tree.
 - Split the space recursively into sub-squares until only a few points left. Space exponential in dimension; time exponential in dimension.
- □ Kd-tree. Linear space; exponential query time is still possible.

Nearest Neighbor Search





Indexing with data partitioning methods:

□ R-tree.

Bottom-up; heuristically construct minimum bounding regions for points Works well for low dimensions (< 10).

□ Rf-tree, X-tree, . . .

Document Representation and Search

The nearest neighbor problem cannot be solved efficiently in high dimensions by partitioning methods.

"Existing methods are outperformed on average by a simple sequential scan, if the number of dimensions exceeds around 10."

[Weber 99, Gionis/Indyk/Motwani 99-04]

Document Representation and Search

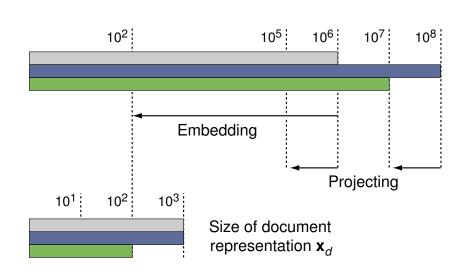
The nearest neighbor problem cannot be solved efficiently in high dimensions by partitioning methods.

"Existing methods are outperformed on average by a simple sequential scan, if the number of dimensions exceeds around 10."

[Weber 99, Gionis/Indyk/Motwani 99-04]

English Wikipedia:

Dictionary	Number of dimensions
1-gram space	3 921 588
4-gram space	274 101 016
8-gram space	373 795 734
Shingling space	75 659 644



Document Representation and Search

Given the representation \mathbf{x}_{d_q} of a query document and a collection D.

- □ Linear comparison under some BOW representation
 - → Similarity ranking (baseline)



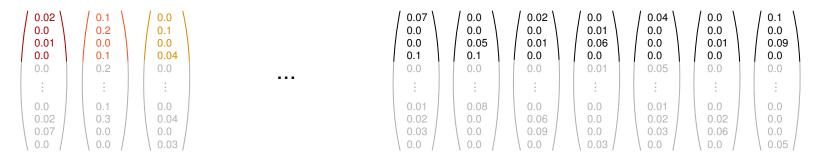
. . .

	0.07 0.0 0.0 0.0 0.1	\	0.0 0.0 0.05 0.1	\	0.02 \ 0.0 0.01 0.0	\	0.0 \ 0.01 0.06 0.0	\	0.04 0.0 0.0 0.0 0.0	\	0.0 0.0 0.01 0.0	\	0.1 0.0 0.09 0.0	١
	0.0	١ ١	0.0		0.0	1 1	0.01	1 1	0.05		0.0	1 1	0.0	
	÷		÷		÷		÷		÷		÷		÷	
	0.01		0.08	11	0.0	$I \setminus I$	0.0	$I \setminus I$	0.01	1 1	0.0	1 \	0.0	
l	0.02		0.0	1 \	0.06	Π	0.0	Π	0.02	1 \	0.02	1 1	0.0	
١	0.03		0.0	'\	0.09	۱ ۱	0.0	/ \	0.03	/ \	0.06	۱ ۱	0.0	I
١	0.0		\ 0.0 /	1	0.0		0.03		0.0		0.0		0.05	ı

Document Representation and Search

Given the representation \mathbf{x}_{d_q} of a query document and a collection D.

- □ Linear comparison under some BOW representation
 - → Similarity ranking (baseline)
- □ Linear comparison under some compact representation
 - \rightarrow Acceptable similarity ranking (85% recall at $\varphi > 0.5$)



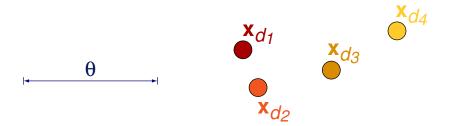
Document Representation and Search

Given the representation \mathbf{x}_{d_q} of a query document and a collection D.

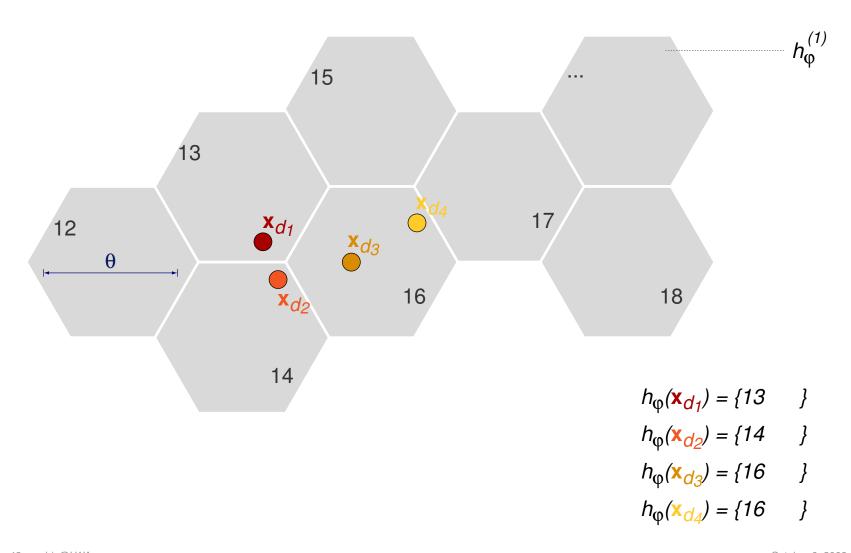
- □ Linear comparison under some BOW representation
 - → Similarity ranking (baseline)
- □ Linear comparison under some compact representation
 - \rightarrow Acceptable similarity ranking (85% recall at $\varphi > 0.5$)
- $lue{}$ Comparison in constant time with a similarity-sensitive hash function h_{arphi}
 - ightharpoonup Binary decision wrt. threshold θ (similar if $\varphi > \theta$ / not similar if $\varphi \leq \theta$)

124298	456723	546781	342509	129842	972653	921345	546719	564214	519461	
0.02	0.1	0.0	0.07	0.0	0.02	0.0	0.04	0.0	0.1	
0.01	0.0	0.0	0.0	0.05	0.01	0.06	0.0	0.01	0.09	
0.0	0.1	0.04	0.1	0.1	0.0	0.0	0.0	0.0	0.0	
:	:	:	 :	:	:	:	:	:	:	
0.0	0.1	0.0	0.01	0.08	0.0	0.0	0.01	0.0	0.0	
0.02	0.3	0.04	0.02	0.0	0.06	0.0	0.02	0.02	0.0	
\ 0.0 /	\ 0.0 /	\ 0.03 /	\ 0.0 /	\ 0.0 /	\ 0.0 /	0.03	\ 0.0 /	\ 0.0 /	0.05	

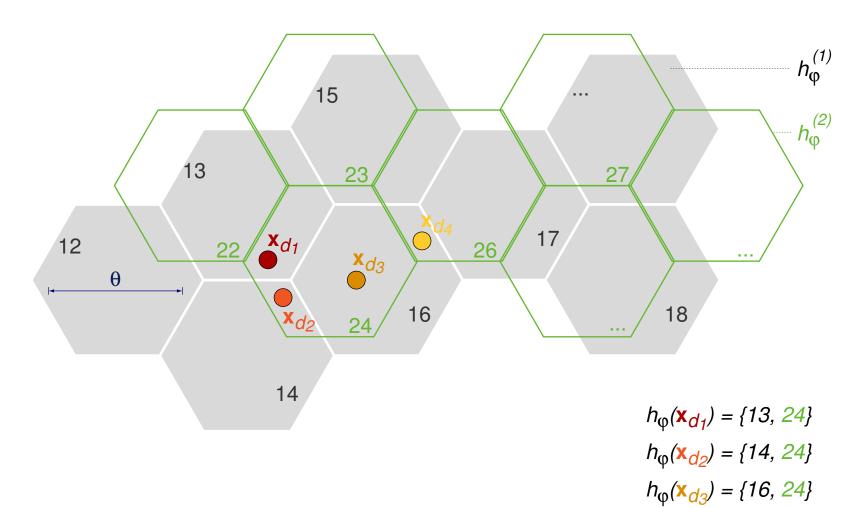
Hash-based Search is a Space Partitioning Method



Hash-based Search is a Space Partitioning Method



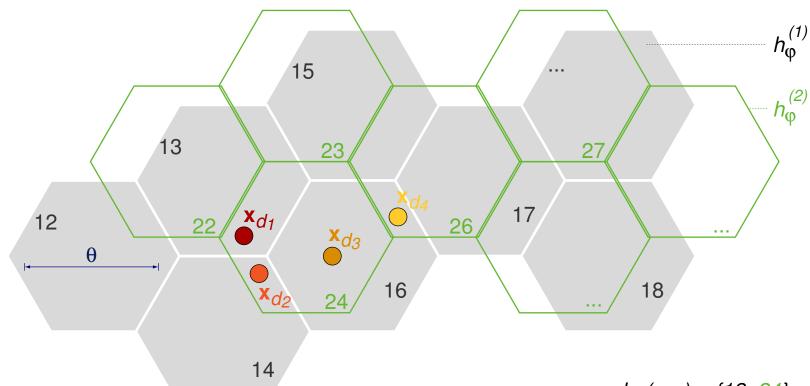
Hash-based Search is a Space Partitioning Method



44 webis@LWA October 6, 2008

 $h_{\Phi}(\mathbf{x}_{O_4}) = \{16, 26\}$

Hash-based Search is a Space Partitioning Method



Similarity collision condition:

$$(h_{\varphi}^*(\mathbf{x}_{d_1}) \cap h_{\varphi}^*(\mathbf{x}_{d_2})) \neq \emptyset \quad \Leftrightarrow \quad \varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) > \theta$$

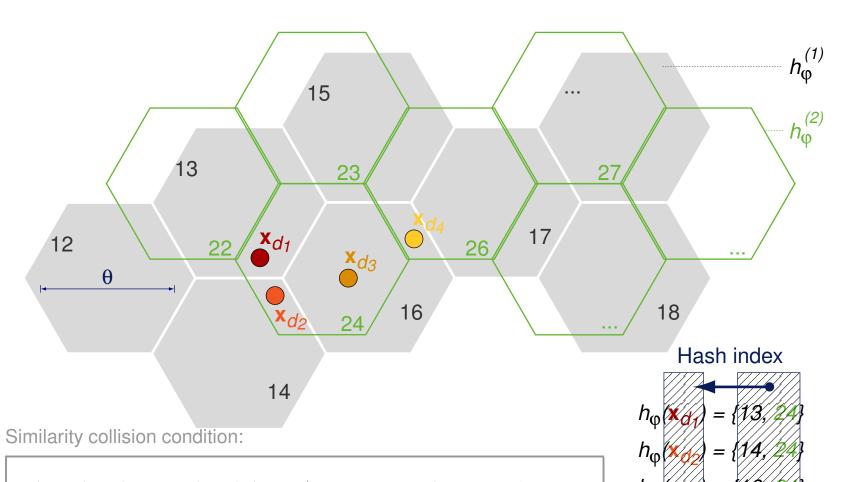
$$h_{\varphi}(\mathbf{x}_{d_1}) = \{13, 24\}$$

 $h_{\varphi}(\mathbf{x}_{d_2}) = \{14, 24\}$

$$h_{\varphi}(\mathbf{x}_{o/3}) = \{16, 24\}$$

$$h_{\phi}(\mathbf{x}_{04}) = \{16, 26\}$$

Hash-based Search is a Space Partitioning Method



 $(h_{\varphi}^*(\mathbf{x}_{d_1}) \cap h_{\varphi}^*(\mathbf{x}_{d_2})) \neq \emptyset \quad \Leftrightarrow \quad \varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) > \theta$

Issues about Hash-based Search

- □ Hash-based search reduces a cont. similarity relation to a binary relation.
- □ Hash-based search is a space partitioning method.
- \Box Space partitioning is realized by a similarity-sensitive hash function h_{φ} .
- \Box Equal codes under h_{φ} indicate similar objects with a high probability.

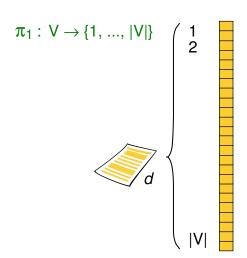
Precision:
$$h_{\varphi}(\mathbf{x}_{d_1}) \cap h_{\varphi}(\mathbf{x}_{d_2}) \neq \emptyset \implies P(\varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) > \theta)$$
 is high

 $\ \square \ h_{\varphi}$ maps similar objects on equal codes with a high probability.

- \Box h_{φ} must be multi-valued if D is partly unknown.
- \Box A perfectly similarity-sensitive hash function h_{ω}^* may exist for each D.

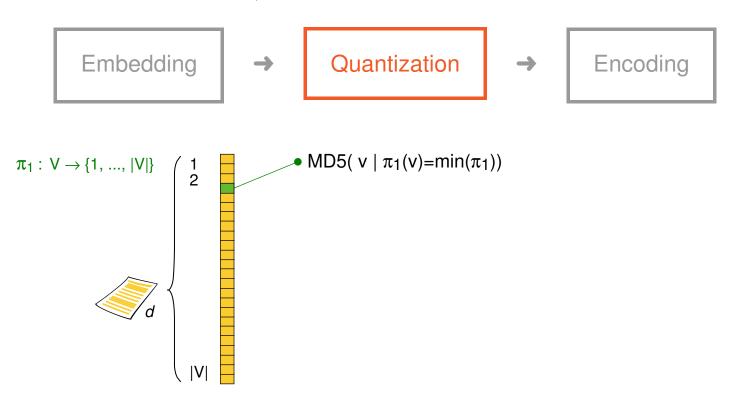
Construction Principles for h_{φ} : Shingling [Broder 2000]





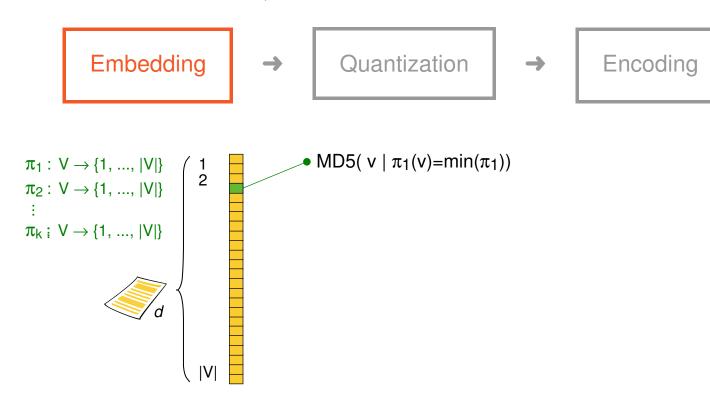
Synchronized random projection

Construction Principles for h_{φ} : Shingling [Broder 2000]



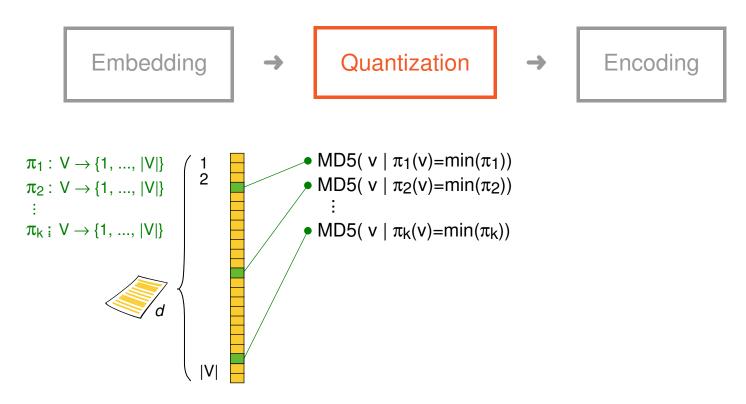
Synchronized random projection

Construction Principles for h_{φ} : Shingling [Broder 2000]



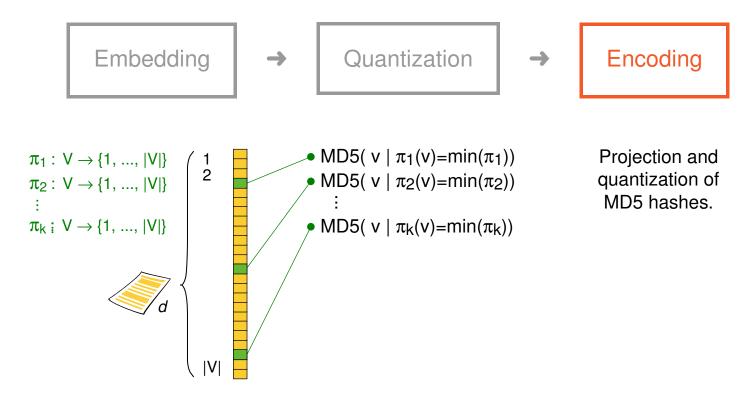
Synchronized random projection

Construction Principles for h_{φ} : Shingling [Broder 2000]



Synchronized random projection

Construction Principles for h_{φ} : Shingling [Broder 2000]

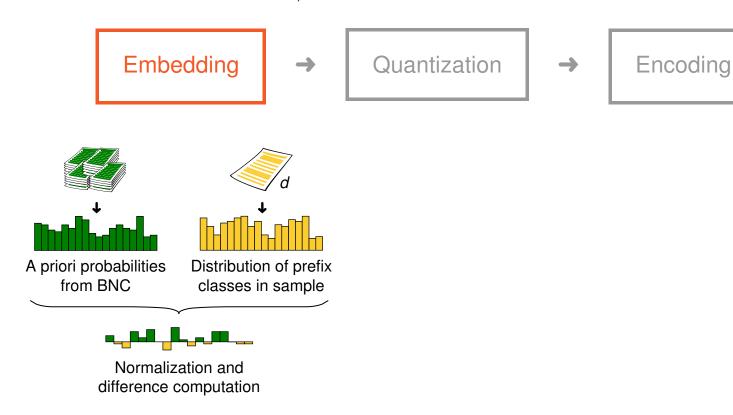


Synchronized random projection

"Super-shingling"

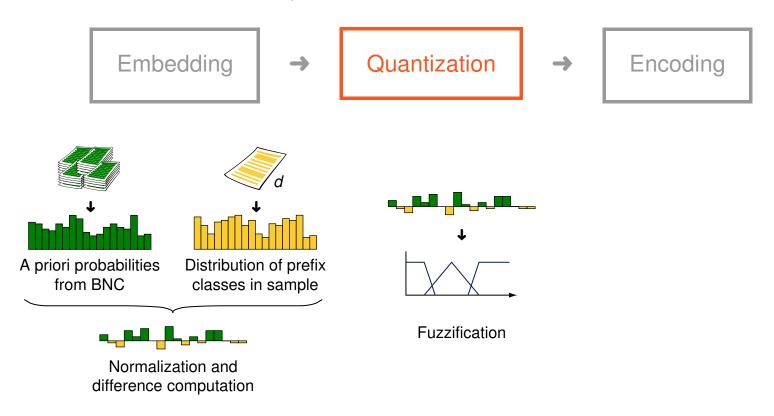
→ Fingerprint = {2643256, 325567} = $h_{\varphi}(\mathbf{x}_d)$

Construction Principles for h_{φ} : Fuzzy-Fingerprinting



Documents from the British National Corpus

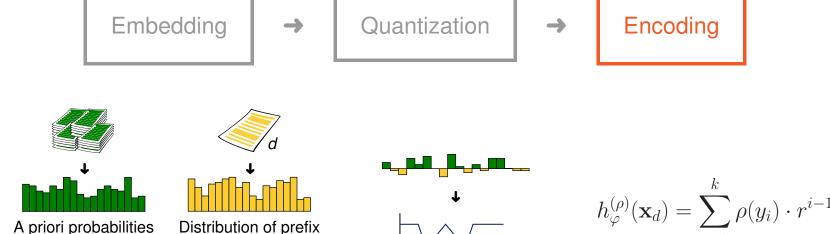
Construction Principles for h_{φ} : Fuzzy-Fingerprinting



Documents from the British National Corpus

Construction Principles for h_{φ} : Fuzzy-Fingerprinting

classes in sample



Normalization and

difference computation

Fuzzification

$$h_{\varphi}^{(\rho)}(\mathbf{x}_d) = \sum_{i=1}^{\kappa} \rho(y_i) \cdot r^{i-1}$$

Documents from the British National Corpus

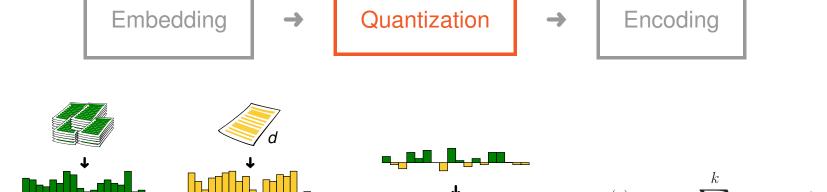
→ Fingerprint = {2643256,

from BNC

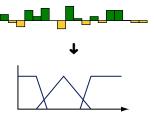
Construction Principles for h_{φ} : Fuzzy-Fingerprinting

Distribution of prefix

classes in sample



Normalization and difference computation



$$h_{\varphi}^{(\rho)}(\mathbf{x}_d) = \sum_{i=1}^k \rho(y_i) \cdot r^{i-1}$$

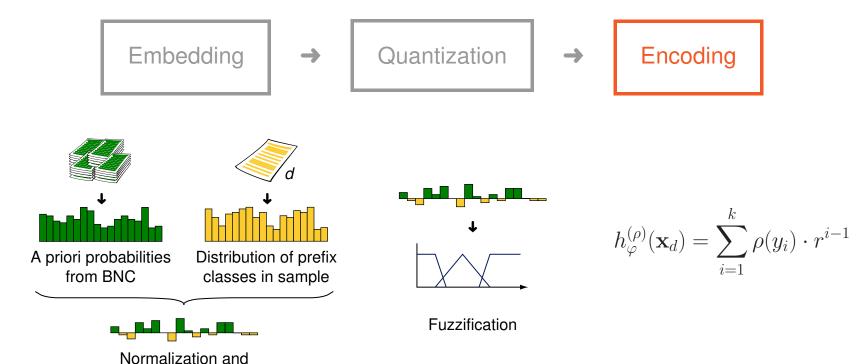
Documents from the British National Corpus

→ Fingerprint = {2643256,

A priori probabilities

from BNC

Construction Principles for h_{φ} : Fuzzy-Fingerprinting



Documents from the British National Corpus

difference computation

→ Fingerprint = {2643256, 325567} = $h_{\varphi}(\mathbf{x}_d)$

Properties of h_{φ}

Code length controls precision.

The collision probability $P(h_{\varphi}(\mathbf{x}_{d_1}) \cap h_{\varphi}(\mathbf{x}_{d_2}) \neq \emptyset \mid \varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) \leq \theta)$ goes down if

- \Box the number k of random vectors (p-stable LSH)
- \Box the number k of prefix classes (Fuzzy-fingerprinting)
- **u** ...

is increased.

Properties of h_{φ}

Code length controls precision.

The collision probability $P(h_{\varphi}(\mathbf{x}_{d_1}) \cap h_{\varphi}(\mathbf{x}_{d_2}) \neq \emptyset \mid \varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) \leq \theta)$ goes down if

- \Box the number k of random vectors (p-stable LSH)
- \Box the number k of prefix classes (Fuzzy-fingerprinting)
- **u** ...

is increased.

Code multiplicity controls recall.

The collision probability $P(h_{\varphi}(\mathbf{x}_{d_1}) \cap h_{\varphi}(\mathbf{x}_{d_2}) \neq \emptyset \mid \varphi(\mathbf{x}_{d_1}, \mathbf{x}_{d_2}) > \theta)$ goes up if

- \Box the number l of vector sets (p-stable LSH)
- \Box the number l of fuzzification schemes (Fuzzy-fingerprinting)
- **u** ...

is increased.

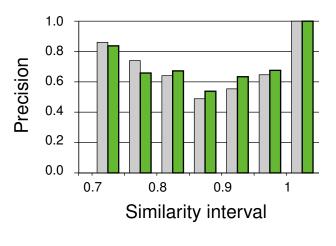
Retrieval Models for Source Code

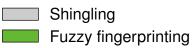
Fingerprint-based Models

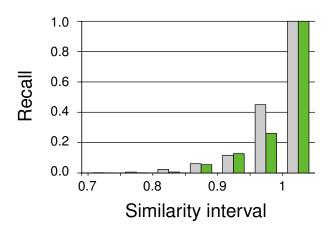
Corpus: as before

Experiment (plot below): 200 queries against fingerprinted corpus

Baseline: greedy string tiling







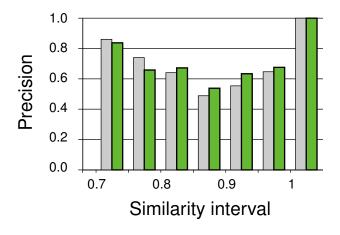
Retrieval Models for Source Code

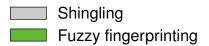
Fingerprint-based Models

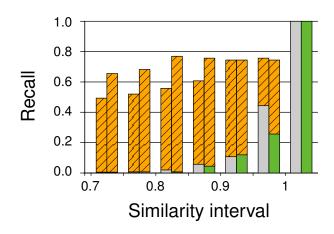
Corpus: as before

Experiment (plot below): 200 queries against fingerprinted corpus

Baseline: greedy string tiling





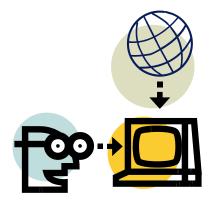


Retrieval of text documents

Summary

Summary

- 1. Survey of retrieval models for high-similarity search in source code.
- 2. We propose the longest common subsequence for the class of structure-based string models:
 - □ better suited for short source code fragments
 - $\Box \varphi$ computation in $O(|\mathbf{d}|^2)$ instead of in $O(|\mathbf{d}|^3)$
- 3. We investigate the use of hash-based search high-similarity search in source code:
 - basis is the class of structure-based string models
 - □ real-world order of magnitudes become possible
 - the ad-hoc application of existing technology leads to unsatisfying recall



Thank you!