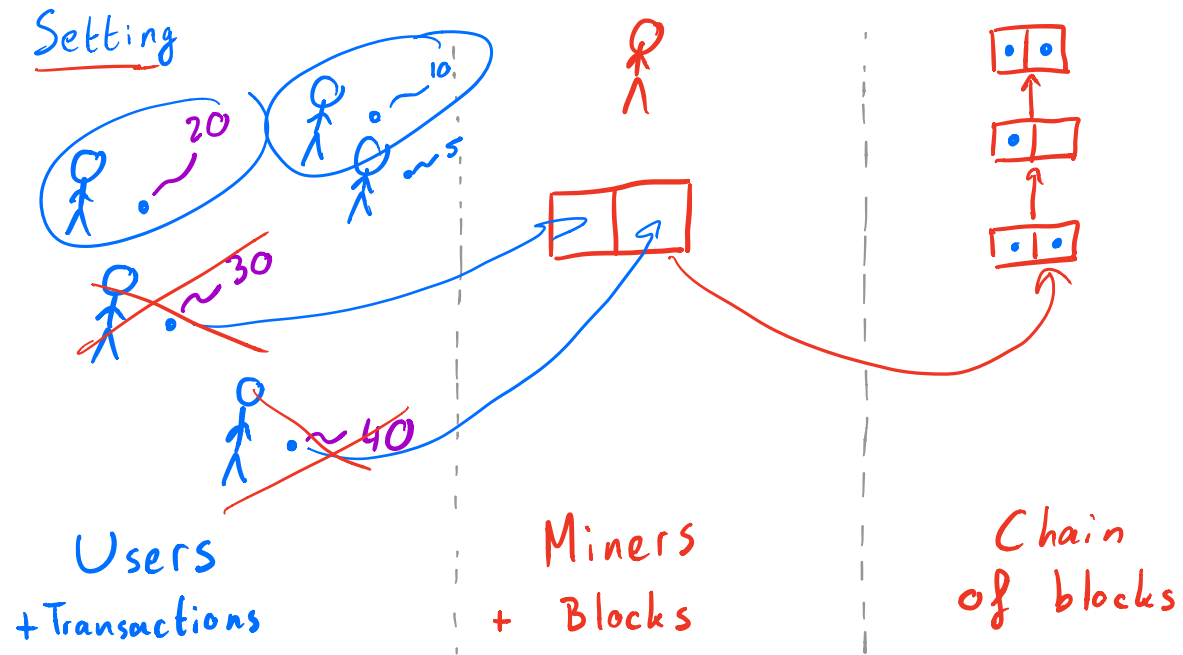


Ethereum's new transaction

fee market design

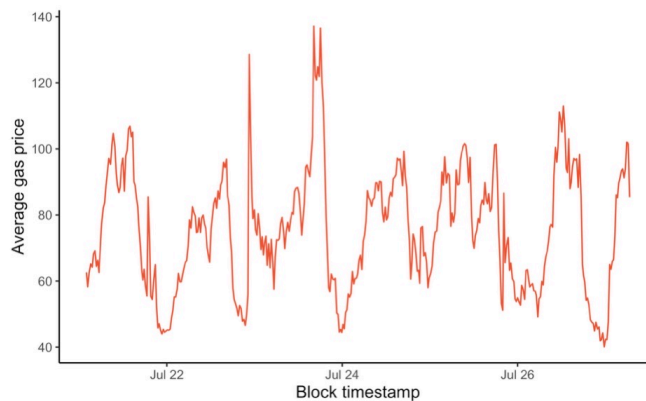
EIP 1559



→ Transaction inclusion is a first-price auction.

→ User bids f ↔ Miner receives f

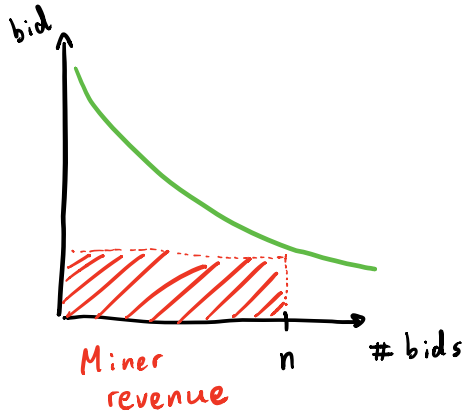
→ FPA ↔ inefficient !



Why 2nd price auctions fail

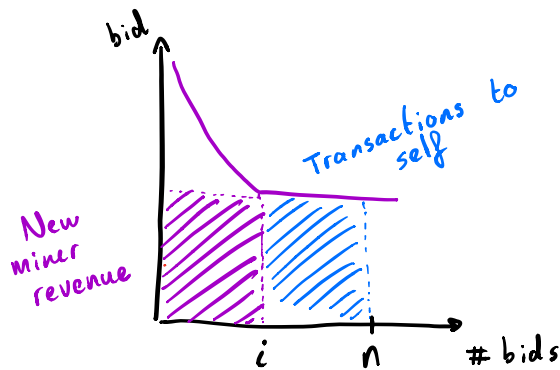
→ Bids in a block are f_1, \dots, f_n

→ Miner receives $n \times f_n$



Problem!

→ Miner can include transactions to itself and increase n -th bid!



Given bids

$$f_1 \geq f_2 \geq \dots \geq f_n$$

Miner optimises

$$\max_{f_i} i \times f_i$$

→ Waste of space!
+ not incentive compatible

\$ 1,000,000	→	\$ 1,000,000
\$ 0.10		\$ 1,000,000
\$ 2 x 0.10		\$ 1,000,000

The key problem

As long as miners decide the entry price, they can game it

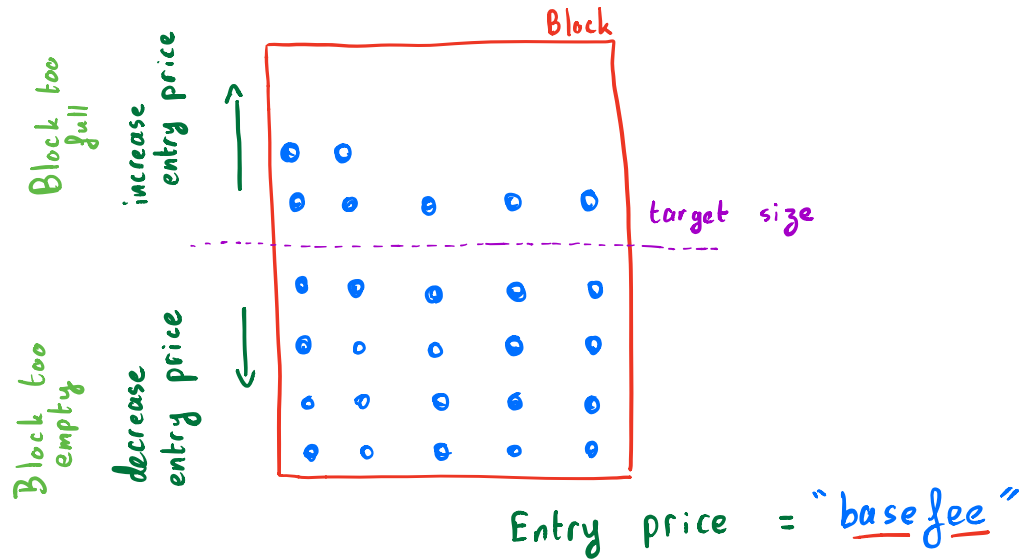
→ Solution : make the entry price "objective".

Dynamic congestion pricing



- Same idea as ERP
- Observe congestion
- When congestion high,
raise entry price
- When congestion low,
lower entry price

Congestion observability

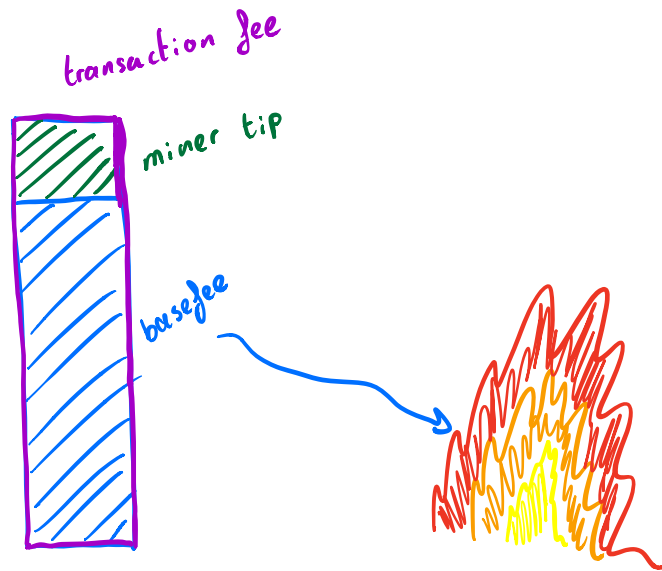


What users pay

Users must pay at least basefee.

Can add a tip to the miner

Basefee is burnt!



Comparing both paradigms

FPA

- Users observe historical prices (+ current mempool state) and decide on f
- Inclusion depends on other bidders directly,
unpredictable

Bad UX!

EIP 1559

- Users observe basefee
- Tip can be set as small compensation
- Most likely included in the next block. Good UX!
(except during spikes!)

Simulations

Stationary case

→ Demand is a Poisson process, $\mathbb{E} = \lambda$

→ Users have a value $v \sim F_v$

cost for waiting $c \sim F_c$

$$v - 5c - f \quad \uparrow \text{basefee + tip}$$

→ See notebook

Strategic users

During demand shifts, users want to get ahead.

Enter a bidding war

→ See notebook

Questions

→ Efficiency of the mechanism
vs. FPA, another 2nd price proposal, "escalator"

→ Miner collusion
Artificially keep basefee low to obtain larger tips

→ Equilibrium behaviour of the users
how to set tip / max fee

→ Approximations (wallet defaults)

→ Miner transaction pool behaviour