


TLA⁺ Video Course – Lecture 4

Leslie Lamport

DIE HARD

This video should be viewed in conjunction with a Web page.
To find that page, search the Web for *TLA+ Video Course*.

The TLA⁺ Video Course
Lecture 4
Die Hard



In this video, you'll learn how TLC can save your life. . . if you ever find yourself in the middle of a Hollywood action movie.

This will require you to learn some more about TLA+, TLC, and the Toolbox — which could turn out to be useful even outside of Hollywood.

[slide 2]

THE DIE HARD PROBLEM

Die Hard 3

Die Hard 3 is

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A 1995 action film.

Die Hard 3 is a 1995 action film starring Bruce Willis and Samuel L. Jackson as the heroes.

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The heroes had to put exactly 4 gallons of water in a jug.

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1 U.S. gallon = $\frac{1}{64}$ hogshead

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1 U.S. gallon = $\frac{1}{64}$ hogshead = 3.785411784 liters

Die Hard 3 is a 1995 action film starring Bruce Willis and Samuel L. Jackson as the heroes.

To disarm a bomb, they had to put exactly 4 gallons of water in a jug.

Die Hard 3

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The heroes had to put exactly 4 gallons of water in a jug.

They had a 3 gallon jug, a 5 gallon jug, and a water faucet.

They were given a 3 gallon jug, a 5 gallon jug, and a water faucet.

Die Hard 3

A 1995 action film.

The heroes had to put exactly 4 gallons of water in a jug.

They had a 3 gallon jug, a 5 gallon jug, and a water faucet.

Search the Web for: *Die Hard Jugs Problem YouTube*.

They were given a 3 gallon jug, a 5 gallon jug, and a water faucet.

You can watch the relevant scene by searching the Web for *Die Hard Jugs Problem YouTube*.

There were no markings on the jugs.

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They needed exacty 4 gallons.

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There were no markings on the jugs.

They needed exacty 4 gallons.

Not 3.99 or 4.01.

There were no markings on the jugs.

They needed exacty 4 gallons.

Not 3.99 or 4.01 gallons.

GETTING STARTED

Getting Started on a Spec

When we want to write a spec, what should we do first?

Getting Started on a Spec

The best way:

Write a single correct behavior.

When we want to write a spec, what should we do first?

I recommend writing the start of a single correct behavior.

Getting Started on a Spec

The best way:

Write a single correct behavior.

Informally.

When we want to write a spec, what should we do first?

I recommend writing the start of a single correct behavior.

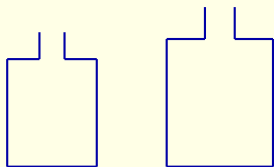
Informally at first.

This isn't a big budget movie.

This isn't a big budget Hollywood movie, and I can't afford big jugs.

This isn't a big budget movie.

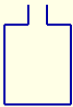
I'll use these cartoon jugs:



This isn't a big budget Hollywood movie, and I can't afford big jugs.

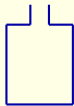
So instead, I'll illustrate the spec with these cartoon jugs.

They start with both jugs empty.



Our heroes start with both jugs empty.

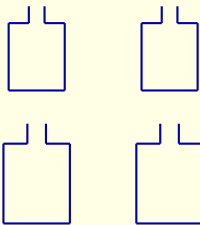
It's possible for them to pour 4 gallons of water into the 5-gallon jug.



Our heroes start with both jugs empty.

It's possible for them to solve their problem by simply pouring 4 gallons of water into the 5-gallon jug.

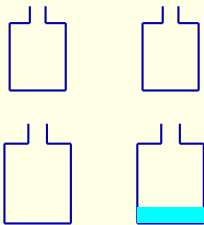
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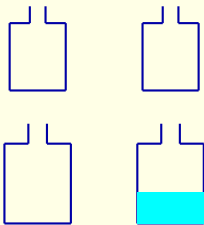
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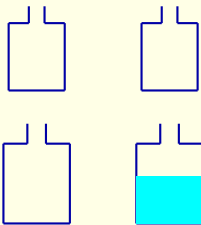
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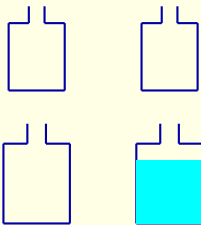
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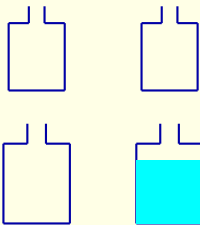
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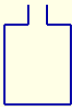
But they'd have to be very lucky to get exactly 4 gallons.

Our heroes start with both jugs empty.

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But they'd have to be very lucky to get exactly 4 gallons.

We only allow behaviors where they know exactly how much water is in each jug.



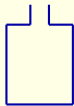
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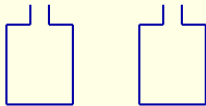
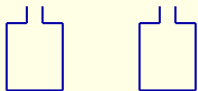
So we only allow behaviors in which they always know exactly how much water is in each jug.

Initially, both jugs empty.



Again, they start with both jugs empty.

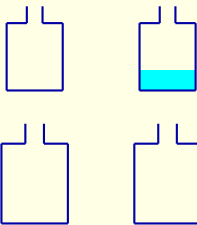
Fill 3-gal. jug.



Again, they start with both jugs empty.

The only thing they can do now is fill a jug. Suppose they fill the 3 gallon jug.

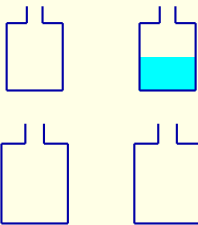
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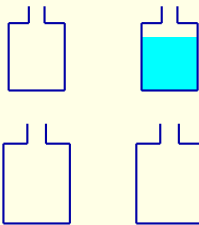
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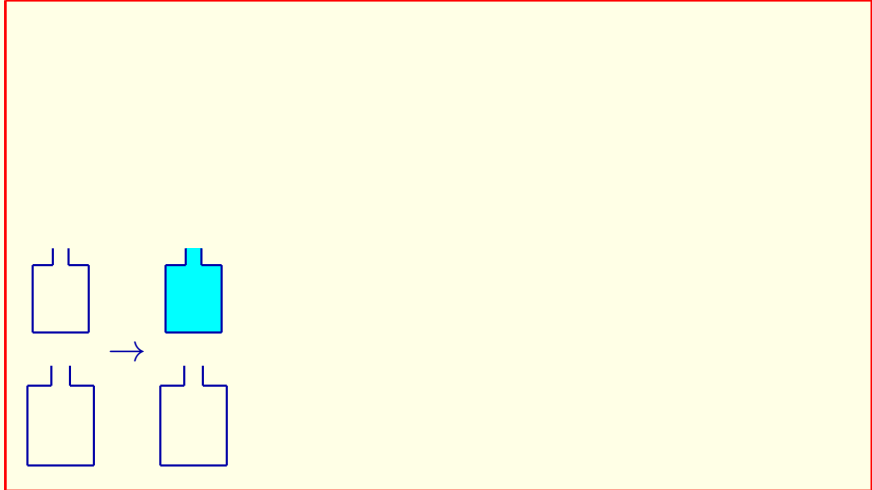
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Fill 3-gal. jug.



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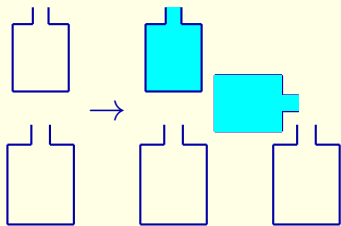
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Empty 3-gal. jug into 5-gal. jug.

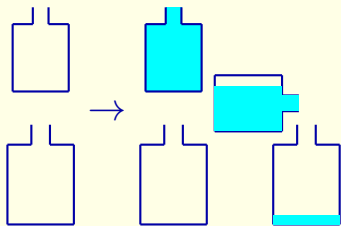


Again, they start with both jugs empty.

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Next, they empty the water from the 3 gallon jug into the 5 gallon jug.

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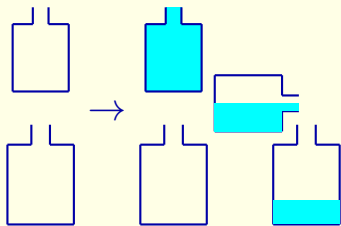


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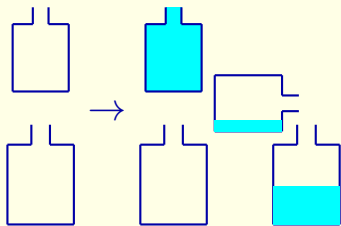


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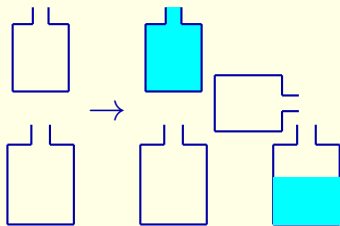


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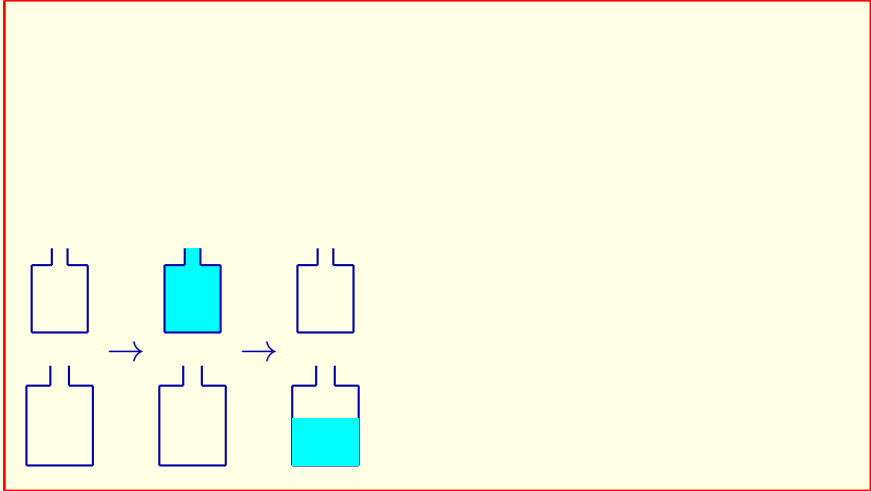
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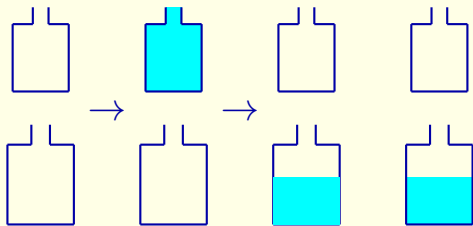


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Fill 3-gal. jug.



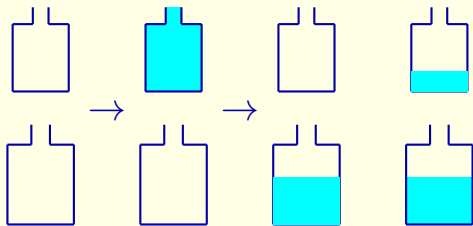
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Now they fill the 3 gallon jug.

Fill 3-gal. jug.



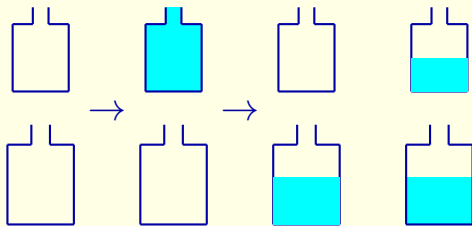
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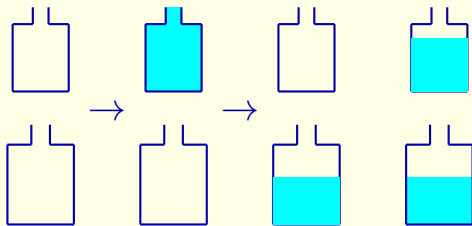
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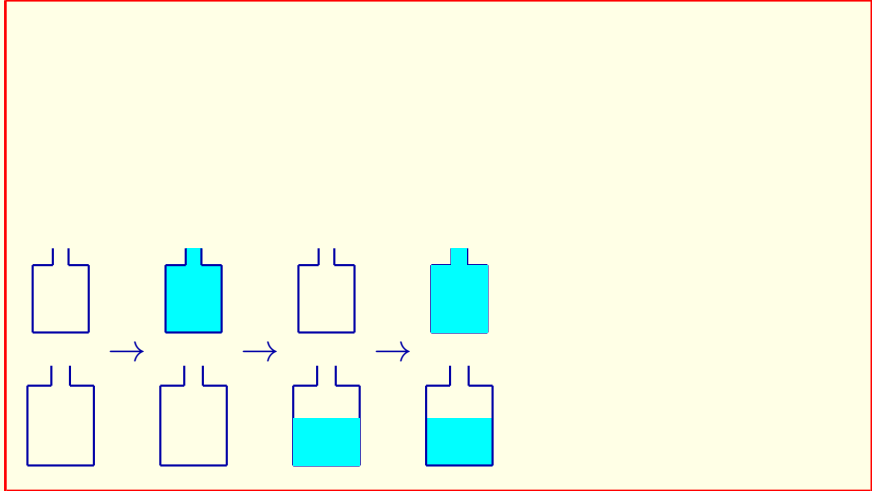


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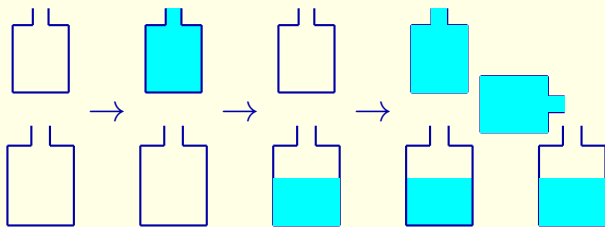
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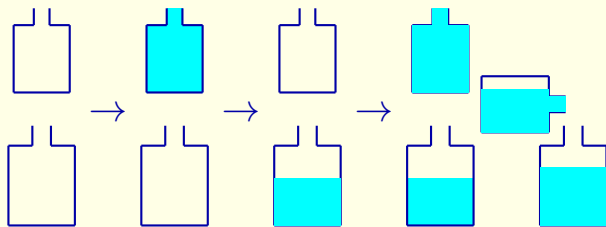
Now they fill the 3 gallon jug.

Fill 5-gal. jug from 3-gal. jug.



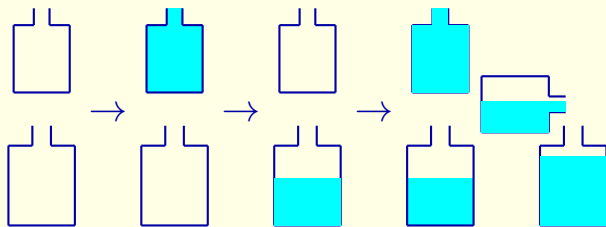
Now they fill the 5 gallon jug from the 3 gallon jug.

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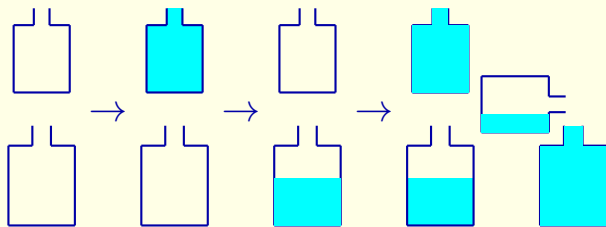
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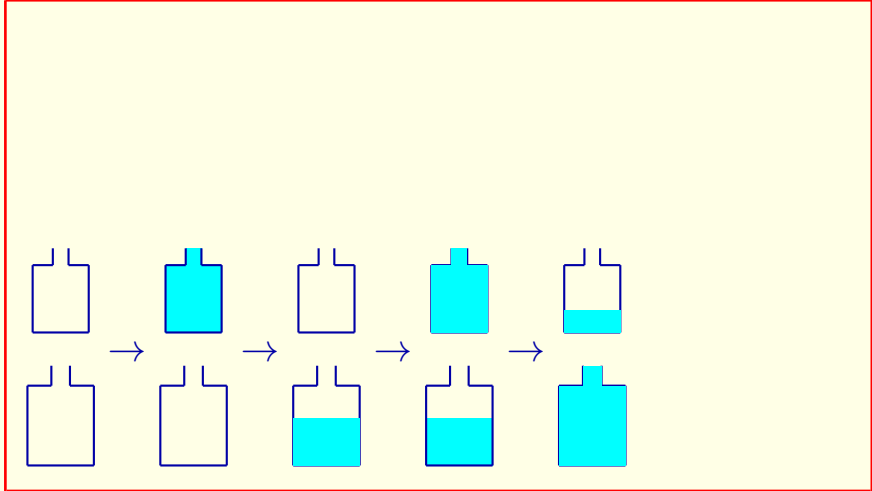


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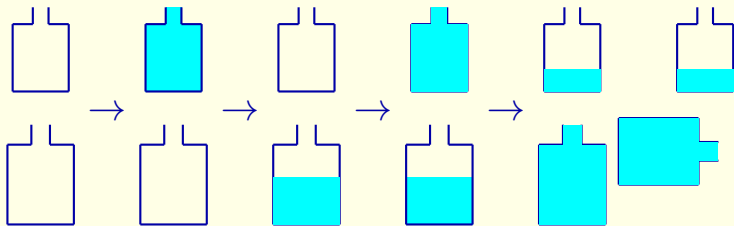


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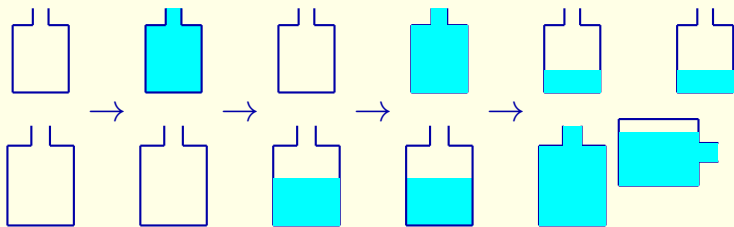
Empty 5-gal. jug.



Now they fill the 5 gallon jug from the 3 gallon jug.

They then empty the 5-gallon jug onto the ground.

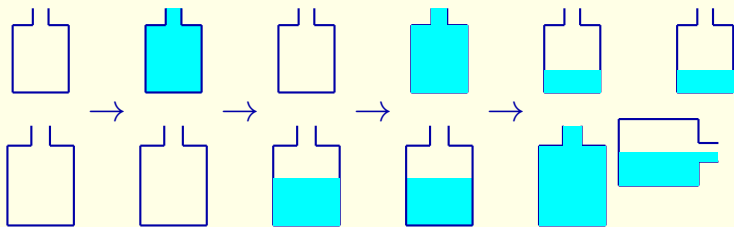
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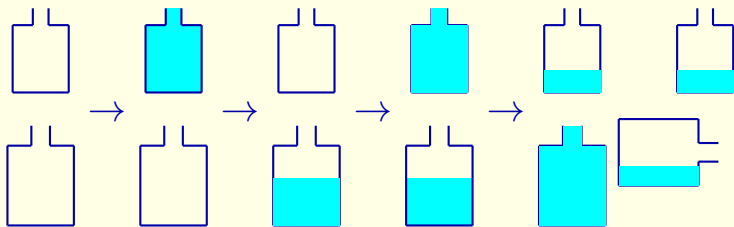
Empty 5-gal. jug.



Now they fill the 5 gallon jug from the 3 gallon jug.

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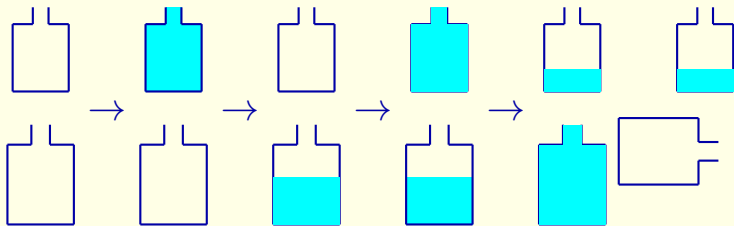
Empty 5-gal. jug.



Now they fill the 5 gallon jug from the 3 gallon jug.

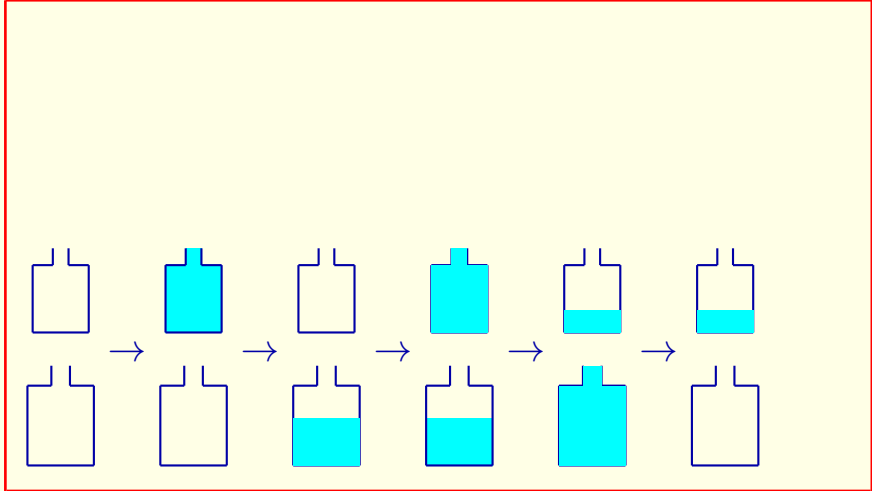
They then empty the 5-gallon jug onto the ground.

Empty 5-gal. jug.



Now they fill the 5 gallon jug from the 3 gallon jug.

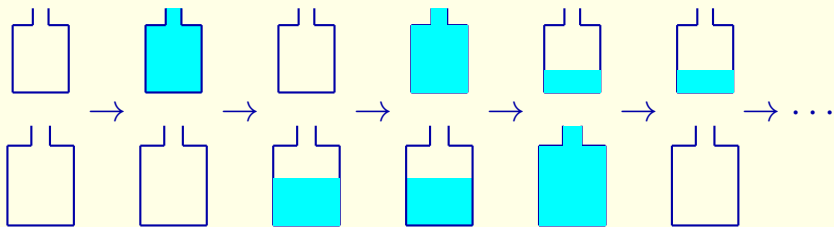
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Now they fill the 5 gallon jug from the 3 gallon jug.

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And so on.

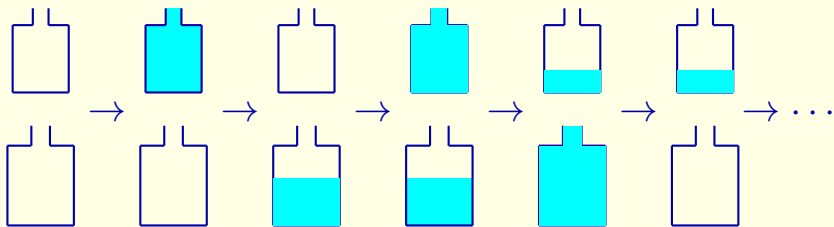


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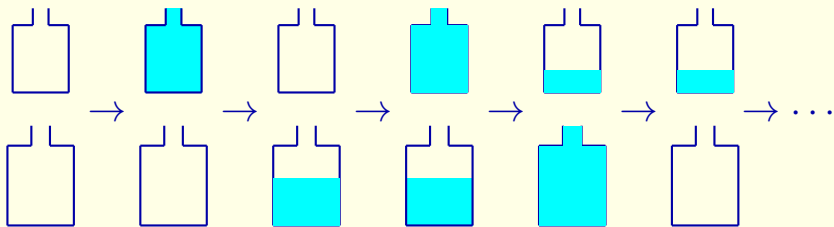
This is just one possible behavior.



This is just one of many possible ways a behavior can begin.

This is just one possible behavior.

Let's write it more formally.



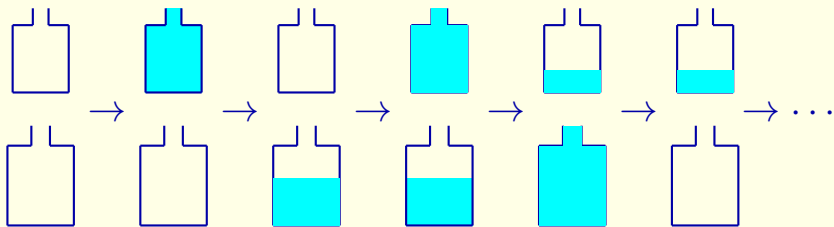
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Let's write it more formally.

Let values of *small* and *big* represent number of gallons in each jug.

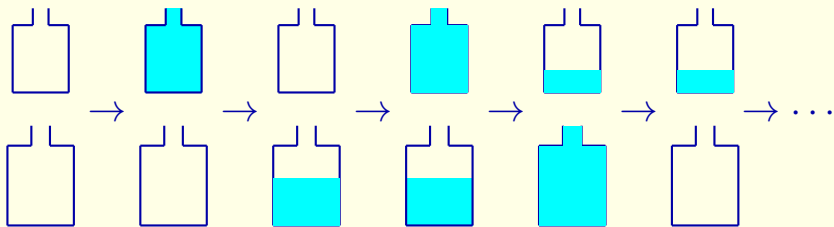


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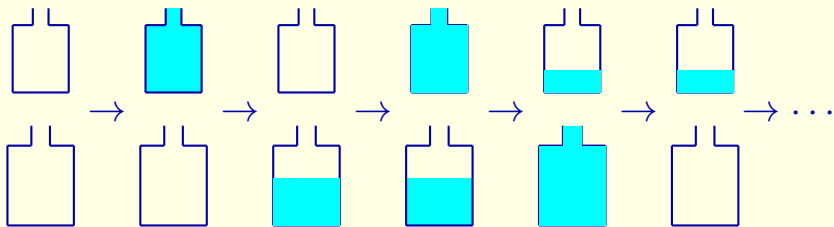
Let the values of the variables *small* and *big* represent the number of gallons of water in each jug.

$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix}$



Initially, both jugs have 0 gallons of water.

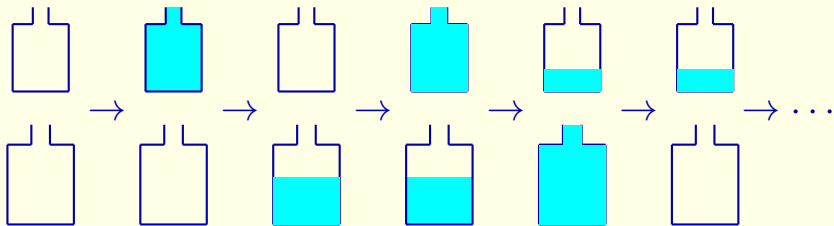
Fill small jug.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \quad \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix}$$


Initially, both jugs have 0 gallons of water.

Filling the small jug puts 3 gallons of water in it.

Empty small jug into big jug.

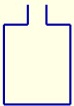
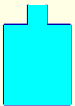
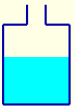
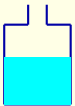
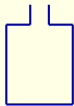
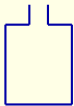
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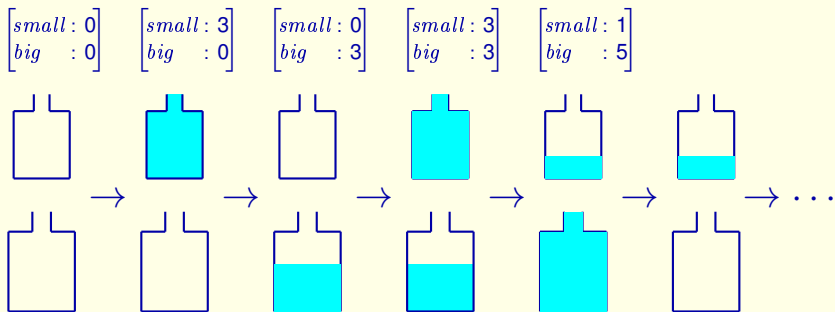
Those 3 gallons are transferred from the small jug to the big jug.

Fill small jug.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix}$$
$$\begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix}$$
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3 gallons are then added to the small jug.

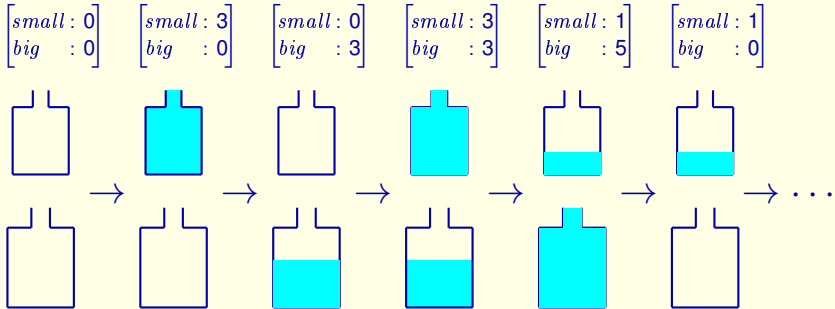
Fill big jug from small jug.



3 gallons are then added to the small jug.

The big jug is then filled from the small jug, putting 5 gallons in the big jug and leaving 1 gallon in the small jug.

Empty big jug.



3 gallons are then added to the small jug.

The big jug is then filled from the small jug, putting 5 gallons in the big jug and leaving 1 gallon in the small jug.

The big jug is then emptied, leaving 0 gallons in it.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

3 gallons are then added to the small jug.

The big jug is then filled from the small jug, putting 5 gallons in the big jug and leaving 1 gallon in the small jug.

The big jug is then emptied, leaving 0 gallons in it.

What did we learn from this behavior?

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

What did we learn by writing this behavior?

We learned two things.

What did we learn from this behavior?

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

1. What the variables are.

What did we learn by writing this behavior?

We learned two things.

First, what the variables are.

What did we learn from this behavior?

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

1. What the variables are.
2. What constitutes a step.

What did we learn by writing this behavior?

We learned two things.

First, what the variables are.

And second, what constitutes a step. For example...

Filling a jug is a single step.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

1. What the variables are.
2. What constitutes a step.

Filling a jug is a single step.

Filling a jug is a single step.

No intermediate states.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 2 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix}$$

1. What the variables are.

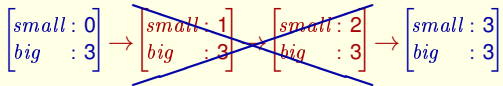
2. What constitutes a step.

Filling a jug is a single step.

There's no intermediate partially-filled state or states.

Filling a jug is a single step.

No intermediate states.



1. What the variables are.

2. What constitutes a step.

Filling a jug is a single step.

There's no intermediate partially-filled state or states.

What did we learn from this behavior?

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

1. What the variables are.
2. What constitutes a step.

Filling a jug is a single step.

There's no intermediate partially-filled state or states.

Simplest abstraction of real jugs and water

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

Filling a jug is a single step.

There's no intermediate partially-filled state or states.

This is the simplest abstraction of the behavior of real jugs and water

Simplest abstraction of real jugs and water

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

for this problem.

Filling a jug is a single step.

There's no intermediate partially-filled state or states.

This is the simplest abstraction of the behavior of real jugs and water for the particular problem faced by our heroes.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

Real specifications are written to eliminate some kinds of errors.

Real specifications are written for a purpose.
Usually to eliminate some particular kinds of errors.

$$\begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 0 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 3 \\ \textit{big} : 3 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 5 \end{bmatrix} \rightarrow \begin{bmatrix} \textit{small} : 1 \\ \textit{big} : 0 \end{bmatrix} \rightarrow \dots$$

Real specifications are written to eliminate some kinds of errors.

Like getting blown up.

Real specifications are written for a purpose.
Usually to eliminate some particular kinds of errors.

For example, to avoid getting blown up.

THE SPECIFICATION

We can now start writing the actual TLA+ specification.

MODULE *DieHard*

Beginning of module.

The spec is in a module called *DieHard*.

MODULE *DieHard*

EXTENDS *Integers*

Imports operators of arithmetic.

The spec is in a module called *DieHard*.

As in our *Simple Program* example, the EXTENDS statement imports operators of arithmetic

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

Declares the variables.

The spec is in a module called *DieHard*.

As in our *Simple Program* example, the EXTENDS statement imports operators of arithmetic

and the VARIABLES statement declares our two variables.

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

TLA⁺ has no type declarations.

The spec is in a module called *DieHard*.

As in our *Simple Program* example, the EXTENDS statement imports operators of arithmetic and the VARIABLES statement declares our two variables.

In TLA+ we don't write type declarations.

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

TLA⁺ has no type declarations.

Type correctness means variables have sensible values.

Type correctness means that all the variables have sensible values.

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

We define a formula that asserts type correctness.

Type correctness means that all the variables have sensible values.

It's a good idea to define a formula that asserts type correctness.

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

We define a formula that asserts type correctness.

Helps to understand spec.

Type correctness means that all the variables have sensible values.

It's a good idea to define a formula that asserts type correctness.

It helps a reader to understand the spec.

MODULE *DieHard*

EXTENDS *Integers*

VARIABLES *small, big*

We define a formula that asserts type correctness.

Helps to understand spec.

TLC can check that it's always true.

Type correctness means that all the variables have sensible values.

It's a good idea to define a formula that asserts type correctness.

It helps a reader to understand the spec.

And TLC can type-check the spec by checking that this formula is always true.

EXTENDS *Integers*

VARIABLES *small, big*

TypeOK \triangleq

We define a formula that asserts type correctness.

Helps to understand spec.

TLC can check that it's always true.

I like to call this formula *TypeOK*.

EXTENDS *Integers*VARIABLES *small, big* $TypeOK \triangleq \wedge small \in 0 .. 3$

I like to call this formula *TypeOK*.

It asserts that the value of *small* is an integer from 0 through 3.

EXTENDS *Integers*VARIABLES *small, big*
$$\begin{aligned} \textit{TypeOK} \triangleq & \wedge \textit{small} \in 0 .. 3 \\ & \wedge \textit{big} \in 0 .. 5 \end{aligned}$$

I like to call this formula *TypeOK*.

It asserts that the value of *small* is an integer from 0 through 3.
and the value of *big* is an integer from 0 through 5.

EXTENDS *Integers*VARIABLES *small, big*

$$\begin{aligned} \textit{TypeOK} \triangleq & \wedge \textit{small} \in 0 .. 3 \\ & \wedge \textit{big} \in 0 .. 5 \end{aligned}$$

This definition is not part of the spec.

I like to call this formula *TypeOK*.

It asserts that the value of *small* is an integer from 0 through 3.
and the value of *big* is an integer from 0 through 5.

This definition is not part of the spec.

EXTENDS *Integers*

VARIABLES *small, big*

$$\begin{aligned} \textit{TypeOK} &\triangleq \wedge \textit{small} \in 0 .. 3 \\ &\quad \wedge \textit{big} \in 0 .. 5 \end{aligned}$$

This definition is not part of the spec.

Removing it doesn't change anything.

I like to call this formula *TypeOK*.

It asserts that the value of *small* is an integer from 0 through 3.
and the value of *big* is an integer from 0 through 5.

This definition is not part of the spec.

Removing it doesn't change anything.

The Initial-State Formula

The initial-state formula.

The Initial-State Formula

$$Init \triangleq$$

The initial-state formula.

As usual, let's name it *Init*.

The Initial-State Formula

$$\begin{aligned} Init &\stackrel{\Delta}{=} \bigwedge big = 0 \\ &\quad \bigwedge small = 0 \end{aligned}$$

The initial-state formula.

As usual, let's name it *Init*.

It asserts that both jugs are empty.

THE NEXT-STATE FORMULA

The next-state formula.

The Next-State Formula

The next-state formula describes all permitted steps.

The next-state formula describes all permitted steps.

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$

The next-state formula describes all permitted steps.

It's usually written as F_1 or F_2 or (and so on) ,

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$,
where each F_i allows a different kind of step.

The next-state formula describes all permitted steps.

It's usually written as F_1 or F_2 or (and so on) ,
where each formula F allows a different kind of step.

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$,
where each F_i allows a different kind of step.

The behavior we wrote has 3 kinds of steps:

The behavior we just wrote has 3 different kinds of steps:

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$,
where each F_i allows a different kind of step.

The behavior we wrote has 3 kinds of steps:

- Fill a jug.

The behavior we just wrote has 3 different kinds of steps:

Steps that fill a jug.

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$,
where each F_i allows a different kind of step.

The behavior we wrote has 3 kinds of steps:

- Fill a jug.
- Empty a jug.

The behavior we just wrote has 3 different kinds of steps:

Steps that fill a jug.

Steps that empty a jug.

The Next-State Formula

The next-state formula describes all permitted steps.

It's usually written as $F_1 \vee F_2 \vee \dots \vee F_n$,
where each F_i allows a different kind of step.

The behavior we wrote has 3 kinds of steps:

- Fill a jug.
- Empty a jug.
- Pour from one jug into the other.

The behavior we just wrote has 3 different kinds of steps:

Steps that fill a jug.

Steps that empty a jug.

And steps that pour from one jug into the other.

$Next \triangleq$

As usual, we call the next-state formula $Next$.

$Next \triangleq$

– Fill a jug.

As usual, we call the next-state formula $Next$.

First we allow steps that fill a jug.

There are two jugs, so we have two possible kinds of steps.

$Next \triangleq \vee \boxed{FillSmall}$ Fill small jug.
 $\vee FillBig$

As usual, we call the next-state formula $Next$.

First we allow steps that fill a jug.

There are two jugs, so we have two possible kinds of steps.

Steps that fill the small jug.

$$Next \triangleq \bigvee FillSmall \\ \bigvee \boxed{FillBig} \quad \text{Fill big jug.}$$

As usual, we call the next-state formula *Next*.

First we allow steps that fill a jug.

There are two jugs, so we have two possible kinds of steps.

Steps that fill the small jug. **And steps that fill the big jug.**

$Next \triangleq \vee FillSmall$
 $\vee FillBig$

– Empty a jug.

Similarly for steps that empty a jug.

$Next \triangleq \vee FillSmall$
 $\vee FillBig$
 $\vee EmptySmall$ – Empty a jug.
 $\vee EmptyBig$

Similarly for steps that empty a jug.

Next \triangleq \vee *FillSmall*
 \vee *FillBig*
 \vee *EmptySmall*
 \vee *EmptyBig*

– Pour from one jug
into the other.

Similarly for steps that empty a jug.

And there are two kinds of steps that pour from one jug to the other.

Next \triangleq \vee *FillSmall*
 \vee *FillBig*
 \vee *EmptySmall*
 \vee *EmptyBig*
 \vee *SmallToBig* From small jug to big jug.
 \vee *BigToSmall*

Similarly for steps that empty a jug.

And there are two kinds of steps that pour from one jug to the other.

Steps that pour from the small jug to the big jug.

Next \triangleq \vee *FillSmall*
 \vee *FillBig*
 \vee *EmptySmall*
 \vee *EmptyBig*
 \vee *SmallToBig*
 \vee *BigToSmall* From big jug to small jug.

Similarly for steps that empty a jug.

And there are two kinds of steps that pour from one jug to the other.

Steps that pour from the small jug to the big jug.

And steps that pour from the big jug to the small jug.

Next \triangleq \vee *FillSmall*
 \vee *FillBig*
 \vee *EmptySmall*
 \vee *EmptyBig*
 \vee *SmallToBig*
 \vee *BigToSmall*

Next \triangleq \vee *FillSmall*
 \vee *FillBig*
 \vee *EmptySmall*
 \vee *EmptyBig*
 \vee *SmallToBig*
 \vee *BigToSmall*

Names must be defined before they are used.

In TLA⁺, names must be defined before they're used.

$Next \triangleq$ \vee $FillSmall$
 \vee $FillBig$
 \vee $EmptySmall$
 \vee $EmptyBig$
 \vee $SmallToBig$
 \vee $BigToSmall$

Names must be defined before they are used.

The definitions of these names must precede this definition of $Next$.

In TLA⁺, names must be defined before they're used.

The definitions of $FillSmall$, $FillBig$, etc. must precede this definition of $Next$.

$FillSmall \triangleq$

We now define $FillSmall$.

Most people first learning TLA⁺ would write this definition.

$FillSmall \triangleq small' = 3$

Most people would write this definition.

We now define *FillSmall*.

Most people first learning TLA⁺ would write this definition.

$$\text{FillSmall} \triangleq \text{small}' = 3$$

Most people would write this definition.

Stop the video now and figure out why it's wrong.

We now define *FillSmall*.

Most people first learning TLA⁺ would write this definition.

Stop the video now and figure out why it's wrong.

$FillSmall \triangleq \boxed{small' = 3}$

If you didn't figure it out, you're thinking of this as setting $small$ to 3.

We now define $FillSmall$.

Most people first learning TLA⁺ would write this definition.

Stop the video now and figure out why it's wrong.

If you didn't figure it out, it means that you're thinking of this as an assignment statement that sets $small$ to 3. It's not.

$FillSmall \triangleq \boxed{small' = 3}$

If you didn't figure it out, you're thinking of this as setting *small* to 3.

It's a formula that's true for some steps and false for others.

It's a formula that's true for some steps and false for others.

$$FillSmall \triangleq \boxed{small' = 3}$$

If you didn't figure it out, you're thinking of this as setting *small* to 3.

It's a formula that's true for some steps and false for others.

It's true for any step in which *small* = 3 in the second state.

It's a formula that's true for some steps and false for others.

It's true for any step in which the value of *small* in the second state is 3.

$$FillSmall \triangleq \boxed{small' = 3}$$

It's a formula that's true for some steps and false for others.

It's true for any step in which the value of *small* in the second state is 3.

$FillSmall \triangleq \boxed{small' = 3}$

$$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : 3 \end{bmatrix}$$

It's a formula that's true for some steps and false for others.

It's true for any step in which the value of *small* in the second state is 3.

It's true for this step that appeared in the behavior we constructed.

$$FillSmall \triangleq \boxed{small' = 3}$$

$$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : 3 \end{bmatrix}$$

$$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : \sqrt{7} \end{bmatrix}$$

It's also true for this step in which *big* equals the square root of 7 in the second state.

$FillSmall \triangleq \boxed{small' = 3}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : 3 \end{bmatrix}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : \sqrt{7} \end{bmatrix}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : "abc" \end{bmatrix}$

It's also true for this step in which *big* equals the square root of 7 in the second state.

And it's also true for this step in which *big* equals the string *abc* in the second state.

$FillSmall \triangleq \boxed{small' = 3}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : 3 \end{bmatrix}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : \sqrt{7} \end{bmatrix}$

$\begin{bmatrix} small : 0 \\ big : 3 \end{bmatrix} \rightarrow \begin{bmatrix} small : 3 \\ big : "abc" \end{bmatrix}$

Should be false
for these.

Of course, these two steps shouldn't be allowed, so $FillSmall$ should equal false for them.

$$\begin{aligned} FillSmall &\stackrel{\Delta}{=} \wedge small' = 3 \\ &\wedge big' = big \end{aligned}$$

This is the correct definition.

Of course, these two steps shouldn't be allowed, so *FillSmall* should equal false for them.

And the correct definition should require the value of *big* to be unchanged.

$$FillSmall \triangleq \begin{array}{l} \wedge small' = 3 \\ \wedge big' = big \end{array}$$

Most people think this shouldn't be needed.

When they first see TLA+, most computer engineers and computer scientists think that this part of the formula shouldn't be needed.

And that you shouldn't have to say what's left unchanged.

My years of experience writing specifications and a couple of thousand years of mathematics say

$$FillSmall \triangleq \begin{array}{l} \wedge small' = 3 \\ \wedge big' = big \end{array}$$

Most people think this shouldn't be needed.

That's a bad idea!

When they first see TLA+, most computer engineers and computer scientists think that this part of the formula shouldn't be needed.

And that you shouldn't have to say what's left unchanged.

My years of experience writing specifications and a couple of thousand years of mathematics say **that's a bad idea.**

$$FillSmall \triangleq \begin{array}{l} \wedge small' = 3 \\ \wedge big' = big \end{array}$$

Most people think this shouldn't be needed.

That's a bad idea! It's not math.

It would leave the simple, elegant realm of mathematics — and enter the more complicated world of programming languages.

$$\begin{aligned} \text{FillSmall} &\triangleq \wedge \text{small}' = 3 \\ &\wedge \text{big}' = \text{big} \end{aligned}$$

$$\begin{aligned} \text{FillBig} &\triangleq \wedge \text{big}' = 5 \\ &\wedge \text{small}' = \text{small} \end{aligned}$$

The definition of *FillBig* is similar.

The definition of *FillBig* is similar.

POURING BETWEEN JUGS

Pouring from one jug into another.

[slide 132]

SmallToBig \triangleq

We now define *SmallToBig*.

SmallToBig \triangleq

There are two cases:

We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.

We now define *SmallToBig*.

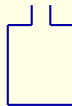
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define $SmallToBig$.

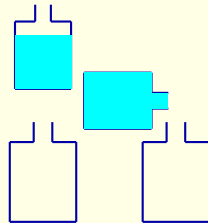
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define $SmallToBig$.

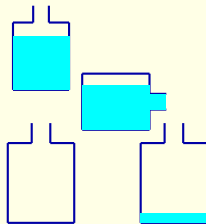
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug. Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define *SmallToBig*.

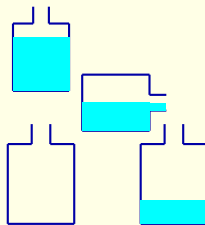
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug. Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define $SmallToBig$.

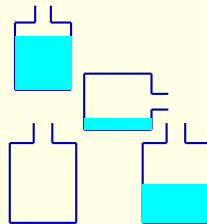
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$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define $SmallToBig$.

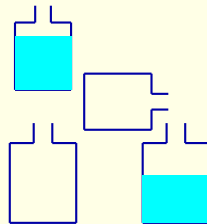
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define $SmallToBig$.

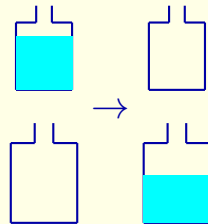
In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.



We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.

We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

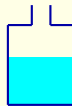
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define $SmallToBig$.

In the behavior we constructed, we saw that there are two cases:

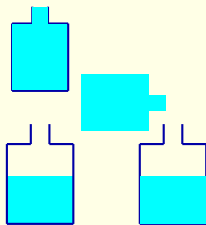
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

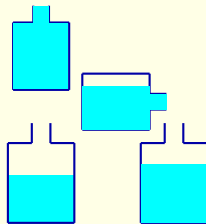
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

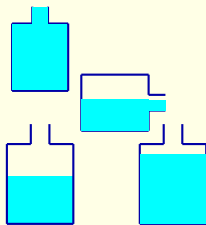
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

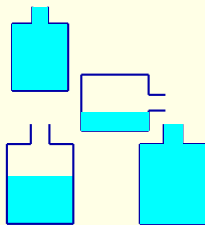
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define $SmallToBig$.

In the behavior we constructed, we saw that there are two cases:

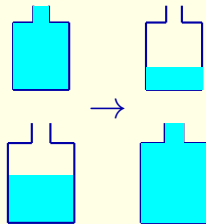
In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

$SmallToBig \triangleq$

There are two cases:

1. There **is** room in *big* for the water in *small*.
2. There **isn't** room in *big* for the water in *small*.



We now define *SmallToBig*.

In the behavior we constructed, we saw that there are two cases:

In case 1, there **is** room in the *big* jug for all the water in the *small* jug.
Here was that case.

In case 2, there **isn't** room in the *big* jug for all the water in the *small* jug.
Here was that case.

SmallToBig \triangleq

There are two cases:

1. There **is** room in *big*
for the water in *small* .
2. There **isn't** room in *big*
for the water in *small* .

$SmallToBig \triangleq$ IF $big + small \leq 5$
THEN 1. There is room: empty $small$.

ELSE 2. There isn't room: fill big .

There are two cases:

1. There **is** room in big
for the water in $small$.
2. There **isn't** room in big
for the water in $small$.

Which case it is depends on the total amount of water in the two jugs.

$SmallToBig \triangleq$ IF $big + small \leq 5$
THEN $\wedge big' = big + small$

ELSE 2. There isn't room: fill big .

Put water from $small$ into big

Which case it is depends on the total amount of water in the two jugs.

In case 1, we put all the water from the small jug into the big jug,

$SmallToBig \triangleq$ IF $big + small \leq 5$
THEN $\wedge big' = big + small$
 $\wedge small' = 0$
ELSE 2. There isn't room: fill big .

Put water from $small$ into big , emptying $small$.

Which case it is depends on the total amount of water in the two jugs.

In case 1, we put all the water from the small jug into the big jug, which empties the small jug.

$$\begin{aligned} \textit{SmallToBig} &\triangleq \text{IF } \textit{big} + \textit{small} \leq 5 \\ &\quad \text{THEN } \wedge \textit{big}' = \textit{big} + \textit{small} \\ &\quad \quad \wedge \textit{small}' = 0 \\ &\quad \text{ELSE } \wedge \textit{big}' = \\ &\quad \quad \wedge \textit{small}' = \end{aligned}$$

Which case it is depends on the total amount of water in the two jugs.

In case 1, we put all the water from the small jug into the big jug, which empties the small jug.

Case 2

[slide 154]

$$\begin{aligned} \textit{SmallToBig} &\triangleq \text{IF } \textit{big} + \textit{small} \leq 5 \\ &\quad \text{THEN } \wedge \textit{big}' = \textit{big} + \textit{small} \\ &\quad \quad \wedge \textit{small}' = 0 \\ &\quad \text{ELSE } \wedge \textit{big}' = ? \\ &\quad \quad \wedge \textit{small}' = ? \end{aligned}$$

Problem: Complete the definition of *SmallToBig*

Which case it is depends on the total amount of water in the two jugs.

In case 1, we put all the water from the small jug into the big jug, which empties the small jug.

Case 2 is left as a problem.

$$\begin{aligned} \text{SmallToBig} &\stackrel{\Delta}{=} \text{IF } big + small \leq 5 \\ &\quad \text{THEN } \wedge big' = big + small \\ &\quad \quad \wedge small' = 0 \\ &\quad \text{ELSE } \wedge big' = ? \\ &\quad \quad \wedge small' = ? \end{aligned}$$

Problem: Complete the definition of *SmallToBig* and write the definition of *BigToSmall*.

Which case it is depends on the total amount of water in the two jugs.

In case 1, we put all the water from the small jug into the big jug, which empties the small jug.

Case 2 is left as a problem. As is writing the definition of *BigToSmall*.

$SmallToBig \triangleq$ IF $big + small \leq 5$
THEN $\wedge big' = big + small$
 $\wedge small' = 0$
ELSE $\wedge big' = ?$
 $\wedge small' = ?$

Problem: Complete the definition of $SmallToBig$
and write the definition of $BigToSmall$.

Stop the video and solve it now.

Stop the video and solve it now, writing down your solution.

$SmallToBig \triangleq$ IF $big + small \leq 5$
THEN $\wedge big' = big + small$
 $\wedge small' = 0$
ELSE $\wedge big' = ?$
 $\wedge small' = ?$

Problem: Complete the definition of $SmallToBig$
and write the definition of $BigToSmall$.

Stop the video and solve it now.

You'll check your solution later.

Stop the video and solve it now, writing down your solution.

You'll check your solution later.

SAVING OUR HEROES

We'll now use TLC to save our heroes.

Open the Toolbox.

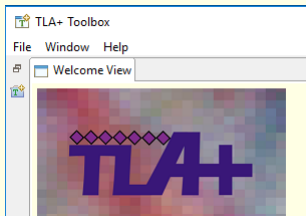
Open the Toolbox.

Open the Toolbox.

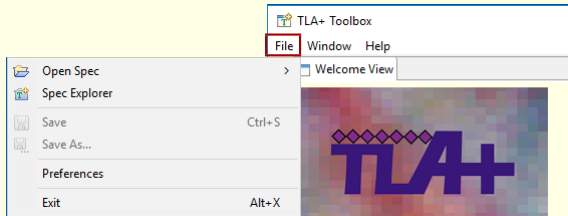
Open a new spec named *DieHard*.

Open the Toolbox.

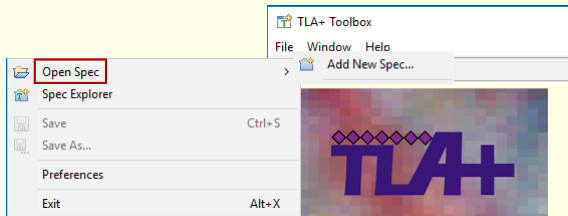
And then open a new spec named *DieHard*.



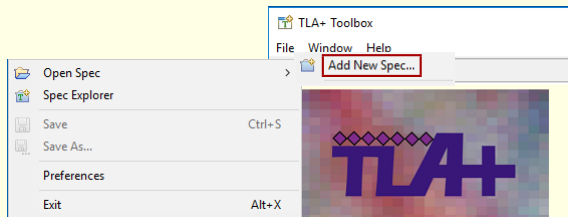
Remember you click on



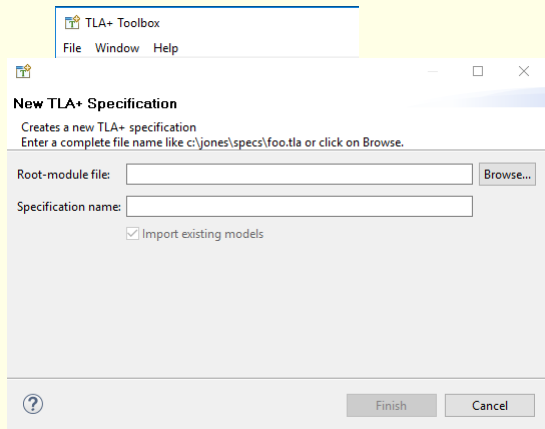
Remember you click on *File*. Then on



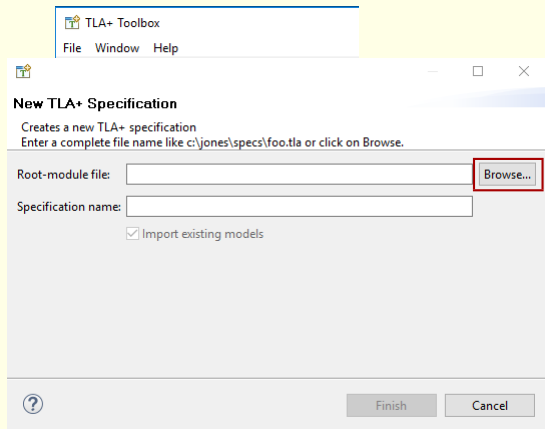
Remember you click on *File*. Then on *Open Spec*. Then on



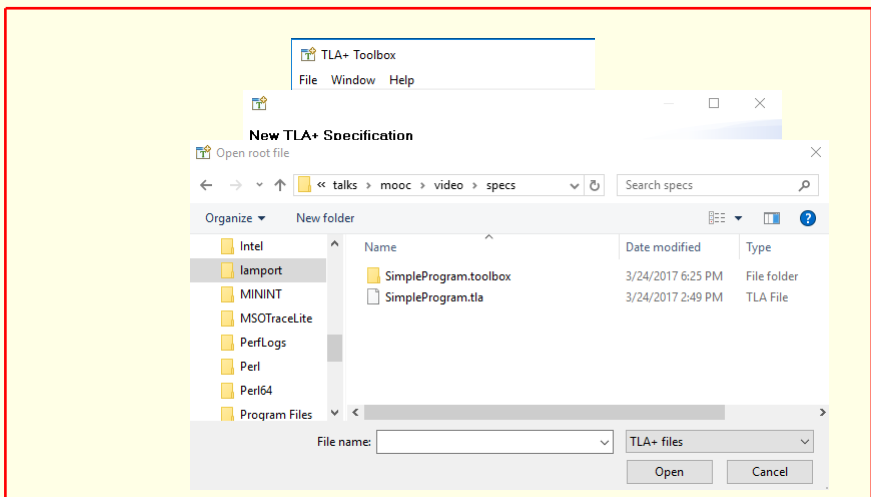
Remember you click on *File*. Then on *Open Spec*. Then on ***Add New Spec***, which opens



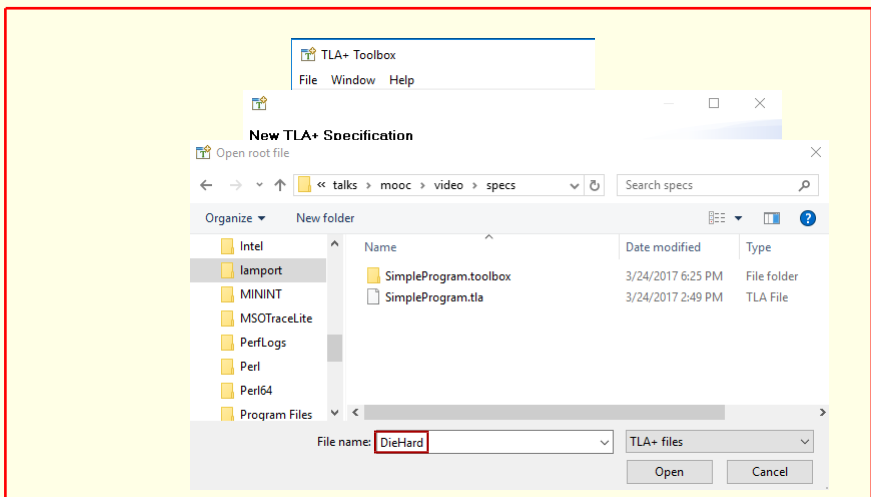
Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens **This window**. Then click on



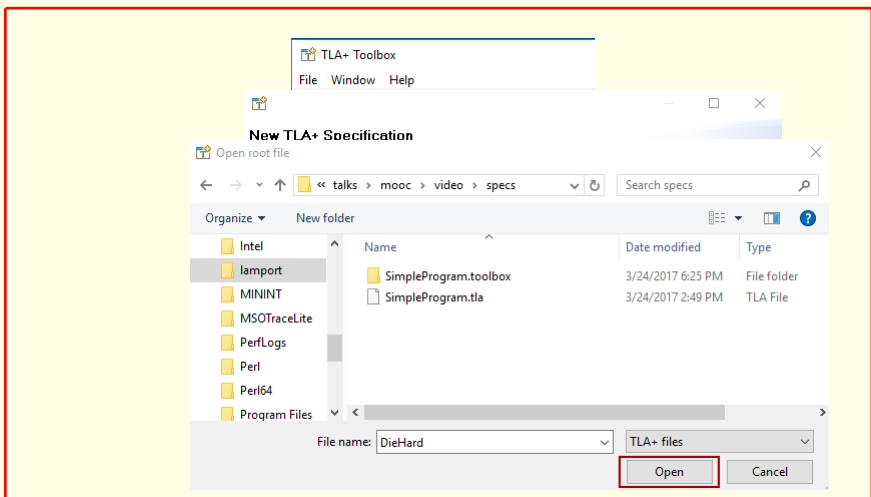
Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises



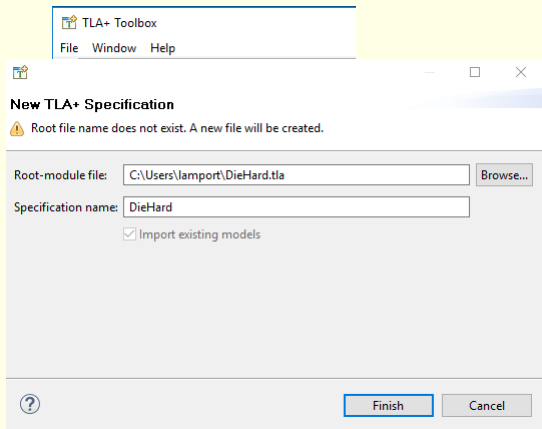
Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises a file browser window—probably on the folder in which you put the *SimpleProgram* spec. Select any folder and enter



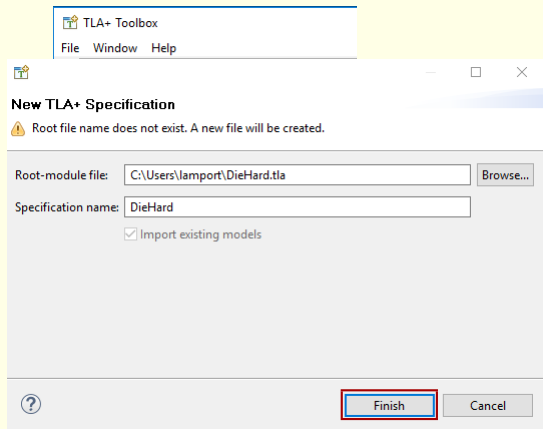
Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises a file browser window—probably on the folder in which you put the *SimpleProgram* spec. Select any folder and enter the file name *DieHard*. Then click on



Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises a file browser window—probably on the folder in which you put the *SimpleProgram* spec. Select any folder and enter the file name *DieHard*. Then click on *Open*



Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises a file browser window—probably on the folder in which you put the *SimpleProgram* spec. Select any folder and enter the file name *DieHard*. Then click on *Open* and then on



Remember you click on *File*. Then on *Open Spec*. Then on *Add New Spec*, which opens This window. Then click on *Browse* which raises a file browser window—probably on the folder in which you put the *SimpleProgram* spec. Select any folder and enter the file name *DieHard*. Then click on *Open* and then on *Finish* which opens

```
1 ----- MODULE DieHard -----  
2  
3  
4 =====
```

An empty spec named diehard.

```
1 ----- MODULE DieHard -----  
2  
3  
4 =====
```

Stop the video and copy the body of the spec.

An empty spec named diehard.

Stop the video now and copy the body of the specification that we just wrote.

```
1 ----- MODULE DieHard -----  
2   
3   
4 -----
```

Stop the video and copy the body of the spec.

Paste it here.

An empty spec named diehard.

Stop the video now and copy the body of the specification that we just wrote.

and then paste the text in the module here.

```
1 ----- MODULE DieHard -----
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7           /\ big   \in 0..5
8
9 Init == /\ big   = 0
10         /\ small = 0
11
12 FillSmall == /\ small' = 3
13              /\ big'   = big
14
15 FillBig == /\ big'   = 5
16             /\ small' = small
17
18 EmptySmall == /\ small' = 0
19               /\ big'   = big
20
21 EmptyBig == /\ big'   = 0
22             /\ small' = small
```

The module contains the complete definitions of *SmallToBig* and *BigToSmall*.

But don't look at them yet.

And here's what you should see.


```
1 ----- MODULE DieHard -----
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7          /\ big   \in 0..5
8
9 Init == /\ big   = 0
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12 FillSmall == /\ small' = 3
13             /\ big'   = big
14
15 FillBig == /\ big'   = 5
16            /\ small' = small
17
18 EmptySmall == /\ small' = 0
19              /\ big'   = big
20
21 EmptyBig == /\ big'   = 0
22            /\ small' = small
```

The module contains the complete definitions of *SmallToBig* and *BigToSmall*.

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10         /\ small = 0
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12 FillSmall == /\ small' = 3
13              /\ big'   = big
14
15 FillBig == /\ big'   = 5
16            /\ small' = small
17
18 EmptySmall == /\ small' = 0
19              /\ big'   = big
20
21 EmptyBig == /\ big'   = 0
22            /\ small' = small
```

The module contains the complete definitions of *SmallToBig* and *BigToSmall*.

But don't look at them yet.

And here's what you should see.

The module contains the complete definitions of *SmallToBig* and *BigToSmall* But don't look at them until after we've saved our heroes.

```
1 ----- MODULE DieHard -----
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7           /\ big   \in 0..5
8
9 Init == /\ big   = 0
10         /\ small = 0
11
12 FillSmall == /\ small' = 3
13              /\ big'   = big
14
15 FillBig == /\ big'   = 5
16            /\ small' = small
17
18 EmptySmall == /\ small' = 0
19              /\ big'   = big
20
21 EmptyBig == /\ big'   = 0
22            /\ small' = small
```

Save the module.

(Type ct1+S.)

And here's what you should see.

The module contains the complete definitions of *SmallToBig* and *BigToSmall*. But don't look at them until after we've saved our heroes.

And first, you have to save the module – which you can do

```
1 ----- MODULE DieHard -----
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7           /\ big   \in 0..5
8
9 Init == /\ big   = 0
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12 FillSmall == /\ small' = 3
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14
15 FillBig == /\ big'   = 5
16            /\ small' = small
17
18 EmptySmall == /\ small' = 0
19              /\ big'   = big
20
21 EmptyBig == /\ big'   = 0
22            /\ small' = small
```

Save the module.
(Type Ctl+S.)

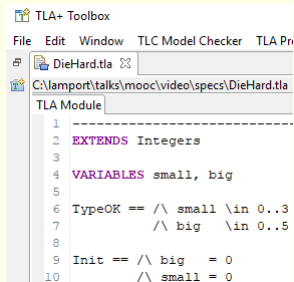
And here's what you should see.

The module contains the complete definitions of *SmallToBig* and *BigToSmall*. But don't look at them until after we've saved our heroes.

And first, you have to save the module – which you can do by typing Control S.

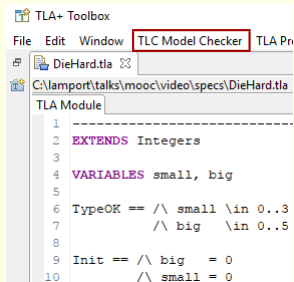
[slide 180]

To run TLC, we create a model.



```
TLA+ Toolbox
File Edit Window TLC Model Checker TLA Pr
DieHard.tla
C:\lampport\talks\moooc\video\specs\DieHard.tla
TLA Module
1
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7           /\ big \in 0..5
8
9 Init == /\ big = 0
10        /\ small = 0
```

To run TLC, we create a model by



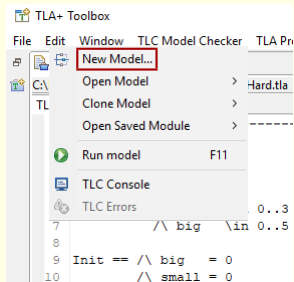
```
TLA+ Toolbox
File Edit Window TLC Model Checker TLA Pr
DieHard.tla
C:\lampport\talks\moooc\video\specs\DieHard.tla
TLA Module
1
2 EXTENDS Integers
3
4 VARIABLES small, big
5
6 TypeOK == /\ small \in 0..3
7           /\ big \in 0..5
8
9 Init == /\ big = 0
10        /\ small = 0
```

To run TLC, we create a model.

To run TLC, we create a model by

Clicking on the TLC Model Checker menu

To run TLC, we create a model.

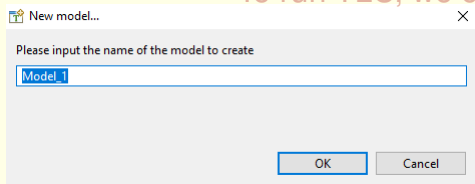


To run TLC, we create a model by

Clicking on the TLC Model Checker menu

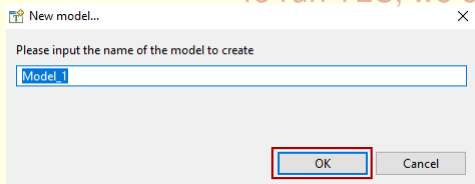
Selecting *New Model*

To run TLC, we create a model.



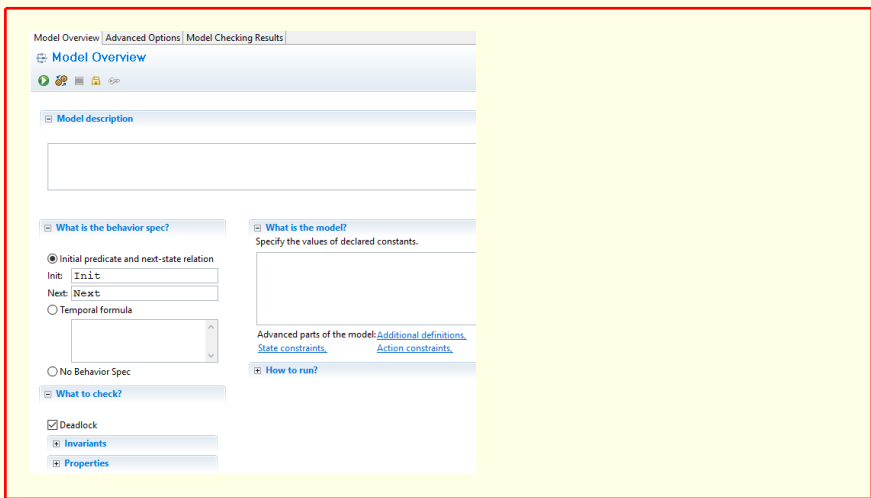
To run TLC, we create a model by
Clicking on the TLC Model Checker menu
Selecting *New Model*
Entering a model name and

To run TLC, we create a model.



To run TLC, we create a model by
Clicking on the TLC Model Checker menu
Selecting *New Model*
Entering a model name and
Clicking OK.

[slide 185]



This raises the Model Overview page

Model Overview | Advanced Options | Model Checking Results

Model Overview

Model description

What is the behavior spec?

Initial predicate and next-state relation

Init:

Next:

Temporal formula

No Behavior Spec

What to check?

Deadlock

Invariants

Properties

What is the model?

Specify the values of declared constants.

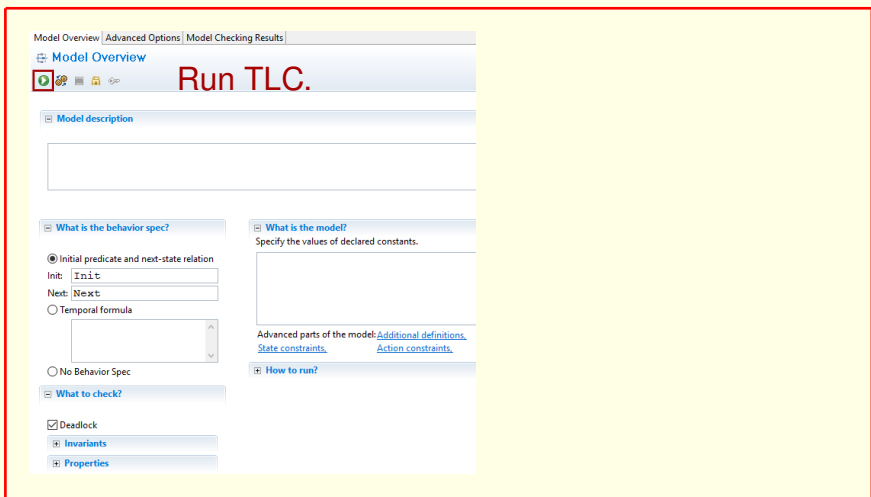
Advanced parts of the model: [Additional definitions](#), [State constraints](#), [Action constraints](#).

How to run?

Initial and next-state formulas

This raises the Model Overview page

Where the Toolbox has filled in the initial formula and the next-state formula.



This raises the Model Overview page

Where the Toolbox has filled in the initial formula and the next-state formula.

Let's now run TLC by clicking on this button.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
 End time: Thu Apr 20 11:54:49 PDT 2017
 Last checkpoint time:
 Current status: Not running
 Errors detected: No errors
 Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression: Value:

TLC quickly finishes, displaying the

Model Overview | Advanced Options | **Model Checking Results**

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

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State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

TLC quickly finishes, displaying the *Model Checking Results* page

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression: Value:

TLC reports no errors.

TLC quickly finishes, displaying the *Model Checking Results* page which reports that it found no errors. We didn't ask TLC to check anything, so

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running

Errors detected: No errors

Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph)

Time	Diameter	States Found	Distinct States	Queue Size	Coverage at 201
2017-04-20 11:54:49	9	97	16	0	Module DieHard DieHard DieHard DieHard DieHard

Evaluate Constant Expression

Expression: Value:

TLC reports no errors.

This means it could run the spec.

TLC quickly finishes, displaying the *Model Checking Results* page which reports that it found no errors. We didn't ask TLC to check anything, so this just means that the spec is one that it could execute.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression: Value:

TLC found 16
reachable states.

TLC quickly finishes, displaying the *Model Checking Results* page which reports that it found no errors. We didn't ask TLC to check anything, so this just means that the spec is one that it could execute.

TLC also reports that it found 16 reachable states

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression: Value:

TLC found 16
reachable states.

(States occurring in some
behavior allowed by spec.)

TLC quickly finishes, displaying the *Model Checking Results* page which reports that it found no errors. We didn't ask TLC to check anything, so this just means that the spec is one that it could execute.

TLC also reports that it found 16 reachable states which are states that occur in some behavior allowed by the spec.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
End time: Thu Apr 20 11:54:49 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

Let's check type correctness

Let's now check type correctness – which means

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
 End time: Thu Apr 20 11:54:49 PDT 2017
 Last checkpoint time:
 Current status: Not running
 Errors detected: No errors
 Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

Let's check type correctness
 (every reachable state
 satisfies *TypeOK*).

Let's now check type correctness – which means that every reachable state
 satisfies formula *TypeOK*.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
 End time: Thu Apr 20 11:54:49 PDT 2017
 Last checkpoint time:
 Current status: Not running
 Errors detected: No errors
 Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph)

Time	Diameter	States Found	Distinct States	Queue Size	Coverage at 201
2017-04-20 11:54:49	9	97	16	0	Module DieHard DieHard DieHard DieHard

Evaluate Constant Expression

Expression:

Value:

Let's check type correctness
 (every reachable state
 satisfies *TypeOK*).

Let's now check type correctness – which means that every reachable state satisfies formula *TypeOK*.

Remember that this formula asserts that each variable has a reasonable value.

Model Overview | **Advanced Options** | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 11:54:48 PDT 2017
 End time: Thu Apr 20 11:54:49 PDT 2017
 Last checkpoint time:
 Current status: Not running
 Errors detected: No errors
 Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 11:54:49	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

Let's now check type correctness – which means that every reachable state satisfies formula *TypeOK*.

Remember that this formula asserts that each variable has a reasonable value.

To do this, we must go back to the *Model Overview* page.

Model Overview



Model description

What is the behavior spec?

Initial predicate and next-state relation

Init:

Next:

Temporal formula

No Behavior Spec

What to check?

Deadlock

Invariants

Properties

What is the model?

Specify the values of declared constants.

Advanced parts of the model: [Additional definitions](#),
[State constraints](#), [Action constraints](#).

How to run?

Model Overview | Advanced Options | Model Checking Results

Model Overview

Model description

What is the behavior spec?

Initial predicate and next-state relation

Init:

Next:

Temporal formula

No Behavior Spec

What to check?

Deadlock

Invariants

Properties

What is the model?

Specify the values of declared constants.

Advanced parts of the model: [Additional definitions.](#) [State constraints.](#) [Action constraints.](#)

How to run?

A formula true in every reachable state is called an *invariant*.

A formula that is true in every reachable state is called an *invariant*. To have TLC check an invariant

Model Overview | Advanced Options | Model Checking Results

Model Overview

Model description

What is the behavior spec?

Initial predicate and next-state relation

Init:

Next:

Temporal formula

No Behavior Spec

What is the model?

Specify the values of declared constants.

Advanced parts of the model: [Additional definitions](#), [State constraints](#), [Action constraints](#).

How to run?

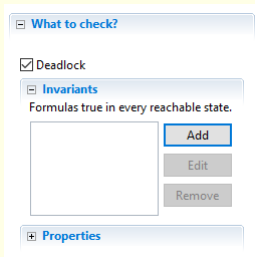
What to check?

Deadlock

Invariants

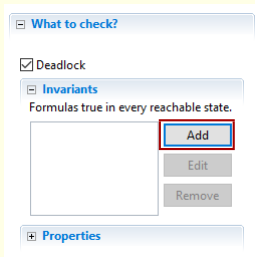
Properties

A formula that is true in every reachable state is called an *invariant*. To have TLC check an invariant Open the *Invariants* section of the model overview page.



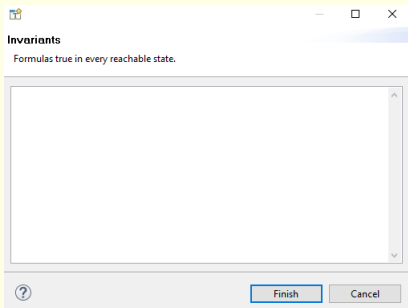
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



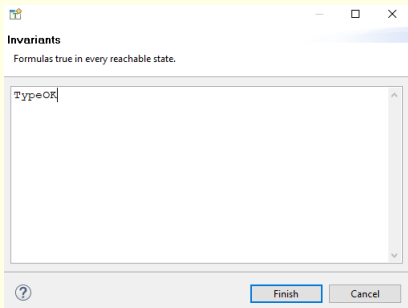
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



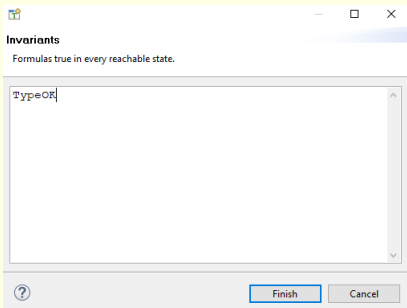
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



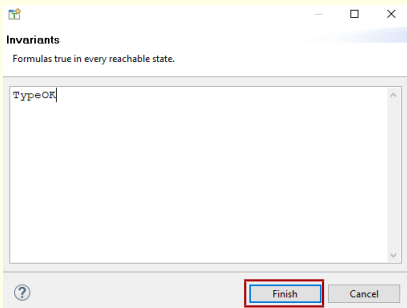
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



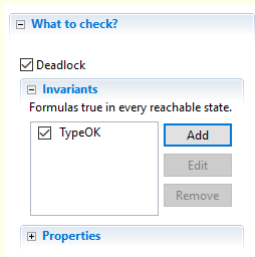
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



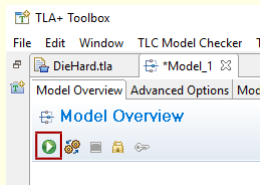
Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



Click on *Add*.

Enter *TypeOK*. And click on *Finish*.



Run TLC.





Click on *Add*.

Enter *TypeOK*. And click on *Finish*.

And run TLC on the model again.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 15:59:31 PDT 2017
 End time: Thu Apr 20 15:59:32 PDT 2017
 Last checkpoint time:
 Current status: Not running
 Errors detected: No errors
 Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-20 15:59:32	9	97	16	0	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

The *Model Checking Results* page

Model Overview | Advanced Options | Model Checking Results

Model Checking Results

General

Start time: Thu Apr 20 15:59:31 PDT 2017
End time: Thu Apr 20 15:59:32 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: No errors
Fingerprint collision probability: calculated: 7.0E-17, observed: 3.7E-18

Statistics

State space progress (click column header for graph)

Time	Diameter	States Found	Distinct States	Queue Size	Coverage at 201
2017-04-20 15:59:32	9	97	16	0	Module DieHard DieHard DieHard DieHard

Evaluate Constant Expression

Expression:

Value:

TLC reports no errors.

The *Model Checking Results* page

again shows that TLC found no errors.

Saving Our Heroes

Now we're ready to save our heroes.

Saving Our Heroes

The four gallons must be in the big jug.

Now we're ready to save our heroes.

The four gallons of water our heroes need must be in the big jug.

Saving Our Heroes

The four gallons must be in the big jug.

We let TLC check if $big \neq 4$ is an invariant.

Now we're ready to save our heroes.

The four gallons of water our heroes need must be in the big jug.

We let TLC check if $big \text{ not equal to } 4$ is an invariant.

Saving Our Heroes

The four gallons must be in the big jug.

We let TLC check if $big \neq 4$ is an invariant.

If it isn't, TLC will show us a behavior ending in a state with $big \neq 4$ false.

If it isn't, TLC will show us a behavior ending in a state with $big \neq 4$ false – a behavior that tells our heroes what they have to do to put 4 gallons in the big jug.

Saving Our Heroes

The four gallons must be in the big jug.

We let TLC check if $big \neq 4$ is an invariant.

\neq is written in ASCII as `/=` or `#`

If it isn't, TLC will show us a behavior ending in a state with $big \neq 4$ false – a behavior that tells our heroes what they have to do to put 4 gallons in the big jug.

In TLA+, *not equal* is written in ASCII as either forward slash equal-sign or sharp (also called *pound sign*).

Saving Our Heroes

The four gallons must be in the big jug.

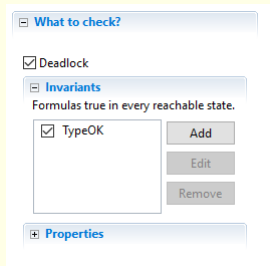
We let TLC check if $big \neq 4$ is an invariant.

We add this invariant to the model.

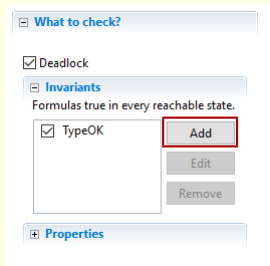
If it isn't, TLC will show us a behavior ending in a state with $big \neq 4$ false – a behavior that tells our heroes what they have to do to put 4 gallons in the big jug.

In TLA+, *not equal* is written in ASCII as either forward slash equal-sign or sharp (also called *pound sign*).

We now add this invariant to the model.

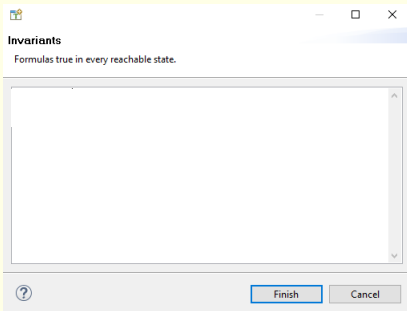


In the Invariants section of the model overview page



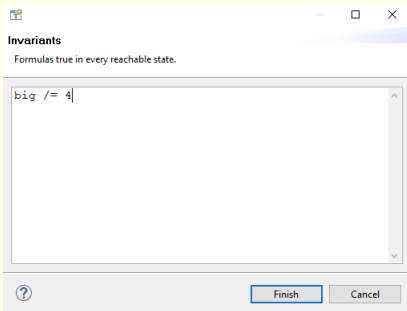
In the Invariants section of the model overview page

We add



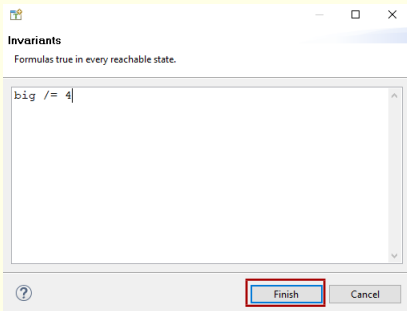
In the Invariants section of the model overview page

We add another invariant



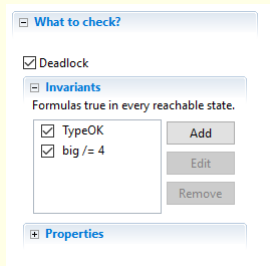
In the Invariants section of the model overview page

We add another invariant **big not equal to 4**



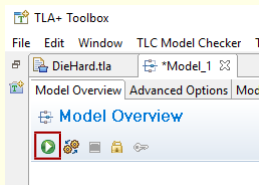
In the Invariants section of the model overview page

We add another invariant **big not equal to 4**



In the Invariants section of the model overview page

We add another invariant **big not equal to 4**



Run TLC.

In the Invariants section of the model overview page

We add another invariant big not equal to 4

And we run TLC.

Model Overview | Advanced Options | Model Checking Results

Model Checking Results 1 warning detected

General

Start time: Fri Apr 21 11:30:52 PDT 2017
End time: Fri Apr 21 11:30:52 PDT 2017
Last checkpoint time:
Current status: Not running
Errors detected: **1 Error**
Fingerprint collision probability:

Statistics

State space progress (click column header for graph) Coverage at 201

Time	Diameter	States Found	Distinct States	Queue Size	Module
2017-04-21 11:30:52	7	73	14	1	DieHard
					DieHard
					DieHard
					DieHard
					DieHard

Evaluate Constant Expression

Expression:

Value:

This time TLC reports an error.

The screenshot shows a window titled "TLC Errors" for "Model_1". The main text reads "Invariant big \neq 4 is violated." Below this is an "Error-Trace Exploration" section with an "Error-Trace" table. The table lists a sequence of states and transitions, with the final state highlighted in red to indicate the error.

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

This time TLC reports an error.

And the Toolbox opens this error window.

The screenshot shows the TLC Errors window for a model named 'Model_1'. The error message is 'Invariant big /= 4 is violated'. Below this, the 'Error-Trace Exploration' section shows an 'Error-Trace' table with 7 states. Each state is a snapshot of the 'big' and 'small' variables. The invariant is violated in the 2nd, 3rd, 4th, 5th, 6th, and 7th states.

Name	Value
▼ ▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

This time TLC reports an error.

And the Toolbox opens this error window.

Which tells us that the invariant was violated

The screenshot shows the TLC Errors window for a model named 'Model_1'. The message states: 'Invariant big /= 4 is violated.' Below this, the 'Error-Trace Exploration' section displays a table of the error trace.

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

This time TLC reports an error.

And the Toolbox opens this error window.

Which tells us that the invariant was violated And displays this error trace.

TLC Errors

Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▼ ▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

A behavior ending in state

$$\begin{bmatrix} big & : & 4 \\ small & : & 3 \end{bmatrix}$$

This time TLC reports an error.

And the Toolbox opens this error window.

Which tells us that the invariant was violated And displays this error trace.

The error trace is a behavior satisfying the spec ending in this state

TLC Errors

Model_1

Invariant: `big /= 4` is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▼ ▲ <Initial predicate>	State (num = 1)
big	0
small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
big	5
small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
big	2
small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
big	2
small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
big	0
small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
big	5
small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
big	4
small	3

A behavior ending in state

$$\begin{bmatrix} big & : & 4 \\ small & : & 3 \end{bmatrix}$$

with the invariant false.

a state in which the invariant equals false.

TLC Errors Model_1

Invariant big \neq 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 0 \\ small & : & 2 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 0 \\ small & : & 2 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 2 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

The complete behavior:

$$\begin{aligned}
 & \begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix} \\
 & \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 0 \\ small & : & 2 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 2 \end{bmatrix} \\
 & \rightarrow \begin{bmatrix} big & : & 4 \\ small & : & 3 \end{bmatrix}
 \end{aligned}$$

a state in which the invariant equals false.

It shows this complete behavior

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
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▼ ▲ <Action line 15, col 12 to li	State (num = 2)
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▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

$$\begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 0 \\ small & : & 2 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 2 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} big & : & 4 \\ small & : & 3 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior
gallons of water in the big jug.

From this behavior, our heroes should be able to see how to get 4

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▼ ▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

$$\begin{aligned}
 & \begin{bmatrix} big & : & 0 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix} \\
 & \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 0 \\ small & : & 2 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 5 \\ small & : & 2 \end{bmatrix} \\
 & \rightarrow \begin{bmatrix} big & : & 4 \\ small & : & 3 \end{bmatrix}
 \end{aligned}$$

a state in which the invariant equals false.

It shows this complete behavior

From this behavior, our heroes should be able to see how to get 4

gallons of water in the big jug. But they might not be the brightest bulbs on the block, and they may need help figuring out how to get from one state to the next. The Toolbox provides that help.

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

$$\begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

a state in which the invariant equals false.

It shows this complete behavior

From this behavior, our heroes should be able to see how to get 4 gallons of water in the big jug. But they might not be the brightest bulbs on the block, and they may need help figuring out how to get from one state to the next. The Toolbox provides that help.

To see why this step is allowed by the spec

TLC Errors Model_1

Invariant big /= 4 is violated.

Error-Trace Exploration

Error-Trace

Name	Value
▲ <Initial predicate>	State (num = 1)
■ big	0
■ small	0
▼ ▲ <Action line 15, col 12 to li	State (num = 2)
■ big	5
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 3)
■ big	2
■ small	3
▼ ▲ <Action line 18, col 15 to li	State (num = 4)
■ big	2
■ small	0
▼ ▲ <Action line 30, col 15 to li	State (num = 5)
■ big	0
■ small	2
▼ ▲ <Action line 15, col 12 to li	State (num = 6)
■ big	5
■ small	2
▼ ▲ <Action line 30, col 15 to li	State (num = 7)
■ big	4
■ small	3

$$\begin{bmatrix} big & : & 5 \\ small & : & 0 \end{bmatrix} \rightarrow \begin{bmatrix} big & : & 2 \\ small & : & 3 \end{bmatrix}$$

Double-click here.

a state in which the invariant equals false.

It shows this complete behavior

From this behavior, our heroes should be able to see how to get 4 gallons of water in the big jug. But they might not be the brightest bulbs on the block, and they may need help figuring out how to get from one state to the next. The Toolbox provides that help.

To see why this step is allowed by the spec Double-click here to find the part of the next-state formula that allows this step

[slide 242]

```

27         ELSE /\ big' = 5
28             /\ small' = small - (5 - big)
29
30 BigToSmall == IF big + small =< 3
31               THEN /\ big' = 0
32                   /\ small' = big + small
33               ELSE /\ big' = small - (3 - big)
34                   /\ small' = 3
35
36 Next == \/ FillSmall
37         \/ FillBig

```

a state in which the invariant equals false.

It shows this complete behavior [From this behavior](#), our heroes should be able to see how to get 4 gallons of water in the big jug. But they might not be the brightest bulbs on the block, and they may need help figuring out how to get from one state to the next. The Toolbox provides that help.

To see why this step is allowed by the spec [Double-click here](#) to find the part of the next-state formula that allows this step **And even Hollywood actors should be able to figure out**

```

27         ELSE /\ big' = 5
28             /\ small' = small - (5 - big)
29
30 BigToSmall == IF big + small =< 3
31               THEN /\ big' = 0
32                   /\ small' = big + small
33               ELSE /\ big' = small - (3 - big)
34                   /\ small' = 3
35
36 Next == \/ FillSmall
37        \/ FillBig

```

a state in which the invariant equals false.

It shows this complete behavior [From this behavior](#), our heroes should be able to see how to get 4 gallons of water in the big jug. But they might not be the brightest bulbs on the block, and they may need help figuring out how to get from one state to the next. The Toolbox provides that help.

To see why this step is allowed by the spec [Double-click here](#) to find the part of the next-state formula that allows this step. And even Hollywood actors should be able to figure out **that they have to pour the big jug into the small jug.**

[slide 244]

SmallToBig **AND** *BigToSmall*

Formulas *SmallToBig* and *BigToSmall*.

```
SmallToBig ==
```

```
BigToSmall ==
```

Now that we've saved our heroes, let's take a look at the definitions of *SmallToBig* and *BigToSmall*.

```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ big'   = 5
                   /\ small' = small - (5 - big)
```

Now that we've saved our heroes, let's take a look at the definitions of *SmallToBig* and *BigToSmall*.

Let's start with *SmallToBig*.

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

\leq is typed as `=<`

Notice that *less than or equal* is represented in ASCII as equal-sign less-than.


```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                   /\ small' = 0
                ELSE /\ big'   = 5
                   /\ small' = small - (5 - big)
```

Notice that *less than or equal* is represented in ASCII as equal-sign less-than.

```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                    /\ small' = 0
                ELSE /\ big'   = 5
                    /\ small' = small - (5 - big)
```

big is filled.

Notice that *less than or equal* is represented in ASCII as equal-sign less-than.

Remember that this is the case in which the *big* jug is filled from the small one.

```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                    /\ small' = 0
                ELSE /\ big'   = 5
                    /\ small' = small - (5 - big)
```

amount poured

Notice that *less than or equal* is represented in ASCII as equal-sign less-than.

Remember that this is the case in which the *big* jug is filled from the small one.

The amount poured into the *big* jug is removed from the small jug.

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                    /\ small' = 0
               ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big*

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big* is set to 5 here.

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big* is set to 5 here.

That's because she thinks of this as two assignment statements. But you know that it's actually a formula that specifies allowed steps.

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

$$A \wedge B = B \wedge A$$

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big* is set to 5 here.

That's because she thinks of this as two assignment statements. But you know that it's actually a formula that specifies allowed steps. And that *and* is commutative, so

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                    /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                    /\ big' = 5
```

$$A \wedge B = B \wedge A$$

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big* is set to 5 here.

That's because she thinks of this as two assignment statements. But you know that it's actually a formula that specifies allowed steps. And that *and* is commutative, so Changing the order of the two sub-formulas makes no difference.


```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                    /\ small' = 0
                ELSE /\ small' = small - (5 - big)
                    /\ big'   = 5
```

Someone who hasn't seen TLA+ before would think this is wrong because this value of *big* is set to 5 here.

That's because she thinks of this as two assignment statements. But you know that it's actually a formula that specifies allowed steps. And that *and* is commutative, so Changing the order of the two sub-formulas makes no difference.

```
SmallToBig == IF big + small =< 5
               THEN /\ big'  = big + small
                   /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                   /\ big'   = 5

BigToSmall == IF big + small =< 3
               THEN /\ big'   = 0
                   /\ small' = big + small
               ELSE /\ big'   = small - (3 - big)
                   /\ small' = 3
```

You can look at the definition of *BigToSmall* in the module by yourself later.

```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                   /\ big'   = 5
```

```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                   /\ big'   = 5
```

WARNING!

Here's a warning about writing specs.

```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                   /\ big'   = 5
```

= is also commutative

Here's a warning about writing specs.

The equality operator is also commutative.

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                   /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                   /\ big' = 5
```

= is also commutative

SO $small' = 0$

Here's a warning about writing specs.

The equality operator is also commutative. *so small prime equals 0*

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                   /\ 0 = small'
               ELSE /\ small' = small - (5 - big)
                   /\ big' = 5
```

= is also commutative

so $\text{small}' = 0$ is equivalent to $0 = \text{small}'$.

Here's a warning about writing specs.

The equality operator is also commutative. *so $\text{small prime equals } 0$ is completely equivalent to $0 \text{ equals small prime}$.*

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                   /\ 0 = small'
               ELSE /\ small' = small - (5 - big)
                   /\ big' = 5
```

These two specs are equivalent.

Here's a warning about writing specs.

The equality operator is also commutative. *so $small\ prime\ equals\ 0$ is completely equivalent to $0\ equals\ small\ prime$.*

These two specs are equivalent.


```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                    /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                    /\ big' = 5
```

These two specs are equivalent.

Here's a warning about writing specs.

The equality operator is also commutative. *so $small\ prime\ equals\ 0$ is completely equivalent to $0\ equals\ small\ prime$.*

These two specs are equivalent.

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ 0 = small'
                ELSE /\ small' = small - (5 - big)
                    /\ big' = 5
```

These two specs are equivalent.

Here's a warning about writing specs.

The equality operator is also commutative. so *small prime equals 0* is completely equivalent to *0 equals small prime*.

These two specs are equivalent.

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                   /\ 0 = small'
               ELSE /\ small' = small - (5 - big)
                   /\ big' = 5
```

These two specs are equivalent.

The TLAPS proof system treats them the same.

And the TLAPS proof system treats them exactly the same.

```
SmallToBig == IF big + small =< 5
               THEN /\ big' = big + small
                    /\ small' = 0
               ELSE /\ small' = small - (5 - big)
                    /\ big' = 5
```

These two specs are equivalent.

But TLC handles only this one.

And the TLAPS proof system treats them exactly the same.

But TLC handles only this one.

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ 0 = small'
                ELSE /\ small' = small - (5 - big)
                    /\ big' = 5
```

These two specs are equivalent.

It reports an error on this one.

And the TLAPS proof system treats them exactly the same.

But TLC handles only this one.

It reports an error if you run it on this one.

There are many ways to write a correct specification.

There are many ways to write a correct specification.

There are many ways to write a correct specification.

TLC can almost always handle the ones most engineers naturally write.

There are many ways to write a correct specification.

TLC can almost always handle the ones most engineers naturally write.

There are many ways to write a correct specification.

TLC can almost always handle the ones most engineers naturally write.

Later, you'll learn what specs TLC can handle.

There are many ways to write a correct specification.

TLC can almost always handle the ones most engineers naturally write.

Later, you'll learn what specs TLC can handle.

For now, follow this simple rule:

For now, just follow this simple rule:

For now, follow this simple rule:

Use a primed variable v' only
in one of these two kinds of
formulas:

$$v' = \dots \quad \text{and} \quad v' \in \dots$$

For now, just follow this simple rule:

Use a primed variable v -prime only in one of these two kinds of formulas

For now, follow this simple rule:

Use a primed variable v' only
in one of these two kinds of
formulas:

$$v' = \boxed{\dots} \quad \text{and} \quad v' \in \boxed{\dots}$$

no primed variables

For now, just follow this simple rule:

Use a primed variable v -prime only in one of these two kinds of formulas
where dot-dot-dot is an expression not containing primes.

For now, follow this simple rule:

Use a primed variable v' only in one of these two kinds of formulas:

$$v' = \boxed{\dots} \quad \text{and} \quad v' \in \boxed{\dots}$$

no primed variables

We'll relax this rule later.

For now, just follow this simple rule:

Use a primed variable v -prime only in one of these two kinds of formulas where dot-dot-dot is an expression not containing primes.

We'll relax this rule later.

CHECKING YOUR DEFINITIONS

Let's now check your definitions of *SmallToBig* and *BigToSmall*.

Your definitions of *SmallToBig* and *BigToSmall* are probably not exactly like mine.

Your definitions are probably not exactly the same as mine.

Your definitions of *SmallToBig* and *BigToSmall*
are probably not exactly like mine.

But they may still be correct.

Your definitions are probably not exactly the same as mine.

But they may still be correct.

Your definitions of *SmallToBig* and *BigToSmall* are probably not exactly like mine.

But they may still be correct.

Math provides many ways of writing the same formula.

Your definitions are probably not exactly the same as mine.

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Let's check your definitions.

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Your definitions of *SmallToBig* and *BigToSmall* are probably not exactly like mine.

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Math provides many ways of writing the same formula.

Let's check your definitions.

But first, let's see how we find errors.

Let's check your definitions.

But first, let's see how we find errors.

Parsing Errors

When writing a spec, our first mistakes are found by the parser.

Parsing Errors

```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                   /\ small' = 0
                ELSE /\ big'   = 5
                   /\ small' = small - (5 - big)
```

Modify the spec.

When writing a spec, our first mistakes are found by the parser.

For example, in the Toolbox, modify the definition of *SmallToBig* by deleting

Parsing Errors

```
SmallToBig == IF big + small =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

Modify the spec.

When writing a spec, our first mistakes are found by the parser.

For example, in the Toolbox, modify the definition of *SmallToBig* by deleting this plus sign.

Parsing Errors

```
SmallToBig == IF big small =< 5
                THEN /\ big'   = big + small
                    /\ small' = 0
                ELSE /\ big'   = 5
                    /\ small' = small - (5 - big)
```

When writing a spec, our first mistakes are found by the parser.

For example, in the Toolbox, modify the definition of *SmallToBig* by deleting this plus sign.

Parsing Errors

```
SmallToBig == IF big small =< 5
                THEN /\ big'   = big + small
                   /\ small' = 0
                ELSE /\ big'   = 5
                   /\ small' = small - (5 - big)
```

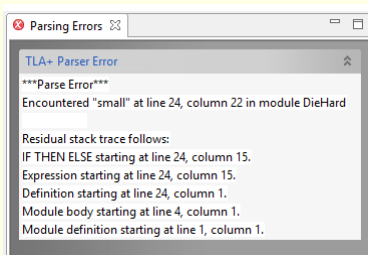
Save the spec.

When writing a spec, our first mistakes are found by the parser.

For example, in the Toolbox, modify the definition of *SmallToBig* by deleting this plus sign.

Now save the spec.

Parsing Errors



When writing a spec, our first mistakes are found by the parser.

For example, in the Toolbox, modify the definition of *SmallToBig* by deleting this plus sign.

Now save the spec.

The Toolbox runs the parser, which raises this error window.

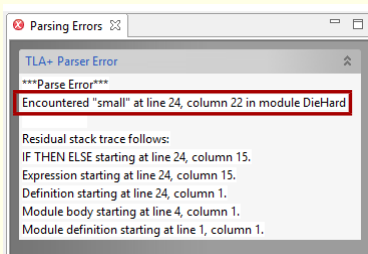
Parsing Errors

```
24 SmallToBig == IF big small =< 5
25                 THEN /\ big'   = big + small
26                 /\ small' = 0
27                 ELSE /\ big'   = 5
28                 /\ small' = small - (5 - big)
29
```

And it puts this error mark in the module editor.

Parsing Errors

Click
here.



And it puts this error mark in the module editor.

Clicking here in the error window

Parsing Errors

```
24 SmallToBig == IF big small =< 5
25           THEN /\ big' = big + small
26           /\ small' = 0
27           ELSE /\ big' = 5
28           /\ small' = small - (5 - big)
29
```

And it puts this error mark in the module editor.

Clicking here in the error window Highlights this part of the module and jumps to it.

Another common error found by parsing:

```
SmallToBig == IF big + small =< 5
                THEN /\ big'   = big + small
                    /\ small' = 0
                ELSE /\ big'   = 5
                    /\ small' = small - (5 - big)
```

Here's another common error found by the parser.

Another common error found by parsing:

```
SmallToBig == IF big + smell =< 5
                THEN /\ big' = big + small
                    /\ small' = 0
                ELSE /\ big' = 5
                    /\ small' = small - (5 - big)
```

An identifier not defined or declared.

Here's another common error found by the parser.

An identifier not yet defined or declared. This is usually a typo.

TLC “Execution Errors”

Errors TLC finds trying to “execute” the spec.

After there are no more parsing errors, TLC can often find errors while trying to *quote* execute the spec. (We’ll see in a later video how TLC does that.)

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After there are no more parsing errors, TLC can often find errors while trying to *quote* execute the spec. (We'll see in a later video how TLC does that.)

For example, change

TLC “Execution Errors”

Errors TLC finds trying to “execute” the spec.

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After there are no more parsing errors, TLC can often find errors while trying to *quote* execute the spec. (We’ll see in a later video how TLC does that.)

For example, change **this five** to **quote five**.

TLC “Execution Errors”

Errors TLC finds trying to “execute” the spec.

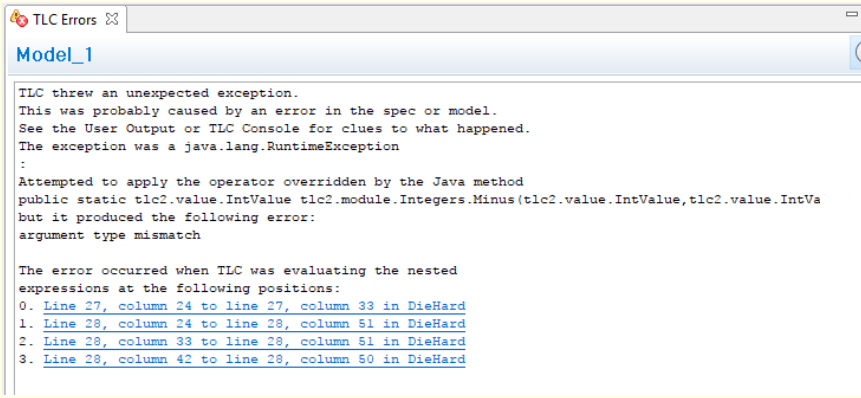
```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ big'   = 5
                   /\ small' = small - ("5" - big)
```

After there are no more parsing errors, TLC can often find errors while trying to *quote* execute the spec. (We'll see in a later video how TLC does that.)

For example, change this five to quote five. **And save the spec.**

Running TLC now produces

TLC “Execution Errors”



TLC Errors

Model_1

TLC threw an unexpected exception.
This was probably caused by an error in the spec or model.
See the User Output or TLC Console for clues to what happened.
The exception was a java.lang.RuntimeException
:
Attempted to apply the operator overridden by the Java method
public static tlc2.value.IntValue tlc2.module.Integers.Minus(tlc2.value.IntValue, tlc2.value.IntVa
but it produced the following error:
argument type mismatch

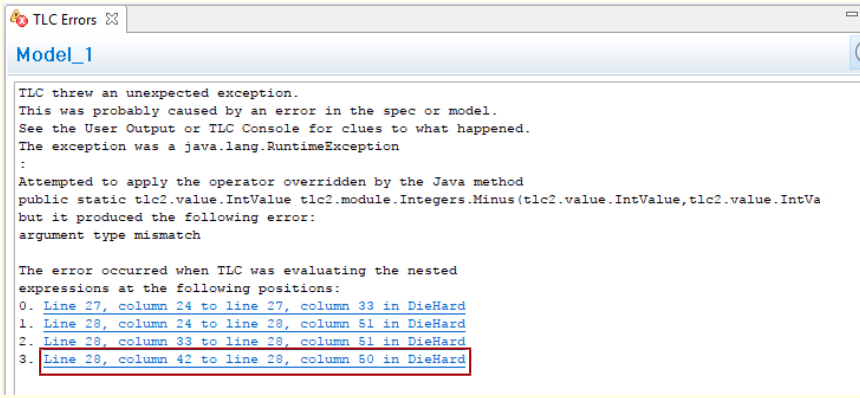
The error occurred when TLC was evaluating the nested
expressions at the following positions:

0. [Line 27, column 24 to line 27, column 33 in DieHard](#)
1. [Line 28, column 24 to line 28, column 51 in DieHard](#)
2. [Line 28, column 33 to line 28, column 51 in DieHard](#)
3. [Line 28, column 42 to line 28, column 50 in DieHard](#)

this error. You can read the complete error report later if you're curious.

For now, just

TLC “Execution Errors”



TLC Errors

Model_1

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For now, just click **here**, which selects and goes to

TLC “Execution Errors”

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SmallToBig == IF big + small =< 5  
              THEN /\ big'   = big + small  
                  /\ small' = 0  
              ELSE /\ big'   = 5  
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```



this error. You can read the complete error report later if you're curious.

For now, just click [here](#), which selects and goes to this part of the module.

Checking Your Definitions

Now, check your definitions of *SmallToBig* and *BigToSmall*.

Checking Your Definitions

```
SmallToBig == IF big + small =< 5
               THEN /\ big'   = big + small
                   /\ small' = 0
               ELSE /\ big'   = 5
                   /\ small' = small - (5 - big)

BigToSmall == IF big + small =< 3
               THEN /\ big'   = 0
                   /\ small' = big + small
               ELSE /\ big'   = small - (3 - big)
                   /\ small' = 3
```

Now, check your definitions of *SmallToBig* and *BigToSmall*.

First

Checking Your Definitions

```
(*  
SmallToBig == IF big + small =< 5  
               THEN /\ big'   = big + small  
                   /\ small' = 0  
               ELSE /\ big'   = 5  
                   /\ small' = small - (5 - big)  
  
BigToSmall == IF big + small =< 3  
               THEN /\ big'   = 0  
                   /\ small' = big + small  
               ELSE /\ big'   = small - (3 - big)  
                   /\ small' = 3  
  
*)
```

Comment out my definitions

Now, check your definitions of *SmallToBig* and *BigToSmall*.

First comment out my definitions by

Checking Your Definitions

(*)

```
SmallToBig == IF big + small =< 5
              THEN /\ big'   = big + small
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BigToSmall == IF big + small =< 3
              THEN /\ big'   = 0
                  /\ small' = big + small
              ELSE /\ big'   = small - (3 - big)
                  /\ small' = 3
```

(*)

Comment out my definitions

Now, check your definitions of *SmallToBig* and *BigToSmall*.

First comment out my definitions by adding these comment delimiters.

Checking Your Definitions

```
(*  
SmallToBig == IF big + small =< 5  
               THEN /\ big'   = big + small  
                   /\ small' = 0  
               ELSE /\ big'   = 5  
                   /\ small' = small - (5 - big)  
  
BigToSmall == IF big + small =< 3  
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                   /\ small' = big + small  
               ELSE /\ big'   = small - (3 - big)  
                   /\ small' = 3  
  
*)
```

Comment out my definitions and add your own.

Now, check your definitions of *SmallToBig* and *BigToSmall*.

First comment out my definitions by adding these comment delimiters.

And add your own definitions.

Save your definitions

Save your definitions

Save your definitions and correct any errors the parser finds.

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Save your definitions and correct any errors the parser finds.

Run TLC.

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Run TLC.

Your definitions are probably correct if TLC:

- Finds no “execution” errors.

Save your definitions and correct any errors the parser finds.

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Run TLC.

Your definitions are probably correct if TLC:

- Finds no “execution” errors.
- Finds no violation of the invariant *TypeOK* .

Save your definitions and correct any errors the parser finds.

Run TLC.

Your definitions are probably correct if TLC:

- Finds no “execution” errors.
- Finds no violation of the invariant *TypeOK* .

Save your definitions and correct any errors the parser finds.

Run TLC.

Your definitions are probably correct if TLC:

- Finds no “execution” errors.
- Finds no violation of the invariant *TypeOK* .
- Finds a violation of the alleged invariant *big* $\neq 4$.

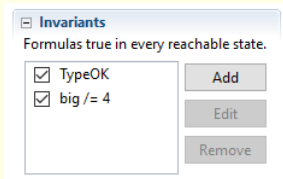
Save your definitions and correct any errors the parser finds.

Run TLC.

Your definitions are probably correct if TLC:

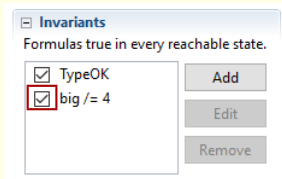
- Finds no “execution” errors.
- Finds no violation of the invariant *TypeOK* .
- And finds a violation of the alleged invariant *big* $\neq 4$.

To be sure, go here



To be sure, go to the *Invariants* section of the *Model Overview* page and

To be sure, go here

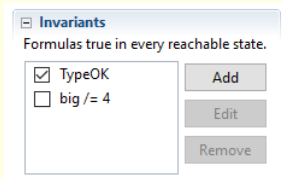


Uncheck this box

To be sure, go to the *Invariants* section of the *Model Overview* page and

Uncheck this box.

To be sure, go here

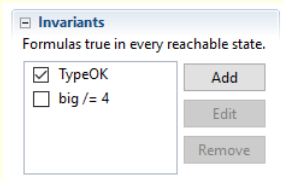


Uncheck this box so only *TypeOK* will be tested.

To be sure, go to the *Invariants* section of the *Model Overview* page and
Uncheck this box.

So only the *TypeOK* invariant will be tested by TLC.

To be sure, go here



Uncheck this box so only *TypeOK* will be tested.

Run TLC again.

To be sure, go to the *Invariants* section of the *Model Overview* page and

Uncheck this box.

So only the *TypeOK* invariant will be tested by TLC.

And run TLC again.

If TLC finds no error

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If TLC finds no error, try to find a different way to write the definitions.


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The best way to learn is by making mistakes.

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The best way to learn is by making mistakes.



Now that we've used TLC to save our heroes from certain death, it's time to leave the glamour of Hollywood for the more romantic subject of marriage and commitment. In the next lecture, we'll examine an algorithm that has been used for many years in weddings and database systems.

[slide 319]

End of Lecture 4

DIE HARD