

Optimization for Machine Learning

(Problems; Algorithms - C)

SUVRIT SRA

Massachusetts Institute of Technology

PKU Summer School on Data Science (July 2017)



Course materials

- <http://svr.it.de/teaching.html>
- Some references:
 - *Introductory lectures on convex optimization* – Nesterov
 - *Convex optimization* – Boyd & Vandenberghe
 - *Nonlinear programming* – Bertsekas
 - *Convex Analysis* – Rockafellar
 - *Fundamentals of convex analysis* – Urruty, Lemaréchal
 - *Lectures on modern convex optimization* – Nemirovski
 - *Optimization for Machine Learning* – Sra, Nowozin, Wright
 - *Theory of Convex Optimization for Machine Learning* – Bubeck
 - *NIPS 2016 Optimization Tutorial* – Bach, Sra
- Some related courses:
 - EE227A, Spring 2013, (Sra, UC Berkeley)
 - 10-801, Spring 2014 (Sra, CMU)
 - EE364a,b (Boyd, Stanford)
 - EE236b,c (Vandenberghe, UCLA)
- Venues: NIPS, ICML, UAI, AISTATS, SIOPT, Math. Prog.

Lecture Plan

- Introduction (3 lectures)
- Problems and algorithms (5 lectures)
- Non-convex optimization, perspectives (2 lectures)

Nonsmooth convergence rates

Theorem. (Nesterov.) Let $\mathcal{B} = \{x \mid \|x - x^0\|_2 \leq D\}$. Assume, $x^* \in \mathcal{B}$. There exists a convex function f in $C_L^0(\mathcal{B})$ (with $L > 0$), such that for $0 \leq k \leq n - 1$, the lower-bound

$$f(x^k) - f(x^*) \geq \frac{LD}{2(1+\sqrt{k+1})},$$

holds for **any algorithm** that generates x^k by linearly combining the previous iterates and subgradients.

Exercise: So design problems where we can do better!

Composite problems

Composite objectives

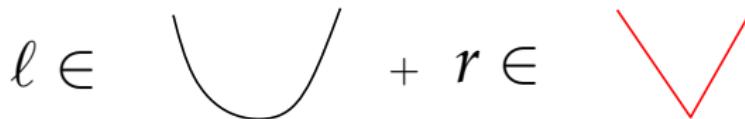
Frequently ML problems take the **regularized** form

$$\text{minimize } f(x) := \ell(x) + r(x)$$

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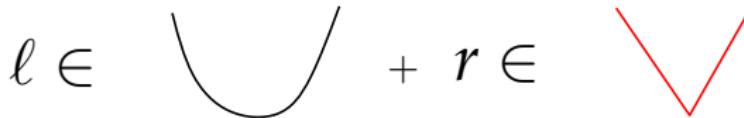
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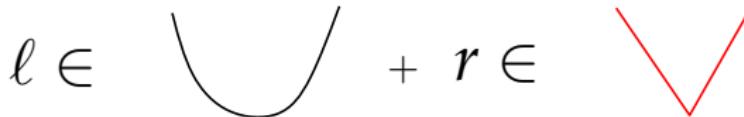
Example: $\ell(x) = \frac{1}{2} \|Ax - b\|^2$ and $r(x) = \lambda \|x\|_1$

Lasso, L1-LS, compressed sensing

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Lasso, L1-LS, compressed sensing

Example: $\ell(x)$: Logistic loss, and $r(x) = \lambda \|x\|_1$

L1-Logistic regression, sparse LR

Composite objective minimization

$$\text{minimize } f(x) := \ell(x) + r(x)$$

subgradient: $x^{k+1} = x^k - \alpha_k g^k, g^k \in \partial f(x^k)$

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but: f is *smooth* plus *nonsmooth*

we should **exploit:** smoothness of ℓ for better method!

Proximal Gradient Method

$$\min \quad f(x) \quad x \in \mathcal{X}$$

Projected (sub)gradient

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Why? If we can compute $\text{prox}_h(x)$ easily, prox-grad converges as fast gradient methods for smooth problems!

Proximity operator

Projection

$$P_{\mathcal{X}}(y) := \operatorname{argmin}_{x \in \mathbb{R}^n} \frac{1}{2} \|x - y\|_2^2 + \mathbb{1}_{\mathcal{X}}(x)$$

Proximity operator

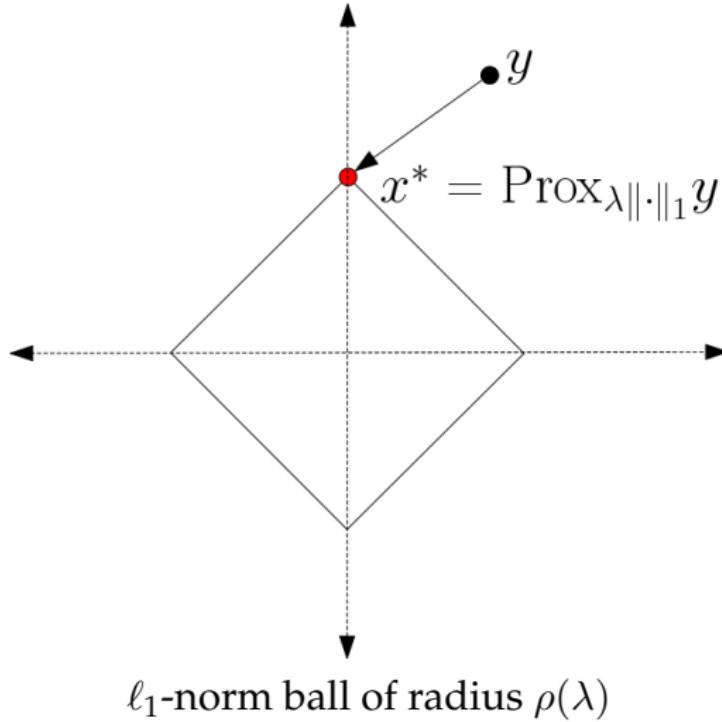
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$$P_{\mathcal{X}}(y) := \operatorname{argmin}_{x \in \mathbb{R}^n} \frac{1}{2} \|x - y\|_2^2 + \mathbb{1}_{\mathcal{X}}(x)$$

Proximity: Replace $\mathbb{1}_{\mathcal{X}}$ by a closed convex function

$$\operatorname{prox}_r(y) := \operatorname{argmin}_{x \in \mathbb{R}^n} \frac{1}{2} \|x - y\|_2^2 + r(x)$$

Proximity operator



Proximity operators

Exercise: Let $r(x) = \|x\|_1$. Solve $\text{prox}_{\lambda r}(y)$.

$$\min_{x \in \mathbb{R}^n} \quad \frac{1}{2} \|x - y\|_2^2 + \lambda \|x\|_1.$$

Hint 1: The above problem decomposes into n independent subproblems of the form

$$\min_{x \in \mathbb{R}} \quad \frac{1}{2} (x - y)^2 + \lambda |x|.$$

Hint 2: Consider the two cases: either $x = 0$ or $x \neq 0$

Exercise: Moreau decomposition $y = \text{prox}_h y + \text{prox}_{h^*} y$
(notice analogy to $V = S + S^\perp$ in linear algebra)

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Above fixed-point eqn suggests iteration

$$x_{k+1} = \text{prox}_{\alpha_k h}(x_k - \alpha_k \nabla f(x_k))$$

Convergence*

Proximal-gradient works, why?

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Gradient mapping: the “gradient-like object”

$$G_\alpha(x) = \frac{1}{\alpha}(x - P_{\alpha h}(x - \alpha \nabla f(x)))$$

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Gradient mapping: the “gradient-like object”

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- ▶ Our lemma shows: $G_\alpha(x) = 0$ if and only if x is optimal
- ▶ So G_α analogous to ∇f
- ▶ If x locally optimal, then $G_\alpha(x) = 0$ (nonconvex f)

Convergence analysis

Assumption: Lipschitz continuous gradient; denoted $f \in C_L^1$

$$\|\nabla f(x) - \nabla f(y)\|_2 \leq L\|x - y\|_2$$

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- ♣ Objective function has “bounded curvature”
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Lemma (Descent). Let $f \in C_L^1$. Then,

$$f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{L}{2} \|y - x\|_2^2$$

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For convex f , compare with

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle.$$

Descent lemma

Proof. Since $f \in C_L^1$, by Taylor's theorem, for the vector $z_t = x + t(y - x)$ we have

$$f(y) = f(x) + \int_0^1 \langle \nabla f(z_t), y - x \rangle dt.$$

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Add and subtract $\langle \nabla f(x), y - x \rangle$ on rhs we have

$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \int_0^1 \langle \nabla f(z_t) - \nabla f(x), y - x \rangle dt$$

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Bounds $f(y)$ around x with quadratic functions

Descent lemma – corollary

$$f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{L}{2} \|y - x\|_2^2$$

Let $y = x - \alpha G_\alpha(x)$, then

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Corollary. So if $0 \leq \alpha \leq 1/L$, we have

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Lemma Let $y = x - \alpha G_\alpha(x)$. Then, for any z we have

$$f(y) + h(y) \leq f(z) + h(z) + \langle G_\alpha(x), x - z \rangle - \frac{\alpha}{2} \|G_\alpha(x)\|_2^2.$$

Exercise: Prove! (hint: f, h are convex, $G_\alpha(x) - \nabla f(x) \in \partial h(y)$)

Convergence analysis

We've actually shown $x' = x - \alpha G_\alpha(x)$ is a descent method.
Write $\phi = f + h$; plug in $z = x$ to obtain

$$\phi(x') \leq \phi(x) - \frac{\alpha}{2} \|G_\alpha(x)\|_2^2.$$

Exercise: Why this inequality suffices to show convergence.

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Use $z = x^*$ in corollary to obtain progress in terms of iterates:

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$$\sum_{i=1}^{k+1} (\phi(x_i) - \phi^*) \leq \frac{L}{2} \sum_{i=1}^{k+1} [\|x_i - x^*\|_2^2 - \|x_{i+1} - x^*\|_2^2]$$

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Since $\phi(x_k)$ is a decreasing sequence, it follows that

$$\phi(x_{k+1}) - \phi^* \leq \frac{1}{k+1} \sum_{i=1}^{k+1} (\phi(x_i) - \phi^*) \leq \frac{L}{2(k+1)} \|x_1 - x^*\|_2^2.$$

This is the well-known $O(1/k)$ rate.

► But for C_L^1 convex functions, optimal rate is $O(1/k^2)$!

Accelerated Proximal Gradient

$$\min \phi(x) = f(x) + h(x)$$

Let $x^0 = y^0 \in \text{dom } h$. For $k \geq 1$:

$$x^k = \text{prox}_{\alpha_k h}(y^{k-1} - \alpha_k \nabla f(y^{k-1}))$$
$$y^k = x_k + \frac{k-1}{k+2}(x^k - x^{k-1}).$$

Framework due to: Nesterov (1983, 2004); also Beck, Teboulle (2009).

Simplified analysis: Tseng (2008).

- Uses extra “memory” for interpolation
- Same computational cost as ordinary prox-grad
- Convergence rate theoretically optimal

$$\phi(x^k) - \phi^* \leq \frac{2L}{(k+1)^2} \|x^0 - x^*\|_2^2.$$

Proximal splitting methods

$$\ell(x) + f(x) + h(x)$$

- ▶ Direct use of prox-grad not easy
- ▶ Requires computation of: $\text{prox}_{\lambda(f+h)}$ (i.e., $(I + \lambda(\partial f + \partial h))^{-1}$)

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Example:

$$\min \quad \frac{1}{2} \|x - y\|_2^2 + \underbrace{\lambda \|x\|_2}_{f(x)} + \underbrace{\mu \sum_{i=1}^{n-1} |x_{i+1} - x_i|}_{h(x)}.$$

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- But good feature: prox_f and prox_h separately easier
- Can we exploit that?

Proximal splitting – operator notation

- If $(I + \partial f + \partial h)^{-1}$ hard, but $(I + \partial f)^{-1}$ and $(I + \partial h)^{-1}$ “easy”

Proximal splitting – operator notation

- ▶ If $(I + \partial f + \partial h)^{-1}$ hard, but $(I + \partial f)^{-1}$ and $(I + \partial h)^{-1}$ “easy”
- ▶ Let us derive a fixed-point equation that “splits” the operators

Proximal splitting – operator notation

- If $(I + \partial f + \partial h)^{-1}$ hard, but $(I + \partial f)^{-1}$ and $(I + \partial h)^{-1}$ “easy”
- Let us derive a fixed-point equation that “splits” the operators

Assume we are solving

$$\min f(x) + h(x),$$

where both f and h are convex but potentially nondifferentiable.

Notice: We implicitly assumed: $\partial(f + h) = \partial f + \partial h$.

Proximal splitting

$$0 \in \partial f(x) + \partial h(x)$$

Proximal splitting

$$\begin{aligned} 0 &\in \partial f(x) + \partial h(x) \\ 2x &\in (I + \partial f)(x) + (I + \partial h)(x) \end{aligned}$$

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- ▶ Not a fixed-point equation yet
- ▶ We need one more idea

Douglas-Rachford splitting

Reflection operator

$$R_h(z) := 2 \operatorname{prox}_h(z) - z$$

Douglas-Rachford splitting

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Douglas-Rachford method

$$z \in (I + \partial h)(x), \quad x = \operatorname{prox}_h(z)$$

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$$z = 2\operatorname{prox}_f(R_h(z)) - R_h(z) = R_f(R_h(z))$$

Finally, z is on both sides of the eqn

Douglas-Rachford method

$$0 \in \partial f(x) + \partial h(x) \Leftrightarrow \begin{cases} x = \text{prox}_h(z) \\ z = R_f(R_h(z)) \end{cases}$$

DR method: given z_0 , iterate for $k \geq 0$

$$x_k = \text{prox}_h(z_k)$$

$$v_k = \text{prox}_f(2x_k - z_k)$$

$$z_{k+1} = z_k + \gamma_k(v_k - x_k)$$

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Theorem. If $f + h$ admits minimizers, and (γ_k) satisfy

$$\gamma_k \in [0, 2], \quad \sum_k \gamma_k(2 - \gamma_k) = \infty,$$

then the DR-iterates v_k and x_k converge to a minimizer.

Douglas-Rachford method

For $\gamma_k = 1$, we have

$$z_{k+1} = z_k + v_k - x_k$$

$$z_{k+1} = z_k + \text{prox}_f(2 \text{prox}_h(z_k) - z_k) - \text{prox}_h(z_k)$$

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Dropping superscripts, writing $P \equiv \text{prox}$, we have

$$z \leftarrow Tz$$

$$T = I + P_f(2P_h - I) - P_h$$

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Lemma DR can be written as: $z \leftarrow \frac{1}{2}(R_f R_h + I)z$, where R_f denotes the *reflection operator* $2P_f - I$ (similarly R_h).

Exercise: Prove this claim.

Proximal methods – cornucopia

- Douglas Rachford splitting
- ADMM (special case of DR on dual)
- Proximal-Dykstra
- Proximal methods for $f_1 + f_2 + \cdots + f_n$
- Peaceman-Rachford
- Proximal quasi-Newton, Newton
- Many other variation...

Best approximation problem

$$\min \quad \delta_A(x) + \delta_B(x) \quad \text{where } A \cap B = \emptyset.$$

Can we use DR?

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$$\min \quad \delta_A(x) + \delta_B(x) \quad \text{where } A \cap B = \emptyset.$$

Can we use DR?

Using a clever analysis of Bauschke & Combettes (2004), DR can still be applied! However, it generates diverging iterates which can be “projected back” to obtain a solution to

$$\min \quad \|a - b\|_2 \quad a \in A, b \in B.$$

See: Jegelka, Bach, Sra (NIPS 2013) for an example.

ADMM

Let us see separable objective with constraints

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$$\begin{aligned} \min \quad & f(x) + g(z) \\ \text{s.t.} \quad & Ax + Bz = c. \end{aligned}$$

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- ▶ The constraint prevents a trivial decoupling
- ▶ Introduce **augmented lagrangian** (AL)

$$L_\rho(x, z, y) := f(x) + g(z) + y^T(Ax + Bz - c) + \frac{\rho}{2} \|Ax + Bz - c\|_2^2$$

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