
Scalable and Robust Tensor Ring Decomposition for Large-scale Data (Supplementary Material)

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Before presenting the proof of the propositions, we need the following definition on Tensor Connect Product (TCP), which computes the tensor core merging.

Definition 1 (Tensor Connect Product (TCP) [Wang et al., 2017]). *Let $\mathcal{Z}_k \in \mathbb{R}^{r_k \times I_k \times r_{k+1}}$, $k = 1, \dots, N$ be N 3-order tensors. The tensor connect product (TCP) between \mathcal{Z}_k and \mathcal{Z}_{k+1} is defined as,*

$$\mathcal{Z}^{(k,k+1)} = \text{fold}(\mathbf{L}(\mathcal{Z}_k) \times \mathbf{R}(\mathcal{Z}_{k+1}))$$

where $\text{fold}(\mathbf{X})$ denotes the operation of reshaping the unfolding matrix \mathbf{X} back to tensor \mathcal{X} and

$$\mathbf{L}(\mathcal{X}) = (\mathbf{X}_{(3)})^T \in \mathbb{R}^{(r_k I_k) \times r_{k+1}}$$

$$\mathbf{R}(\mathcal{X}) = \mathbf{X}_{(1)} \in \mathbb{R}^{r_k \times (I_k r_{k+1})}.$$

First, we consider the computation of the Gram matrix using only two core tensors. According to the tensor core merging of two core tensors \mathcal{Z}_k and \mathcal{Z}_{k+1} , we establish the following lemma.

Lemma 1. *Let $\mathcal{Z}_k \in \mathbb{R}^{r_k \times I_k \times r_{k+1}}$, $k = 1, \dots, N$, be 3-rd order tensors. The Gram matrix of $\mathbf{Z}_{[2]}^{(k,k+1)}$ can be computed as*

$$\mathbf{G}_{\mathcal{Z}^{(k,k+1)}} = \mathbf{Z}_{[2]}^{(k,k+1),T} \mathbf{Z}_{[2]}^{(k,k+1)} = \Phi(\mathbf{Q}_k \mathbf{Q}_{k+1}) \quad (1)$$

where $\mathbf{Q}_k(:, i \times r_{k+1} + j) = \text{vec}\{(\mathcal{Z}_k(:, :, i)) \mathcal{Z}_k(:, :, j)^T\}$, with $\text{vec}\{\cdot\}$ denoting the vectorization operation, and $\Phi(\mathbf{X})$ is a reshape operation by which $\mathbf{X} \in \mathbb{R}^{m^2 \times n^2}$ is first divided into $m \times n$ blocks $\{\mathbf{X}_{ij}\}_{i,j=1}^{m,n} \in \mathbb{R}^{m \times n}$, then reshaped as

$$\Phi(\mathbf{X}) = [\text{vec}\{\mathbf{X}_{11}^T\} \text{vec}\{\mathbf{X}_{21}^T\} \dots \text{vec}\{\mathbf{X}_{mn}^T\}]^T.$$

1 PROOF OF LEMMA 1

Proof. From TCP in Definition 1, we can express the fiber-wise relation between mode-2 fibers of \mathcal{Z}_k , \mathcal{Z}_{k+1} and $\mathcal{Z}^{(k,k+1)}$ as

$$\mathcal{Z}^{(k,k+1)}(i, :, j) = \sum_{m=1}^{r_{k+1}} \mathcal{Z}_{k+1}(m, :, j) \otimes \mathcal{Z}_k(i, :, m), \quad (2)$$

with $i \in [1, r_k]$, $j \in [1, r_{k+2}]$ and \otimes denotes the kronecker product. Then, the (i, j) -th entry of the Gram matrix of $\mathbf{Z}_{[2]}^{(k,k+1)}$ can be computed as

$$\left[\mathbf{Z}_{[2]}^{(k,k+1),T} \mathbf{Z}_{[2]}^{(k,k+1)} \right]_{i,j} = \left(\mathcal{Z}^{(k,k+1)}(p_i, :, q_i) \right)^T \mathcal{Z}^{(k,k+1)}(p_j, :, q_j), \quad (3)$$

where $p_i = \lceil i/r_{k+1} \rceil, q_i = \text{mod}(i-1, r_{k+1}) + 1$. Substituting (2) in (3) and using the property

$$(\mathbf{A} \otimes \mathbf{B})(\mathbf{C} \otimes \mathbf{D}) = (\mathbf{AC}) \otimes (\mathbf{BD}), \quad (4)$$

we have that

$$\left[\mathbf{Z}_{[2]}^{(k,k+1),T} \mathbf{Z}_{[2]}^{(k,k+1)} \right]_{i,j} = \mathbf{v}_1^T \mathbf{v}_2, \quad (5)$$

where

$$\begin{aligned} \mathbf{v}_1 &= \text{vec}\{\mathcal{Z}_k(p_i, :, :)^T \mathcal{Z}_k(p_j, :, :)\} \\ \mathbf{v}_2 &= \text{vec}\{\mathcal{Z}_{k+1}(:, :, q_i) \mathcal{Z}_{k+1}(:, :, q_j)^T\}. \end{aligned}$$

Therefore, by defining $\mathbf{Q}_k(:, i \times r_{k+1} + j) = \text{vec}\{(\mathcal{Z}_k(:, :, i)) \mathcal{Z}_k(:, :, j)^T\}$, the Gram matrix of $\mathbf{Z}_{[2]}^{(k,k+1)}$ can be computed as

$$\mathbf{G}_{\mathcal{Z}^{(k,k+1)}} = \mathbf{Z}_{[2]}^{(k,k+1),T} \mathbf{Z}_{[2]}^{(k,k+1)} = \Phi(\mathbf{Q}_k \mathbf{Q}_{k+1}) \quad (6)$$

where $\Phi(\mathbf{X})$ is a reshape operation by which $\mathbf{X} \in \mathbb{R}^{m^2 \times n^2}$ is first divided into $m \times n$ blocks $\{\mathbf{X}_{ij}\}_{i,j}^{m,n} \in \mathbb{R}^{m \times n}$, then reshaped as

$$\Phi(\mathbf{X}) = [\text{vec}\{\mathbf{X}_{11}^T\} \text{vec}\{\mathbf{X}_{21}^T\} \dots \text{vec}\{\mathbf{X}_{mn}^T\}]^T. \quad \square$$

2 PROOF OF PROPOSITION 1

Following (1) in the proof of Lemma 1, for $\mathcal{Z}^{\leq c} \in \mathbb{R}^{r_1 \times \prod_{k=1}^c I_k \times r_{c+1}}$ which is a subchain obtained by merging c cores $\{\mathcal{Z}_k\}_{k=1}^c$, according to TCP, we can express the fiber-wise relation between mode-2 fibers of \mathcal{Z}_c and $\mathcal{Z}^{\leq c-1}$ as

$$\mathcal{Z}^{\leq c}(i, :, j) = \sum_{m=1}^{r_c} \mathcal{Z}_c(m, :, j) \otimes \mathcal{Z}^{\leq c-1}(i, :, m). \quad (7)$$

With the above recursion equation we have

$$\begin{aligned} &\mathcal{Z}^{\leq c}(i, :, j) \\ &= \sum_{m=1}^{r_c} \mathcal{Z}_c(m, :, j) \otimes \left(\sum_{m=1}^{r_{c-1}} \mathcal{Z}_{c-1}(m, :, j) \otimes \dots \left(\sum_{m=1}^{r_2} \mathcal{Z}_2(m, :, j) \otimes \mathcal{Z}_1(i, :, m) \right) \right) \end{aligned} \quad (8)$$

Again, using the property in (4), we can obtain that

$$\left[\mathbf{Z}_{[2]}^{\leq c,T} \mathbf{Z}_{[2]}^{\leq c} \right]_{i,j} = \mathbf{v}_1^T \mathbf{Q}_2^T \dots \mathbf{Q}_{c-1}^T \mathbf{v}_2, \quad (9)$$

where

$$\begin{aligned} \mathbf{v}_1 &= \text{vec}\{\mathcal{Z}_1(p_i, :, :)^T \mathcal{Z}_1(p_j, :, :)\} \\ \mathbf{v}_2 &= \text{vec}\{\mathcal{Z}_c(:, :, q_i) \mathcal{Z}_c(:, :, q_j)^T\} \\ \mathbf{Q}_k(:, i \times r_{k+1} + j) &= \text{vec}\{(\mathcal{Z}_k(:, :, i)) \mathcal{Z}_k(:, :, j)^T\}, k = 2, \dots, c-1 \\ p_i &= \lceil i/r_{c+1} \rceil, q_i = \text{mod}(i-1, r_{c+1}) + 1 \end{aligned}$$

Then, the Gram matrix of $\mathbf{Z}_{[2]}^{\leq c}$ can be computed as

$$\mathbf{G}_{\mathcal{Z}^{\leq c}} = \mathbf{Z}_{[2]}^{\leq c,T} \mathbf{Z}_{[2]}^{\leq c} = \Phi\left(\prod_{k=1}^c \mathbf{Q}_k\right), \quad (10)$$

where $\mathbf{Q}_k(:, i \times r_{k+1} + j) = \text{vec}\{(\mathcal{Z}_k(:, :, i)) \mathcal{Z}_k(:, :, j)^T\}$ for $k > 1$ and

$$\mathbf{Q}_1(:, i \times r_2 + j) = \begin{cases} \text{vec}\{(\mathcal{Z}_1(:, :, i)) \mathcal{Z}_1(:, :, j)^T\}, & c \text{ is even} \\ \text{vec}\{(\mathcal{Z}_1(:, :, j)) \mathcal{Z}_1(:, :, i)^T\}, & c \text{ is odd} \end{cases}$$

References

Wenqi Wang, Vaneet Aggarwal, and Shuchin Aeron. Efficient low rank tensor ring completion. In *Proceedings of the IEEE International Conference on Computer Vision*, pages 5697–5705, 2017.