

• Supplementary File •

A zeroth-order Algorithm for distributed optimization with stochastic stripe observations

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Appendix A Assumptions

Consider the following distributed optimization problem over an n -agent network [9]:

$$\begin{aligned} \min_{\mathbf{x}} \quad & f(\mathbf{x}), \quad f(\mathbf{x}) = \sum_{i=1}^n f_i(x_i) \\ \text{s. t.} \quad & x_i = x_j, x \in X \end{aligned} \tag{A1}$$

where $\mathbf{x} = [x_1^\top, x_2^\top, \dots, x_n^\top]^\top \in \mathcal{R}^{np}$, $x_i \in \mathcal{R}^p$, $f_i : \mathcal{R}^p \rightarrow \mathcal{R}$ is a convex function, X is the global constraint.

First, we provide an assumption about the communication topology between agents over the network.

Consider a time-varying multi-agent network. The communication topology between agents over the network is described by a directed graph $\mathcal{G}(k) = (\mathcal{N}, \mathcal{E}(k), W(k))$, where $\mathcal{N} = \{1, 2, \dots, n\}$ is the agent set, $\mathcal{E}(k) \subset \mathcal{N} \times \mathcal{N}$ represents information communication links at time k , and $W(k) = [w_{ij}(k)]_{i,j}$ represents the adjacency matrix at time k . In addition, denote $\mathcal{N}_i(k) = \{j | (i, j) \in \mathcal{E}(k)\}$ as the neighbors of agent i at time k . Each agent interacts with its neighbors in $\mathcal{G}(k) = (\mathcal{N}, \mathcal{E}(k), W(k))$ at time k . The following assumption is on the communication topology $\mathcal{G}(k) = (\mathcal{N}, \mathcal{E}(k), W(k))$, which is widely used in distributed time-varying network designs ([8], [9]).

Assumption 1. The graph $\mathcal{G}(k) = (\mathcal{N}, \mathcal{E}(k), W(k))$ satisfies:

- (a) There exists a constant η with $0 < \eta < 1$ such that, $\forall k \geq 0$ and $\forall i, j$, $w_{ii}(k) \geq \eta$; $w_{ij}(k) \geq \eta$ if $(j, i) \in \mathcal{E}(k)$.
- (b) $W(k)$ is doubly stochastic, i. e. $\sum_{i=1}^n w_{ij}(k) = 1$ and $\sum_{j=1}^n w_{ij}(k) = 1$.
- (c) There is an integer $\kappa \geq 1$ such that $\forall k \geq 0$ and $\forall (j, i) \in \mathcal{N} \times \mathcal{N}$, $(j, i) \in \mathcal{E}(k) \cup \mathcal{E}(k+1) \cup \dots \cup \mathcal{E}(k+\kappa-1)$.

The following assumption holds for local functions and constraints of Problem (A1):

Assumption 2. (a) Problem (A1) has solutions.

- (b) $f_{L_i}(x)$ and $f_{R_i}(x)$ are convex functions with $f_{L_i}(x) \leq f_{R_i}(x)$.
- (c) X is a non-empty, compact, convex constraint set in \mathcal{R}^p .
- (d) The gradients of $f_{L_i}(x)$ and $f_{R_i}(x)$ are locally Lipschitz continuous with constant L .

Then, we assume the following assumption holds for parameter λ_i of Problem (A1):

Assumption 3. There exists a common $\lambda_0 \in (0, 1)$, such that for all agents $i \in \mathcal{N}$,

$$g_i(x) = \lambda_0 f_{L_i}(x) + (1 - \lambda_0) f_{R_i}(x).$$

Assumption 3 gives the characteristics of stripe observations $Y_{L_i}(x)$ and $Y_{R_i}(x)$, which results from inherent errors of measuring devices or methods over the network. (\hat{x}^*, λ_0) is the optimal solution to Problem (??), and (\hat{x}^*, λ_0) is expected exact solution to Problem (A1) with stripe observations $Y_L(x)$ and $Y_R(x)$ under Assumption 3.

Next, we give different probabilistic choices of each agents' initial preferences $\lambda_i(0)$ for Algorithm 1.

Assumption 4. (a) Each agent i has an initial preference $\lambda_i(0)$ with $\lambda_0 = \sum_{i=1}^n \lambda_i(0)$.

(b) Each agent i has an initial preference $\lambda_i(0)$. $\lambda_i(0)$ s are independent and identically distributed (i.i.d.) random variables with $\mathbb{E}\lambda_i(0) = \lambda_0$.

(c) Each agent i has an initial preference $\lambda_i(0)$. $\lambda_i(0)$ s are independent and identically distributed (i.i.d.) random variables with $\mathbb{E}\lambda_i(0) = \lambda_0$, $\text{var}\lambda_i(0) = \nu^2$.

Assumption 4(a) is an ideal assumption that although each agent gets incomplete information of $\lambda_i(0)$, they could get the complete information of λ_0 through cooperation. Assumption 4(b) and (c) are general probabilistic assumptions of agents preferences.

Still, the following assumption holds for parameters $\{\Delta_i^q(k)\}_{k \geq 0}$ of Algorithm 1:

Condition 1. (Parameter selection)

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(a) Let $\{\Delta_i^q(k)\}_{k \geq 0}$ be a sequence of independent and identically distributed (i. i. d.) random variables, for any fixed (i, q) , all $k \geq 0$ and (i, q) ,

$$|\Delta_i^q(k)| < M_1, \quad \left| \frac{1}{\Delta_i^q(k)} \right| < M_2, \quad \mathbb{E} \left[\frac{1}{\Delta_i^q(k)} \right] = 0;$$

(b) $\{\Delta_i^q(k)\}_{k \geq 0}$ and $\{\Delta_j^r(k)\}_{k \geq 0}$ are mutually independent of each other for $i \neq j$ or $q \neq r$.

Remark 1. In order to ensure the almost sure convergence of Algorithm 1, $\{c(k)\}$ and $\{\iota(k)\}$ should satisfy the stochastic approximation assumption [29]

$$\sum_{k=1}^{\infty} \frac{\iota(k)}{c(k)} = \infty, \quad \sum_{k=1}^{\infty} \frac{\iota^2(k)}{c^2(k)} < \infty,$$

and

$$\sum_{k=1}^{\infty} \iota(k)c(k) < \infty.$$

Therefore, $0 \leq \epsilon < \frac{1}{4}$, $\frac{1}{2} - \epsilon > \delta > \epsilon$.

With the stripe observation environment of the distributed problem, $\frac{\iota(k)}{c(k)}$ chosen in this paper differs from that of [17].

Appendix B Proof of Theorem 3

The following lemma, is essential for the proof of Theorem 3.

Lemma 1. With Assumption 3, $\|g(\mathbf{x}^*, \boldsymbol{\lambda}(\mathbf{0})) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| \leq |\lambda(0) - \lambda_0| \cdot \|f_L(\hat{\mathbf{x}}^*) - f_R(\hat{\mathbf{x}}^*)\|$.

Proof. We get

$$\begin{aligned} & \|g(\mathbf{x}^*, \boldsymbol{\lambda}(\mathbf{0})) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| \\ &= \sum_{i=1}^n \left[\lambda(0) f_{L_i}(x^*) + (1 - \lambda(0)) f_{R_i}(x^*) \right] - \sum_{i=1}^n \left[\lambda_0 f_{L_i}(\hat{x}^*) + (1 - \lambda_0) f_{R_i}(\hat{x}^*) \right] \leq \|\Gamma_{1n} + \Gamma_{2n}\| \end{aligned} \quad (\text{B1})$$

where

$$\begin{aligned} \Gamma_{1n} &= \sum_{i=1}^n \left[\lambda(0) (f_{L_i}(x^*) - f_{L_i}(\hat{x}^*)) + (1 - \lambda(0)) (f_{R_i}(x^*) - f_{R_i}(\hat{x}^*)) \right] \\ \Gamma_{2n} &= \sum_{i=1}^n \left[(\lambda(0) - \lambda_0) (f_{L_i}(\hat{x}^*) - f_{R_i}(\hat{x}^*)) \right]. \end{aligned}$$

According to Assumption 2(b), we obtain

$$\begin{aligned} f_{L_i}(x^*) - f_{L_i}(\hat{x}^*) &\leq \langle \nabla f_{L_i}(x^*), x^* - \hat{x}^* \rangle \\ f_{R_i}(x^*) - f_{R_i}(\hat{x}^*) &\leq \langle \nabla f_{R_i}(x^*), x^* - \hat{x}^* \rangle. \end{aligned} \quad (\text{B2})$$

It yields

$$\begin{aligned} \|\Gamma_{1n}\| &= \left\| \sum_{i=1}^n \left[\lambda(0) (f_{L_i}(x^*) - f_{L_i}(\hat{x}^*)) + (1 - \lambda(0)) (f_{R_i}(x^*) - f_{R_i}(\hat{x}^*)) \right] \right\| \\ &\leq \left\| \sum_{i=1}^n \left[\lambda(0) f_{L_i}(x^*) + (1 - \lambda(0)) f_{R_i}(x^*) \right] \right\| \cdot \|f_{L_i}(\hat{x}^*) - f_{R_i}(\hat{x}^*)\|. \end{aligned} \quad (\text{B3})$$

Since $(x^*, \lambda(0))$ is an optimal solution of the distributed problem, we have

$$\sum_{i=1}^n \left[\lambda(0) f_{L_i}(x^*) + (1 - \lambda(0)) f_{R_i}(x^*) \right] = 0. \quad (\text{B4})$$

Therefore,

$$\|\Gamma_{1n}\| = 0. \quad (\text{B5})$$

Combining (B5) with (B1), yields

$$\|g(\mathbf{x}^*, \boldsymbol{\lambda}(\mathbf{0})) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| \leq |\lambda(0) - \lambda_0| \cdot \|f_L(\hat{\mathbf{x}}^*) - f_R(\hat{\mathbf{x}}^*)\|. \quad (\text{B6})$$

Lemma 1 gives an upper bound of $\|g(\mathbf{x}^*, \boldsymbol{\lambda}(\mathbf{0})) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\|$.

Next, we provide the proof of Theorem 3.

Proof.

- (a) Theorem 3 is a direct conclusion of Lemma 1 and the fact that $\lambda_0 = \lambda(0)$ with Assumption 4(a).
 (b) According to the law of large numbers [23],

$$\lim_{n \rightarrow \infty} P(\|\lambda(0) - \lambda_0\| < \epsilon) = 1.$$

It yields

$$\lim_{n \rightarrow \infty} P\left(\|g(\mathbf{x}^*, \boldsymbol{\lambda}(0)) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| < \epsilon c_0(\hat{\mathbf{x}}^*)\right) = 1.$$

- (c) The following equality holds according to the central limit theorem [23]:

$$\lim_{n \rightarrow \infty} P\left(\frac{\sqrt{n}(\lambda(0) - \lambda_0)}{\nu} \leq x\right) = \phi(x)$$

where $\phi(x)$ is the cumulative distribution function of Gaussian distribution $\mathcal{N}(0, 1)$. With Lemma 1, we get

$$\begin{aligned} & \lim_{n \rightarrow \infty} P\left(\|g(\mathbf{x}^*, \boldsymbol{\lambda}(0)) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| \leq a c_0(\hat{\mathbf{x}}^*)\right) \geq \lim_{n \rightarrow \infty} P\left(|\lambda(0) - \lambda_0| \leq a\right) \\ & = P\left(\frac{\sqrt{n}|\lambda(0) - \lambda_0|}{\nu} \leq \frac{\sqrt{na}}{\nu}\right) = \phi\left(\frac{\sqrt{na}}{\nu}\right) - \phi\left(-\frac{\sqrt{na}}{\nu}\right) = 2\phi\left(\frac{\sqrt{na}}{\nu}\right) - 1 \end{aligned}$$

and

$$\begin{aligned} & \lim_{n \rightarrow \infty} P\left(\|g(\mathbf{x}^*, \boldsymbol{\lambda}(0)) - g(\hat{\mathbf{x}}^*, \boldsymbol{\lambda}_0)\| \geq b c_0(\hat{\mathbf{x}}^*)\right) \leq \lim_{n \rightarrow \infty} P\left(|\lambda(0) - \lambda_0| \geq b\right) \\ & = P\left(\frac{\sqrt{n}|\lambda(0) - \lambda_0|}{\nu} \geq \frac{\sqrt{nb}}{\nu}\right) = 1 - \phi\left(\frac{\sqrt{nb}}{\nu}\right) + \phi\left(-\frac{\sqrt{nb}}{\nu}\right) = 2 - 2\phi\left(\frac{\sqrt{nb}}{\nu}\right). \end{aligned}$$

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