

# Supplementary Information for “Quantum key distribution with simply characterized light sources”

Akihiro Mizutani<sup>\*,1</sup> Toshihiko Sasaki,<sup>2</sup> Yuki Takeuchi,<sup>3</sup> Kiyoshi Tamaki,<sup>4</sup> and Masato Koashi<sup>2</sup>

<sup>1</sup>*Mitsubishi Electric Corporation, Information Technology R&D Center,  
5-1-1 Ofuna, Kamakura-shi, Kanagawa, 247-8501 Japan*

<sup>2</sup>*Photon Science Center, Graduate School of Engineering,  
The University of Tokyo, Bunkyo-ku, Tokyo 113-8656, Japan*

<sup>3</sup>*NTT Communication Science Laboratories, NTT Corporation,*

<sup>3-1</sup>*Morinosato Wakamiya, Atsugi-shi, Kanagawa 243-0198, Japan*

<sup>4</sup>*Graduate School of Science and Engineering for Research,  
University of Toyama, Gofuku 3190, Toyama, 930-8555, Japan*

\*Mizutani.Akihiro@dy.MitsubishiElectric.co.jp

In the Supplementary Information, we prove Lemmas 1 and 2 in the main text.

## I. PROOF OF LEMMA 1

### Lemma 1

$$\hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1 \leq \lambda \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 \quad (\text{I.1})$$

with  $\lambda := 3 + \sqrt{5}$ .

*Proof.* We first explicitly describe  $\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1$  and  $\hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1$  by respectively using Eqs. (18) and (20) in the main text as

$$\begin{aligned} \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 &= \hat{P} \left[ \frac{|001\rangle_A}{\sqrt{2}} \right] \otimes \left( \hat{P}[|1\rangle_B] + \frac{1}{2} \hat{P}[|2\rangle_B] \right) + \hat{P} \left[ \frac{|100\rangle_A}{\sqrt{2}} \right] \otimes \left( \frac{1}{2} \hat{P}[|2\rangle_B] + \hat{P}[|3\rangle_B] \right) \\ &+ \sum_{s=0}^1 \hat{P} \left[ \frac{|010\rangle_A + (-1)^s |100\rangle_A}{\sqrt{2}} \right] \otimes \hat{\Pi}_{1,s \oplus 1} + \sum_{s=0}^1 \hat{P} \left[ \frac{|001\rangle_A + (-1)^s |010\rangle_A}{\sqrt{2}} \right] \otimes \Pi_{2,s \oplus 1}, \end{aligned} \quad (\text{I.2})$$

$$\hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1 = (\hat{P}[|001\rangle_A] + \hat{P}[|100\rangle_A]) \otimes \frac{\hat{P}[|2\rangle_B]}{2} + \hat{P}[|010\rangle_A] \otimes (\hat{P}[|1\rangle_B] + \hat{P}[|3\rangle_B]). \quad (\text{I.3})$$

In Eq. (I.2), it is straightforward to show that

$$\begin{aligned} \sum_{s=0}^1 \hat{P} \left[ \frac{|010\rangle_A + (-1)^s |100\rangle_A}{\sqrt{2}} \right] \otimes \hat{\Pi}_{1,s \oplus 1} &= \hat{P} \left[ \frac{|100\rangle_A |1\rangle_B - \frac{|010\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right] + \hat{P} \left[ \frac{|010\rangle_A |1\rangle_B - \frac{|100\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right], \\ \sum_{s=0}^1 \hat{P} \left[ \frac{|001\rangle_A + (-1)^s |010\rangle_A}{\sqrt{2}} \right] \otimes \hat{\Pi}_{2,s \oplus 1} &= \hat{P} \left[ \frac{|001\rangle_A |3\rangle_B - \frac{|010\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right] + \hat{P} \left[ \frac{|010\rangle_A |3\rangle_B - \frac{|001\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right]. \end{aligned}$$

To upper-bound  $\hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1$  by using  $\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1$ , we remove the four projectors in  $\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1$  that are orthogonal to the range of  $\hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1$ , which results in

$$\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 \geq \frac{1}{2} \left( \hat{P} \left[ \frac{|001\rangle_A}{\sqrt{2}} \right] + \hat{P} \left[ \frac{|100\rangle_A}{\sqrt{2}} \right] \right) \otimes \hat{P}[|2\rangle_B] + \hat{P} \left[ \frac{|010\rangle_A |1\rangle_B - \frac{|100\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right] + \hat{P} \left[ \frac{|010\rangle_A |3\rangle_B - \frac{|001\rangle_A |2\rangle_B}{\sqrt{2}}}{\sqrt{2}} \right]. \quad (\text{I.4})$$

Moreover, we apply the following inequality that holds for any normalized vectors  $|a\rangle$  and  $|b\rangle$  with  $\langle a|b\rangle = 0$ <sup>1</sup>,

$$\hat{P}\left[|a\rangle - \frac{|b\rangle}{\sqrt{2}}\right] \geq \frac{2}{\lambda} \left( \hat{P}[|a\rangle] + \hat{P}\left[\frac{|b\rangle}{\sqrt{2}}\right] \right) - \hat{P}\left[\frac{|b\rangle}{\sqrt{2}}\right] \quad (\text{I.6})$$

with  $\lambda := 3 + \sqrt{5}$ , to the last two projectors of the rhs in Eq. (I.4) and obtain

$$\begin{aligned} \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 &\geq \frac{1}{2} \left( \hat{P}\left[\frac{|001\rangle_A}{\sqrt{2}}\right] + \hat{P}\left[\frac{|100\rangle_A}{\sqrt{2}}\right] \right) \otimes \hat{P}[|2\rangle_B] \\ &+ \frac{1}{\lambda} \left( P[|010\rangle_A |1\rangle_B] + P\left[\frac{|100\rangle_A |2\rangle_B}{\sqrt{2}}\right] \right) - \frac{\hat{P}[|100\rangle_A |2\rangle_B]}{4} \\ &+ \frac{1}{\lambda} \left( P[|010\rangle_A |3\rangle_B] + P\left[\frac{|001\rangle_A |2\rangle_B}{\sqrt{2}}\right] \right) - \frac{\hat{P}[|001\rangle_A |2\rangle_B]}{4} \\ &= \hat{P}_1 \hat{e}_{\text{ph}} \hat{P}_1 / \lambda. \end{aligned} \quad (\text{I.7})$$

This ends the proof of Lemma 1. ■

## II. PROOF OF LEMMA 2

**Lemma 2** For any density operator  $\hat{\sigma}$ ,

$$\text{tr} \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 \hat{\sigma} \leq \text{tr} \hat{e}_{\text{bit}} \hat{\sigma} + \sqrt{\text{tr} \hat{\sigma} \hat{P}_1 \cdot \text{tr} \hat{\sigma} \hat{P}_3}. \quad (\text{II.1})$$

*Proof.* From Eq. (19), for any state  $\hat{\sigma}$  we have

$$\text{tr} \hat{P}_{\text{odd}} \hat{e}_{\text{bit}} \hat{P}_{\text{odd}} \hat{\sigma} \leq \text{tr} \hat{e}_{\text{bit}} \hat{\sigma}, \quad (\text{II.2})$$

which leads to

$$\text{tr} \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 \hat{\sigma} \leq \text{tr} \hat{e}_{\text{bit}} \hat{\sigma} - \text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 + \hat{P}_3 \hat{e}_{\text{bit}} \hat{P}_1) \hat{\sigma}. \quad (\text{II.3})$$

Since  $-\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 + \hat{P}_3 \hat{e}_{\text{bit}} \hat{P}_1) \hat{\sigma} \leq |\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 \hat{\sigma})| + |\text{tr}(\hat{P}_3 \hat{e}_{\text{bit}} \hat{P}_1 \hat{\sigma})| = 2|\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 \hat{\sigma})|^2$ , we derive an upper bound on  $|\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 \hat{\sigma})|$ . From the expression of the POVM element  $\hat{e}_{\text{bit}}^j$  given by Eq. (18), we have

$$\hat{T} := 2\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 = |001\rangle\langle 111|_A \otimes (\hat{\Pi}_{1,1} - \hat{\Pi}_{1,0}) + |100\rangle\langle 111|_A \otimes (\hat{\Pi}_{2,1} - \hat{\Pi}_{2,0}). \quad (\text{II.4})$$

As  $(\hat{\Pi}_{1,1} - \hat{\Pi}_{1,0})^2 = (|1\rangle\langle 1| + |2\rangle\langle 2|)/2$  and  $(\hat{\Pi}_{2,1} - \hat{\Pi}_{2,0})^2 = (|2\rangle\langle 2| + |3\rangle\langle 3|)/2$ , we obtain

$$\begin{aligned} \hat{T}^\dagger \hat{T} &= \hat{P}[|111\rangle_A] \otimes [(\hat{\Pi}_{1,1} - \hat{\Pi}_{1,0})^2 + (\hat{\Pi}_{2,1} - \hat{\Pi}_{2,0})^2] \\ &\leq \hat{I}_{AB}. \end{aligned} \quad (\text{II.5})$$

This inequality implies that the operator norm of  $\hat{T}$  is upper-bounded by 1:

$$\|\hat{T}\|_\infty := \min\{c \geq 0 \text{ s.t. } \forall v \|\hat{T}v\| \leq c\|v\|\} \leq 1, \quad (\text{II.6})$$

where  $\|\cdot\| := \sqrt{\langle \cdot | \cdot \rangle}$ . Next, we define

$$\hat{G} := \hat{P}_3 \hat{\sigma} \hat{P}_1. \quad (\text{II.7})$$

<sup>1</sup> Note that Eq. (I.6) holds because the smallest eigenvalue of the following Hermitian operator:

$$\hat{P}\left[|a\rangle - \frac{|b\rangle}{\sqrt{2}}\right] - \frac{2}{\lambda} \left( \hat{P}[|a\rangle] + \hat{P}\left[\frac{|b\rangle}{\sqrt{2}}\right] \right) + \hat{P}\left[\frac{|b\rangle}{\sqrt{2}}\right] \quad (\text{I.5})$$

is zero.

<sup>2</sup> The last equality comes from the fact that  $|\text{tr} \hat{A}| = |\text{tr} \hat{A}^\dagger|$  holds for any square matrix  $\hat{A}$ .

Its trace norm  $\|\hat{G}\|_1$  is written by using a unitary operator  $\hat{W}$  and is calculated as

$$\begin{aligned} \|\hat{G}\|_1 &= |\text{tr} \hat{G} \hat{W}| = |(\sqrt{\hat{\sigma}} \hat{P}_3, \sqrt{\hat{\sigma}} \hat{P}_1 \hat{W})| \\ &\leq \sqrt{(\sqrt{\hat{\sigma}} \hat{P}_3, \sqrt{\hat{\sigma}} \hat{P}_3)} \sqrt{(\sqrt{\hat{\sigma}} \hat{P}_1 \hat{W}, \sqrt{\hat{\sigma}} \hat{P}_1 \hat{W})} \\ &= \sqrt{\text{tr} \hat{P}_3 \hat{\sigma}} \sqrt{\text{tr} \hat{P}_1 \hat{\sigma}}, \end{aligned} \quad (\text{II.8})$$

where we use the definition of Hilbert-Schmidt inner product in the second equality and use Schwarz inequality in the first inequality. Finally, using Hölder's inequality, Eqs. (II.6) and (II.8) gives

$$\begin{aligned} 2|\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 \hat{\sigma})| &= |\text{tr} \hat{T} \hat{G}| \leq \|\hat{T} \hat{G}\|_1 \leq \|\hat{T}\|_\infty \|\hat{G}\|_1 \\ &\leq \sqrt{\text{tr} \hat{P}_3 \hat{\sigma} \cdot \text{tr} \hat{P}_1 \hat{\sigma}}. \end{aligned} \quad (\text{II.9})$$

Therefore, we obtain

$$-\text{tr}(\hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_3 + \text{tr} \hat{P}_3 \hat{e}_{\text{bit}} \hat{P}_1) \hat{\sigma} \leq \sqrt{\text{tr} \hat{P}_1 \hat{\sigma} \cdot \text{tr} \hat{P}_3 \hat{\sigma}}. \quad (\text{II.10})$$

Finally, by using Eqs. (II.3) and (II.10), we conclude that

$$\text{tr} \hat{P}_1 \hat{e}_{\text{bit}} \hat{P}_1 \hat{\sigma} \leq \text{tr} \hat{e}_{\text{bit}} \hat{\sigma} + \sqrt{\text{tr} \hat{\sigma} \hat{P}_1 \cdot \text{tr} \hat{\sigma} \hat{P}_3}. \quad (\text{II.11})$$

This ends the proof of Lemma 2.  $\blacksquare$