# Distributed Optimal Quantization and Power Allocation for Sensor Detection Via Consensus

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Abstract—We address the optimal transmit power allocation problem (from the sensor nodes (SNs) to the fusion center (FC)) for the decentralized detection of an unknown deterministic spatially uncorrelated signal which is being observed by a distributed wireless sensor network. We propose a novel fully distributed algorithm, in order to calculate the optimal transmit power allocation for each sensor node (SN) and the optimal number of quantization bits for the test statistic in order to match the channel capacity. The SNs send their quantized information over orthogonal uncorrelated channels to the FC which linearly combines them and makes a final decision. What makes this

scheme attractive is that the SNs share with their neighbours just their individual transmit powers at the current states. As a

result, the SN processing complexity is further reduced.

Keywords—Distributed detection, distributed processing, soft decision, wireless sensor networks.

### I. Introduction

Wireless sensor networks (WSNs) are spatially deployed over a field to monitor certain physical or environmental phenomena. Generally, the sensing process is orientated towards estimating various parameters of interest which can be employed to arrive at a certain decision. This decision can then be relayed in a pre-specified manner or can be employed for onfield actuation. We note that the reliable and continued operation of a WSN over many years is often desirable. This is due to the operational environment in which post-deployment access to a sensor node (SN) is at best very limited. Unfortunately, SNs suffer from constrained bandwidth and limited available on-board power. Moreover, due to the locality

$$\mathsf{E}\{T_i|\mathcal{H}_0\} = N\sigma_i^2$$
,  $\mathsf{Var}\{T_i|\mathcal{H}_0\} = 2N\sigma_i^4$ ,  $\mathsf{E}\{T_i|\mathcal{H}_1\} = N\sigma_i^2 (1+\xi_i)$ ,  $\mathsf{Var}\{T_i|\mathcal{H}_1\} = 2N\sigma_i^4 (1+2\xi_i)$ .

of the observed process, cooperation amongst SNs is often required to derive an inference. However, such a cooperation comes at the expense of high bandwidth requirements and signalizing overhead. For instance, a WSN formed by  $\mathbb Z$  sensor nodes would require transmission of  $\mathbb Z(\mathbb Z^2)$  message exchanges to attain full cooperation. Consequently, designing distributed detection algorithms that efficiently utilize the scarce bandwidth and cope with the impairments in a wireless channel is very important.

This work investigates the detection performance of the SN over flat fading wireless transmission links. A centralized solution (taken at the fusion center (FC)) is proposed in [1] where the deterministic signal (s) to be detected is assumed to be known a-priori. We relax this constraint by deriving a scheme that detects an unknown deterministic signal (s) by employing a linear fusion rule at the FC and adopting the modified deflection coefficient as the detection performance criterion. We also propose a fully distributed algorithm where we allocate the SN transmit power for each individual SN using only local information.

The problem of decentralized detection (and estimation) in a WSN has been extensively tackled in [1]-[7], to name but just a few. Recent publications [8]-[9] propose a distributed algorithm for in-network estimation of algebraic connectivity. Interestingly, [9] uses an estimation strategy to adapt the SN transmit power in order to maximize the connectivity of the network, while in this paper we take advantage of the objective function structure and develop a novel distributed algorithm to allocate the SN to FC transmit power. The algorithm is very efficient in terms of convergence and data exchange, also accurate and simple to implement.

Section II describes the system model and we derive an approach that utilizes the SN to FC channel capacity. An optimum linear combining rule is adopted at the FC with the combining weights optimized in Section III. Section III presents the derivation of the decentralized optimum SN transmit power allocation and our proposed algorithm. Finally, simulation results are given in Section IV and conclusions in Section V.

# II. SYSTEM MODEL AND QUANTIZED DECISION COMBINING

Consider the problem of detecting the presence of a deterministic signal  $\mathbb{Z}(\mathbb{Z})$  by a sensor network consisting of  $\mathbb{Z}$  SNs. The  $\mathbb{Z}^{\mathbb{Z}\mathbb{Z}}$  SN collects  $\mathbb{Z}$  samples of the observed signal

(2n(2)), and so the detection problem can be formulated as a binary hypothesis test as follows:

$$H_0: \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) = \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) \tag{1}$$

$$H_1: \mathbb{Z}_2(\mathbb{Z}) = \mathbb{Z}_2(\mathbb{Z}) + \mathbb{Z}_2(\mathbb{Z})(2)$$
 where  $\mathbb{Z}_2(\mathbb{Z}) \sim$ 

 $N(0222^2)$  is AWGN and 22(2) is the observation of (2(2)),

both at the 200 node. The 200 SN then estimates the energy:

$$T_i = \sum_{n=1}^{N} |x_i(n)|^2, \ i = 1, 2, \dots, M$$
 (3)

which for large 2 can be approximated by a Gaussian  $1 - N\sigma^2 (1 + \xi_1) - V_{2r} (T_1 | \mathcal{H}_1) - 2N\sigma^4 (1 + 2\xi_1)$ 

distribution [10] under both hypothesis. So is not difficult to derive

(4) where  $\sum_{n=1}^{\mathbb{Z}} s_i^2(n)/N\sigma_i^2$  which can be considered as effective observed SNR. Now linear soft decision combining at the FC has superior performance to the hard decision approach, but it entails additional complexity. In addition soft decision combining puts additional demands on both the limited power resources of the SNs and the effective utilization of the SN to FC channel capacity. So here we propose a scheme, where each individual SN has to quantize its observed test statistic ( $\mathbb{Z}$ ) to  $\mathbb{Z}$  bits. The number of quantization bits at the  $\mathbb{Z}$  SN must satisfy the channel capacity constraint:

$$L_i \le \frac{1}{2} \log_2 \left( 1 + \frac{p_i h_i^2}{\zeta_i} \right)_{\text{bits}} \tag{5}$$

$$?^{}_{2} = ?_{2} + ?_{2}$$
 (6)

where  $\mathbb{Z}_{\mathbb{Z}}$  is quantization noise (variance,  $\mathbb{Z}_{\mathbb{Z}}$ ) independent of  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) in ((1) and (2)). Assuming quantization noise with a uniform distribution and  $\mathbb{Z}_{\mathbb{Z}} \in [0\mathbb{Z}\mathbb{Z}]$ , then

$$\sigma_{v_i}^2 = \frac{U^2}{3 \times 2^{2L_i}}.\tag{7}$$

Linearly combining  $\left\{ \hat{T}_i \right\}_{i=1}^M$  at the FC gives

$$T_f = \sum_{i=1}^{M} \alpha_i \hat{T}_i \tag{8}$$

where the weights  $\{\alpha_i\}_{i=1}^M$  will be optimized in Section III. Again, for large  $\mathbb{Z}$ ,  $\mathbb{Z}_{\mathbb{Z}}$  will be approximately Gaussian and so we can derive (9) and (10). We now define  $\mathbb{Z} = \mathbb{E}\{\mathbb{Z}_{\mathbb{Z}} | \mathbb{H}_1\}$ -

$$\{T_f|\mathcal{H}_0\} = N \sum_{i=1}^{M} \alpha_i \left(\sigma_i^2 \xi_i\right)$$
 Eand for a fixed  $\mathbb{Z}_{\mathbb{Z}}$  (probability

of false alarm) we can write [11]:

$$P_d = Q\left(\frac{Q^{-1}(P_{fa})\sqrt{\operatorname{Var}\left\{T_f|\mathcal{H}_0\right\}} - \psi}{\sqrt{\operatorname{Var}\left\{T_f|\mathcal{H}_1\right\}}}\right)$$
(11)

where  $\square$  is the probability of detection. So using (5), (7), (8), (9) and (10) in (11) we get

$$P_{d} = Q \left( \frac{Q^{-1} \left( P_{fa} \right) \sqrt{\sum_{i=1}^{M} \alpha_{i}^{2} \left( 2N\sigma_{i}^{4} + \frac{U^{2}}{3\left( 1 + \frac{p_{i}h_{i}^{2}}{\zeta_{i}} \right)} \right) - \psi}}{\sqrt{\sum_{i=1}^{M} \alpha_{i}^{2} \left[ 2N\sigma_{i}^{4} \left( 1 + 2\xi_{i} \right) + \frac{U^{2}}{3\left( 1 + \frac{p_{i}h_{i}^{2}}{\zeta_{i}} \right)} \right]}} \right). \quad (12)$$

The formula in (12) imposes a relationship between the probability of detection, the power allocated to each transmission link (SN to the FC) and the weight (21 in (8)) for each individual link.

### III. DECENTRALIZED OPTIMUM WEIGHT COMBINING AND POWER ALLOCATION

We would now like to find the optimum weighting vector  $(\alpha^{\square})$  and the optimum power allocation vector  $(\mathfrak{p}^{\square})$  that achieves the best possible 2 (see definitions later), under the constraint of a maximum transmit power budget (2a). However, maximizing (12) w.r.t. α and p is difficult and no closed form solution can be found. From (12) it is straightforward to observe that the 22 is a monotonically increasing function of the deflection coefficient. Moreover, Var{2 | H0} 2 Var{2 | H1}. Employing these two facts, it is intuitive to approximate the optimization problem of (12) by maximization of the deflection coefficient which is given as:

$$\tilde{d}^{2}(\boldsymbol{\alpha}, \boldsymbol{p}) = \left(\frac{\mathrm{E}\left\{T_{f}|\mathcal{H}_{1}\right\} - \mathrm{E}\left\{T_{f}|\mathcal{H}_{0}\right\}}{\sqrt{\mathrm{Var}\left\{T_{f}|\mathcal{H}_{1}\right\}}}\right)^{2} = \frac{\left(\boldsymbol{b}^{T}\boldsymbol{\alpha}\right)^{2}}{\boldsymbol{\alpha}^{T}\mathbf{R}\boldsymbol{\alpha}}$$
(13)

where 
$$\mathbf{b} = [N\sigma_1^2\xi_1, N\sigma_2^2\xi_2, \dots, N\sigma_M^2\xi_M]^T_{\mathbf{Q}}$$
 =

 $[2_1??_2???????]^2, p = [2_1??_2???????]^2$ 

$$\operatorname*{R}=2N\mathrm{diag}\left(\sigma_{1}^{4}\left(1+2\xi_{1}\right)+\frac{\sigma_{v_{1}}^{2}}{2N},...,\ \sigma_{M}^{4}\left(1+2\xi_{M}\right)+\frac{\sigma_{v_{M}}^{2}}{2N}\right).$$

Note that the dependence of  $\mathbb{Z}^2(\alpha\mathbb{Z}p)$  on the transmit power vector p enters (13) through the  $\left\{\sigma_{v_i}^2\right\}_{i=1}^{M}$  terms via (5) and (7). Now, our optimization problem is:

$$(\boldsymbol{\alpha}^{o}, \boldsymbol{p}^{o}) = \underset{\max}{\operatorname{arg max}} 2^{\circ} 2 (\alpha 2 p)'$$

$$p \qquad \qquad \alpha, \qquad \qquad (P1)$$

The straightforward solution to (P1) is to obtain it in a centralized manner (i.e., at a FC), where the FC has full knowledge of the channel gains (22) which might change over time and need to be updated. The dependence of R on the flat

fading channel coefficients  $\left(\left\{h_i\right\}_{i=1}^M\right)$  enters through

 $\{\sigma_{v_i}^2\}_{i=1}^M$ . In this paper we propose a distributed solution, where the SNs are limited to use local information to be able to decide if they should transmit any information to the FC or stay in sleeping mode.

# A. Optimisation through Decentralized Weight Combining

Letting  $\beta = R^{1/2}\alpha$  in (13), then we have

$$\tilde{d^{2}}\left(\boldsymbol{\beta},\boldsymbol{p}\right) = \frac{\boldsymbol{\beta}^{T}\boldsymbol{D}\boldsymbol{\beta}}{||\boldsymbol{\beta}||^{2}}, \quad \mathbf{D} = \mathbf{R} - \mathbf{D}\mathbf{D}\mathbf{D}\mathbf{D}\mathbf{D}\mathbf{B}\mathbf{R} - \mathbf{D}\mathbf{D}$$
(14)

and  $\alpha^{2} = R^{-1/2}\beta_{2/2}$  in (P1) (assuming p is constant), where  $\beta_{2/2}$ is the eigenvector corresponding to the maximum eigenvalue of D. So we can easily show that:

$$\boldsymbol{\alpha}^{o} = \begin{bmatrix} \frac{N\sigma_{1}^{2}\xi_{1}}{2N\sigma_{1}^{4}(1+2\xi_{1})+\sigma_{v_{1}}^{2}}, \cdots, \frac{N\sigma_{M}^{2}\xi_{M}}{2N\sigma_{M}^{4}(1+2\xi_{M})+\sigma_{v_{M}}^{2}} \end{bmatrix}$$
(15)

Note that (15) establishes a relationship between the optimum weighting vector  $\alpha^{\square}$  and the individual sensor transmit powers through  $\sigma_{v_i}^2$  quantity (7).

## B. Decentralized Optimum Power Allocation

We now propose a novel algorithm aimed at allocating the sensor transmit power to the FC in a fully decentralized fashion.

Substitute 
$$\alpha^{\text{\tiny II}}$$
 (15) into 
$$\sum_{i=1}^{M} \frac{N^2 \sigma_i^4 \xi_i^2}{2N \sigma_i^4 \left(1 + 2\xi_i\right) + \frac{U^2}{3\left(1 + \frac{p_i h_i^2}{\zeta_i}\right)}}$$
 (P1) to maximize  $_{\text{\tiny P}}$  (P2)

$$\sum_{i=1} p_i \le P_t, \ p_i \ge 0, \ i=1,\ldots,M.$$
 subject to  $i=1$ 

Now (P2) can be solved using the Lagrangian: 
$$f(\boldsymbol{p},\lambda_0,\mu) = \sum_{i=1}^{M} \frac{N^2 \sigma_i^4 \xi_i^2}{2N \sigma_i^4 \left(1 + 2\xi_i\right) + \frac{U^2}{3\left(1 + \frac{p_i h_i^2}{\zeta_i}\right)}} - \lambda_0 \left(\sum_{i=1}^{M} p_i - P_t\right) + \sum_{i=1}^{M} \mu_i p_i$$

and imposing the Karush-Kuhn-Tucker (K.K.T) conditions [14]:

$$p_{i}^{o} = \left[ \frac{1}{\sqrt{\lambda_{0}}} \left( \frac{\xi_{i} U \sqrt{3}}{6\sigma_{i}^{2} (1 + 2\xi_{i}) \sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \right) - \frac{U^{2}}{6N\sigma_{i}^{4} (1 + 2\xi_{i}) \frac{h_{i}^{2}}{\zeta_{i}}} - \frac{\zeta_{i}}{h_{i}^{2}} \right]^{+}$$
(24)

where

$$[?] + = \frac{1}{2}$$
 if  $? 0$ 

$$[?] + = \frac{1}{2}$$

$$0?$$
 if  $? ? 0?$ 

As mentioned before, the centralized solution at the FC requires full knowledge of the channel gains (2) which might be timevarying and need to be always updated. It also requires the

$$\operatorname{Var}\left\{\mathbb{Z}_{\mathbb{Z}}\left|H_{0}\right\}\right\} = \frac{X^{\mathbb{Z}}}{\mathbb{Z}_{\mathbb{Z}}^{2}} i \frac{2\mathbb{Z}_{\mathbb{Z}}^{4} + \mathbb{Z}_{\mathbb{Z}_{\mathbb{Z}}}^{2}}{\mathbb{Z}} i \operatorname{Var}\left\{\mathbb{Z}_{\mathbb{Z}}\left|H_{1}\right\}\right\} = \frac{X^{\mathbb{Z}}}{\mathbb{Z}_{\mathbb{Z}}^{2}} 2\mathbb{Z}_{\mathbb{Z}}^{4} \left(1+2\mathbb{Z}_{\mathbb{Z}}\right) + \mathbb{Z}_{\mathbb{Z}_{\mathbb{Z}}}^{2} \mathbb{Z}$$

$$\stackrel{\mathbb{Z}}{\mathbb{Z}} = 1$$

$$(10)$$

$$0 \in \frac{N^{2} \sigma_{i}^{4} \xi_{i}^{2}}{\left(2N \sigma_{i}^{4} \left(1 + 2\xi_{i}\right) + \frac{U^{2}}{3\left(1 + \frac{p_{i}h_{i}^{2}}{\zeta_{i}}\right)}\right)^{2}} \times \frac{U^{2} \times \frac{h_{i}^{2}}{\zeta_{i}}}{3\left(1 + \frac{p_{i}h_{i}^{2}}{\zeta_{i}}\right)^{2}} - \lambda_{0} + \mu_{i}}$$

$$\lambda_{0} \left(\sum_{i=1}^{M} p_{i} - P_{t}\right) = 0$$

$$\sum_{i=1}^{M} p_{i} - P_{t} \leq 0$$

$$\lambda_{0} \geq 0, \ \mu_{i}p_{i} = 0, \ i = 1, 2, \dots, M$$

$$(22)$$

$$f(\boldsymbol{p},\lambda_0,\mu) = \sum_{i=1}^{M} f_i(p_i,\lambda_0) = \sum_{i=1}^{M} \left(\frac{N^2 \sigma_i^4 \xi_i^2}{2N \sigma_i^4 \left(1+2\xi_i\right) + \frac{U^2}{3\left(1+\frac{p_i h_i^2}{\zeta_i}\right)}} - \lambda_0 p_i + \frac{\lambda_0}{M} P_t\right).$$

Now, (P2) is converted into 2 separable problems that can be solved in parallel using the dual ascent algorithm:

$$2 \operatorname{min} 2 \operatorname{min} 2$$

$$\lambda_0 [k+1] = \lambda_0 [k] + \epsilon [k] \left( \sum_{i=1}^M p_i [k+1] - P_t \right).$$
 (a2)

For this formulation we can see that the only step that requires an exchange of values among the sensors is the (a2) step which

requires the computation of 
$$\sum_{i=1}^{M} p_i [k+1]$$
 at each

sensor node. Because of the communication topology for the 2 SNs (i.e., not fully connected), we will use the average consensus algorithm [12] to ensure the availability of this term at each SN. In this paper, we assume ideal exchange of information between sensors that are connected. Solving the K.K.T conditions in (22) and (23) gives a solution for the optimum 22:

variance of AWGN (2a) and each of the local SNRs (2a). Moreover, the FC has to broadcast back to each individual SN the allocated SN transmit power which might be decoded with error due to fading. Furthermore, when the FC is battery operated, the centralized solution (at the FC) becomes inefficient and not scalable as the number of SNs increases. On the other hand, the proposed distributed algorithm (Algorithm 1) is fully scalable in terms of data exchange and SN processing complexity. As is shown in the simulation result it is also very

accurate. We now define e[k] accurate. We now define  $p_i[k+1] = \frac{1}{M} \sum_{i=1}^{M} p_i[k+1]$ . The convergence criteria that we use in here is the relative absolute difference:

 $\frac{||p[k+1]-p[k]||}{||p[k]||} \le \kappa$ , where 2 is a positive small constant and p[2] is the vector of the SN transmit power at the 222 iteration.

It is also possible to exchange among the SNs the  $\mathbb{Z}_{\mathbb{Z}}$ ,  $\mathbb{Z}_{\mathbb{Z}}$ , 2 nd 2 quantities ∀2 where each SN will store them in the corresponding vectors h,  $\zeta$ ,  $\sigma^2$ , and  $\xi$  together with their corresponding SN index. When all the quantities will be available at each SN they can be used to allocate the SN transmit power through (24) and  $\mathbb{Z}_0$  can be calculated through the constraint in (P2)

## Algorithm1: Optimizing the sensor transmit power

STEP 1: Set 2 = 0, 2 = 0 equal to a small positive value and initialize  $\mathbb{Z}_0$  [0],  $\forall \mathbb{Z}$ ;

STEP 2: Compute 2 [1], ∀2 using (a1);

STEP 3: Run consensus over  $\mathbb{Z}_{\mathbb{Z}}[1]$  to get  $p_i$ [1];

STEP 4: Compute  $\mathbb{Z}_0$  [1] using (a2);

STEP 5: Set ② = 1;

STEP 6: Repeat until convergence

$$p_{i}[k+1] \!=\! \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right]^{+} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! - \frac{\zeta_{i}}{h_{i}^{2}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{4}(1\!+\!2\xi_{i})\frac{h_{i}^{2}}{\zeta_{i}}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{6\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right) \! - \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! \right] \! + \! \left[ \! \frac{1}{\sqrt{\lambda_{0}[k]}} \! \left( \! \frac{\xi_{i}U\sqrt{3}}{2\sigma_{i}} \! \right) \! + \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}} \! \right] \! + \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! + \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! + \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! + \! \frac{U^{2}}{6N\sigma_{i}^{2}(1\!+\!2\xi_{i})\sqrt{\frac{h_{i}^{2}}{\zeta_{i}}}} \! + \! \frac{U^{2}}{$$

Run consensus over 
$$2 \, \mathbb{Z} \, [2+1]$$
 until convergence  $\lambda_0 \, [k+1] = \lambda_0 \, [k] + \epsilon \, [k] \, \Big( M \overline{p_i \, [k+1]} - P_t$ ,

### IV. SIMULATION RESULTS

In this section, the proposed algorithm is evaluated numerically and compared to its centralized counterpart. Also, we choose  $\mathbb{Z}_0[0] = 10^{-8}$ ,  $\forall \mathbb{Z}$ ,  $\mathbb{Z} = 10^{-7}$  and  $\mathbb{Z}[\mathbb{Z}] = \mathbb{Z}_0[\mathbb{Z}]\mathbb{Z}\mathbb{Z}$ .

We let all the  $\mathbb{Z}_{\mathbb{Z}}^2$  terms at each SN be different, such that  $\xi_a = 10 \log_{10} \left( \frac{1}{M} \sum\limits_{i}^{M} \xi_i \right) = \text{-4 dB, unless otherwise stated.}$  In addition we let  $\mathbb{Z}_{\mathbb{Z}} = 0 \mathbb{Z} 1 \ \forall \mathbb{Z}$ . We compare the results with

the matched filter detector<sup>1</sup> (MFD) and use this as a benchmark.

We will also refer to "equal linear combining" in (8) (i.e.,

 $\alpha_i = \frac{1}{\sqrt{M}}, \forall i$ ) and "equal power allocation" in (5) (i.e.,  $p_i = \frac{P_t}{M}, \forall i$ ). Finally, we choose  $\mathbb{Z}_{\mathbb{Z}}$  with equality in (5). In Fig. 1, the middle plot shows the SN transmit power 2 for the 22 SN to the FC channel using two different approaches (i.e., distributed and centralized). The actual channel coefficients (randomly chosen) are in the upper plot in Fig. 1. Clearly, the performance of our proposed distributed method is very close to the centralized one. As expected, both centralized and decentralized methods allocate more power to the best channels. In this way, the nodes that have very bad channels (i.e.,nodes that require very high power to transmit) will be censored (i.e., will not transmit even a single bit). In Fig. ??, we show that for large number of samples (2) the optimum power allocation scheme tends to a uniform power allocation as expected (see the definition of R in Section III). Fig. 3 shows the total power budget (2a) against the mis-detection (1-2a)performance for 6 different schemes. The energy detector (ED) performance tends to converge to the matched filter detector for

$$T_i = \sum\limits_{n=1}^{rac{n}{2}} x_i(n) s_i(n), \ \ orall i=1,2,\ldots,M$$
 .

The global test statistic (  $\mbox{\fontfamily{1}{10}}\mbox{\fontfamil$ 

$$\alpha_i = \frac{\sum\limits_{n=1}^{n} s_i^2(n)}{\sigma_i^2 \sum\limits_{n=1}^{N} s_i^2(n) + \sigma_{v_i}^2}, \ \ \forall i = 1, 2, \dots, M.$$
 The optimum weights have been derived through the Likelihood Ratio Test (LRT)

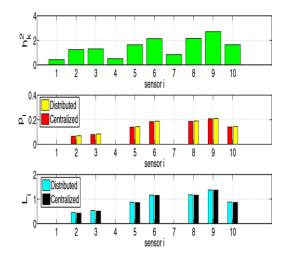


Fig. 1.  $P_{fa}=0.1,\ P_t=1,\ U=3,\ \xi_a=-4\ \text{dB},\ N=10$  Centralized and decentralized sensor transmit power and channel bit allocation for and \$\$\mathbb{Z}\_{\mathbb{Z}}(2)=0.22\$

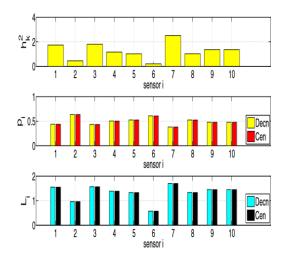


Fig. 2. Centralized and decentralized sensor transmit power and channel bit allocation for  $\mathbb{Z}_{@@} = 0\mathbb{Z}1$ ,  $\mathbb{Z}_{@} = 5$ ,  $\mathbb{Z} = 3$ ,  $\mathbb{Z}_{@} = -1$  dB,  $\mathbb{Z} = 50$  and  $\mathbb{Z}_{@}(\mathbb{Z}) = 0\mathbb{Z}3 \ \forall \mathbb{Z}$ .

a low power budget ( $\mathbb{Z}_{\mathbb{Z}}$ ). Fig. 4 shows the receiver operating characteristic against the sample number ( $\mathbb{Z}$ ). As expected, the matched filter detector outperforms the energy detector but it requires full knowledge of the useful signal. And in Fig. 5, we examine the probability of detection ( $\mathbb{Z}_{\mathbb{Z}}$ ) performance against the total power budget ( $\mathbb{Z}_{\mathbb{Z}}$ ). As  $\mathbb{Z}_{\mathbb{Z}}$  increases, then  $\mathbb{Z}_{\mathbb{Z}}$  improves.

### V. CONCLUSION

We have shown how to perform distributed detection, via SNs transmitting a quantized version of the received energy test statistic to the FC. In addition we have derived the optimal linear combining weights at the FC and proposed a novel distributed algorithm to calculate the optimal transmit power for each SN in order to maximize  $\mathbb{Z}_{\mathbb{Z}}$ . In this way, the SN can allocate its own transmit power by exchanging information

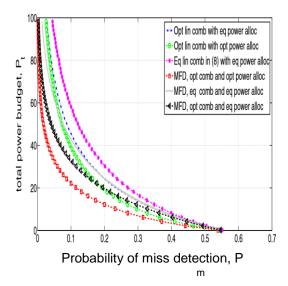


Fig. 3. Total power budget ( $\mathbb{Z}_{\mathbb{Z}}$ ) versus probability of miss-detection ( $1-\mathbb{Z}_{\mathbb{Z}}$ ), with  $\mathbb{Z}_{\mathbb{Z}} = 0\mathbb{Z}$ 1,  $\mathbb{Z} = 3$ ,  $\mathbb{Z}_{\mathbb{Z}} = -4$  dB,  $\mathbb{Z} = 5$  and  $\mathbb{Z} = 100$ .

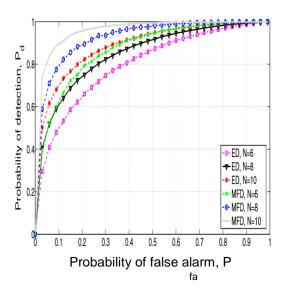


Fig. 4. Probability of detection ( $\mathbb{Z}_{\mathbb{Z}}$ ) versus probability of false alarm ( $\mathbb{Z}_{\mathbb{Z}}$ ), with  $\mathbb{Z}=3$ ,  $\mathbb{Z}_{\mathbb{Z}}=-4$  dB,  $\mathbb{Z}_{\mathbb{Z}}=1$  and  $\mathbb{Z}=10$ .

with its own neighbours. What makes this scheme very useful and attractive is that the only value that they should exchange among neighbours is their own transmit power at the current state. The algorithm is robust and easy implementable.

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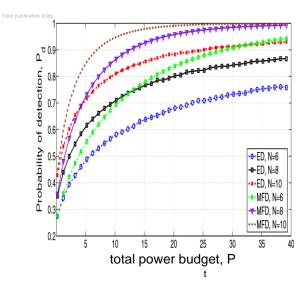


Fig. 5. Probability of detection  $(\mathbb{Z}_{\mathbb{Z}})$  versus total power budget  $(\mathbb{Z}_{\mathbb{Z}})$ , with  $\mathbb{Z}=3$ ,  $\mathbb{Z}_{\mathbb{Z}}=-4$  dB,  $\mathbb{Z}_{\mathbb{Z}\mathbb{Z}}=0\mathbb{Z}1$  and  $\mathbb{Z}=20$ .

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