

Practical Experiments to Evaluate Quality Metrics of MRAM-Based Physical Unclonable Functions

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ABSTRACT: Process variations in the manufacturing of digital circuits can be leveraged to design Physical Unclonable Functions (PUFs) extensively employed in hardware-based security. Different PUFs based on Magnetic Random-Access-Memory (MRAM) devices have been studied and proposed. However, most of such research has been simulation-based, which do not fully capture the physical reality. We present experimental results on a PUF implemented on dies fabricated with a type of the MRAM technology namely Thermally-Assisted-Switching MRAM (TAS-MRAM). To the best of our knowledge, this is the first experimental validation of a TAS-MRAM-based PUF. We demonstrate how voltage values used for writing in the TAS-MRAM cells can make stochastic behaviors required for PUF design. The analysis of the obtained results provides some preliminary findings on the practical application of TAS-MRAM-based PUFs in authentication protocols. In addition, the results show that for key-generation protocols one of the standard error correction methods should be employed if the proposed PUF is used.

INDEX TERMS: Physical Unclonable Functions, magnetic RAM, Thermally Assisted Switching MRAM

I. INTRODUCTION

Physical unclonable functions (PUFs) are primitive and essential circuitry components for hardware-based security. They are cost-effective to generate trustworthy signatures and high quality random numbers for different security protocols such as anti-counterfeiting, identification, authentication, key-generation, etc. [1-2]. Also, PUFs are low-power so they have received significant attention for addressing security issues of the Internet-of-Things (IoT) devices [2-3].

The word *unclonable* in the term *PUF* refers to the most important property of such circuits. It suggests the difficulty of fabricating a circuit or developing an algorithm able to generate the same inputs/outputs of PUF. For this purpose, designers can leverage process variations (PVs) that happen during circuit manufacturing. PVs can make electrical characteristics of manufactured circuits unique in each instance [1].

Different PUFs have been reviewed in previous studies [4]. A well-known class of PUFs in digital circuits/systems is

Memory-PUF in which memory cells are employed to design and implement PUFs. This class has received much interest because memories and state elements are embedded in today's System on Chips (SoCs) [1, 3]. Memory-PUFs can be realized using different memory technologies. Two widespread memory technologies that have been largely studied for PUF design are Static RAM (SRAM) and Dynamic RAM (DRAM). However, PUFs based on these technologies have security issues since they are based on using digital information of "0" or "1" [7]. In such PUFs, for the first response, "0" or "1" is read in some specific cells. The response for those cells is the same for the next response queries. As a result, such PUFs can be broken when the hacker accesses them [7]. Contrariwise, PUFs based on technologies such as Resistive RAM (ReRAM) and Magnetoresistive RAM (MRAM) use the resistance value of memory cells. Thus, a given cell can be either '0', '1', 'Z' depending on its resistance value and the resistance ranges being considered as '0', '1', or 'Z' [7].

The nonobviousness of ReRAM or MRAM PUFs compared with SRAM or DRAM PUF makes them more encouraging [7]. However, ReRAM cells are susceptible to environmental and voltage fluctuations, which results in a 5-20% error rate [8].

MRAM is one of several new emerging technologies aiming to become a “universal” memory device. This technology has the potential to gain importance for non-volatile memories because of promising properties such as non-volatility, low fabrication cost, high speed, low power consumption, and high reliability [10].

In MRAM devices the electron spin is used to store information. Different types of MRAM devices have been realized, such as Toggle MRAM (T-MRAM), Spin-Transfer Torque MRAM (STT-MRAM), Thermally-Assisted-Switching MRAM (TAS-MRAM), etc. In each type, a unique method is used to change the electron spin.

Different PUFs based on different MRAM types have been studied and proposed [9-12]. However, most of the studied MRAM-based PUFs have been conducted on simulation environments. Despite the valuable knowledge obtained in such studies, the lack of practical experiments is very tangible. However, there are a few experimental studies but they are done on STT-MRAM devices. Research on designing other types of MRAM-PUFs is limited and to the best of our knowledge, this is the first experimental validation of TAS-MRAM-based PUF.

In this work, we present experiments, results, and analyses on the stochastic switching behavior of an MRAM-based PUF using some fabricated TAS-MRAM dies. These dies are designed and fabricated by Crocus Technology [13]. The measured quality metrics and analysis results in this study provide some preliminary findings showing that TAS-MRAM can be used in authentication or key-generation protocols instead of SRAM PUFs. This is due to 1) higher speed and lower power of TAS-MRAM rather than of SRAM, 2) the acceptable error rate of the proposed PUF, which is more or less like the SRAM-PUFs error rate. The error rate results in this study show that the proposed PUF like SRAM-PUFs can be used in key generation protocols if one of the common error correction methods is used in that protocol. This is because even a 1-bit mismatch between two keys is not acceptable.

The remainder of the paper is organized as follows: Section II presents a necessary background on PUFs and MRAM devices. Section III represents the proposed PUF design in this work. Section IV explains the experiments and exhibits obtained results. Some discussions are prepared in Section V about different MRAM-PUFs. Finally, Section VI provides concluding remarks.

II. BACKGROUND

A. PHYSICAL UNCLONABLE FUNCTIONS

As mentioned, in order to design PUFs designers can leverage PVs that happen during circuit manufacturing [1-3]. PVs cause fluctuations in physical dimensions of transistors and

interconnections, and thus in their magnetic and electrical characteristics in each instance of a fabricated circuit layout [1].

PUFs can be categorized into two classes: weak PUFs and strong PUFs [1, 9]. Weak PUFs have one or few practical and valid inputs, the so-called *challenge* in literature [1, 9]. On the contrary, strong PUFs have a large number of challenges within a finite timeframe [1, 9]. Weak PUFs employed within IoT devices have at least one unclonable output, the so-called *response*, which is unique per each fabricated device (PUF instance) while applying an identical challenge [2-3]. Therefore, they can be utilized as the device ID, secret authentication signature, or secret cryptographic key [14]. As a result, there is no need for storing the keys, which can be stolen, into untrustworthy Flash or One-Time-Programmable memories. A large number of challenge-response pairs (CRPs) in strong PUFs have enough entropy to be able to replace Hardware Security Modules in Public Key Infrastructure (PKI) [9, 15]

In order to introduce some properties of PUFs in this section and Section IV, we employ a notation as follows:

- \mathcal{P} : the structure or layout of a PUF
- Π : a fabricated instance of \mathcal{P} ;
- $|\Pi|$: the number of all the fabricated Π s
- $C \leftarrow \{0,1\}^N$: a valid challenge for each Π with N bits
- \check{C} : the set of all valid challenges
- $|\check{C}|$: the number of challenges in \check{C}
- $R \leftarrow \{0,1\}^M$: response for each Π with M bits
- $R_{ij} = \Pi_i(C_j)$ ($1 \leq i \leq |\Pi|$, $1 \leq j \leq |\check{C}|$): a response obtained by applying C_j , the j th C of \check{C} , to Π_i , the i th Π .

The quality of a PUF is measured according to different parameters. Three most important ones are explained in the following:

1) INTER-PUF VARIATION, OR UNIQUENESS (UQ)

For a weak PUF, UQ shows the average difference between the responses corresponding to one identical valid challenge applied to every instance of the PUF. For a strong PUF, all or several valid challenges are applied to every instance of the PUF. This parameter is shown in (1). For the ideal PUF, UQ is 50% [1, 9].

$$UQ = \frac{1}{|\Pi| \cdot |\check{C}|} \sum_{i=1}^{|\Pi|-1} \sum_{j=1}^{|\check{C}|} \Pi_i(C_j) - \Pi_{i+1}(C_j) \quad (1)$$

2) INTRA-PUF VARIATION, OR REPRODUCIBILITY (RE)

RE is the variation of the responses of an instance PUF over time. For a weak PUF, RE shows the average difference between responses corresponding to one identical challenge repeatedly applied to each instance of a PUF, e.g., for $T = 1000$ times. For a strong PUF, instead of one challenge, all or several challenges are applied to all the instances. RE is shown in (2). If a PUF is ideal, its responses are stable and robust for

environmental conditions and multiple reading; thus RE is 0% for the ideal PUF [1, 9]. In real words, less than 10% RE is acceptable [9], and one can employ error correction methods to decrease RE to the order of magnitude of -6 or even -9 [16].

$$RE = \frac{1}{T \cdot |\Pi| \cdot |\check{C}|} \sum_{i=1}^{|\Pi|} \sum_{j=1}^{|\check{C}|} \sum_{t=1}^{T-1} T_t(\Pi_i(C_j)) - T_{t+1}(\Pi_i(C_j)) \quad (2)$$

3) UNIFORMITY (UF)

One can easily calculate the fraction exists in the number of '0' to the number '1' in the responses corresponding to an identical challenge applied to every instance of a weak PUF. UF is the average of this fraction among all the instances. Similar to the previous parameters, all or several challenges are utilized for strong PUFs. UF is shown in (3). For the ideal PUF, UF is 1 [1, 9].

$$UF = \frac{1}{|\Pi| \cdot |\check{C}|} \sum_{i=1}^{|\Pi|} \sum_{j=1}^{|\check{C}|} \frac{\text{No. of zeroes in } r_i^{c_j}}{\text{No. of ones in } r_i^{c_j}} \quad (3)$$

These parameters are employed in this work to show the quality of the proposed PUF.

B. MRAM TECHNOLOGY AND TAS-MRAM DEVICES

MRAM exploits the Tunnel magnetoresistance (TMR) effects due to Magnetic tunnel junctions (MTJs). In Fig. 1a, a simplified schematic of an MTJ is presented. It consists of one insulator, the so-called *tunnel barrier*, interpolated between two ferromagnetic layers. The tunnel barrier is thin enough such that electrons can tunnel from one ferromagnetic layer into the other one. The TMR phenomenon is impossible in classical physics, but it is explicable in quantum physics.

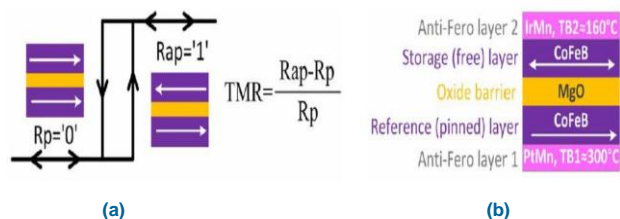


FIGURE 1. a) An MTJ in the parallel and antiparallel state, b) general schematic of an MTJ in TAS-MRAM devices [18]

Another point about MTJs is that their resistances significantly differ when their ferromagnetic layers have a parallel (P) or anti-parallel (AP) magnetic orientation [17]. These two resistances are denoted by R_p and R_{ap} in Fig. 1a. An MTJ resistance is low when the two ferromagnetic layers have a parallel magnetic orientation; otherwise, in the antiparallel configuration, the resistance is high [18]. This can be interpreted that MTJs operate like a switch.

One of the two ferromagnetic layers in each MTJ has a fixed magnetic orientation. This layer is called *pinned layer*

(PL). On the contrary, the magnetic orientation of the other layer, the so-called *free-layer* (FL), is easily changeable. Different methods have been designed to change the magnetic orientation of FL. For example, in STT-MRAM devices, a polarized current is applied to change the magnetic orientation of FL [19].

In TAS-MRAM devices, increasing temperature through MTJ while applying an external magnetic field can change the magnetic orientation of FL [20]. Fig. 1b shows the schematic of an MTJ device fabricated with TAS-MRAM. As can be seen in this figure, the MTJ has two antiferromagnetic layers of AFML₁ and AFML₂, correspond to the PL and FL, respectively. AFML₁ and AFML₂ have the blocking temperatures of TB₁ = 300°C and TB₂ = 160°C, respectively. The magnetic orientation of the FL and PL remains fixed and insensitive to external magnetic fields for temperatures below the blocking temperatures, also called the “temperature of Neel”. To change the magnetic orientation of the FL in Fig. 1b, one must heat AFML₂ up to 160°C. For this purpose, a parameter called “heating voltage” is employed [21].

The writing process in a TAS-MRAM MTJ is depicted in Fig. 2. At first, the MTJ heats up and when it is warm enough, its FL becomes ready to store ‘0’ or ‘1’. Afterward, a field line current is applied to the MTJ. The direction of this current determines the direction of the FL. Finally, a short time-space is required to cool down the MTJ by stopping the heating while maintaining the field line current.

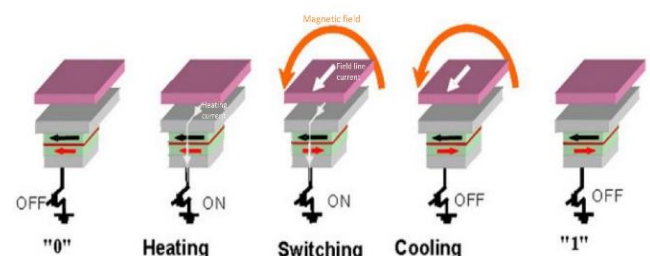


FIGURE 2. The procedure of writing in a TAS-MRAM MTJ [17]

C. PROCESS VARIATION EFFECTS ON TAS-MRAM

Due to PVs physical attributes of TAS-MRAM MTJs, such as the thickness and materials density of the ferromagnetic/insulator layers, are slightly different in every MTJ. One important feature of each MTJ affected by PV is the heating voltage threshold (HVth) of the MTJ, which is the minimum voltage required to sufficient heat up the AFML₂ to change the magnetic direction of the FL [22]. To set/reset a TAS-MRAM MTJ reliably, one needs to employ a value for the heating voltage more than the HVth.

The HVth variability causes stochastic switching behavior in each MTJ while applying a voltage equal to the theoretically calculated HVth. In other words, using the HVth, some MTJs heat up enough to switch and some MTJs do not. As a result, employing a value near to the HVth may cause failures in the set/reset operations [21]. Finding a value that causes a failure

probability of 50% allows designing PUF. However, having exactly 50% is not practically feasible.

III. PROPOSED PUF DESIGN

In this work, we aimed to explore an MRAM-PUF implemented on TAS-MRAM dies. It leverages PV effects on write operations. The dies are designed and fabricated by Crocus Technology [13]. Each die includes 1 *kbits* arranged in a 32×32 array such that each bit is individually addressable and accessible. A microscopic picture of the dies and its holding package (QFN44) are shown in Fig. 3.

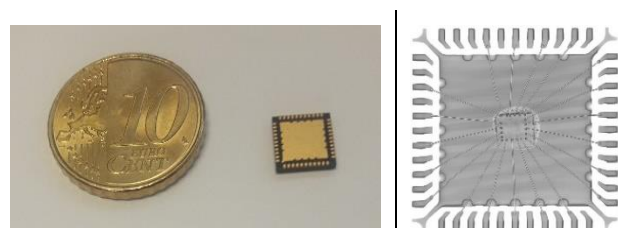


FIGURE 3. QFN44 package (Left) and microscopic picture of the TAS-MRAM die in QFN44 (right)

The architecture of the dies is presented in Fig. 4. Here, IOF, IOM, and IOR are sense pads employed during the read (R) operation. As seen in this figure, IOF is the pad connected to the top of the MTJ of each memory cell. It enables applying the intended voltages, during the read/write operations. IOM is connected right below the MTJ, and IOR is between a poly 500Ω resistance and a Select transistor. This transistor in each cell is driven by one of the outputs of the decoder, which converts the data of the address pines. The total impedance of an MTJ holding '0' is changed from $R_{min} = 2k\Omega$ to $R_{max} = 4k\Omega$ by the write '1' operation (W_1). Likewise, the write '0' operation (W_0) changes the MTJ impedance from R_{max} to R_{min} . Both the operations require three voltages: V_{Heat} , V_{Field1} , and V_{Field2} . The first one is needed to heat locally the selected MTJ, whereas the second and third ones allow changing the magnetic orientation of FL in the desired state after heating. To have a certain W_0 , one needs to apply 2.2 V to V_{Heat} and then V_{Field1} and V_{Field2} 3.3 V and 0 V, respectively. Likewise, a certain W_1 needs to apply 3.3 V, 0 V, and 2.2 V to V_{Heat} , V_{Field1} , and V_{Field2} , respectively. The duration of these three signals, T_{Heat} , T_{Field1} , and T_{Field2} must be 30 ns. In these cases, one can be sure that the write operations are done without any failure. We call these operations "Certain Write 0" (CW_0) and "Certain Write 1" (CW_1). The requirement of the read (R) operation are $V_{Field1} = V_{Field2} = 0$ V, $V_{Heat} = 0.3$ V, and $T_{Heat} = 30$ ns.

In the proposed PUF, we use CW_1 and then the "Uncertain Write 0" (UW_0) in which a voltage in the range of 1 V to 1.8 V is selected for applying to V_{Heat} . In this case, the probability that a cell has the logic '0' after performing the sequence of CW_1 - UW_0 depends on: 1) the selected voltage for V_{Heat} , 2) PVs that affect the MTJ. To have an efficient PUF, one must select a value from the mentioned range such that it results in a failure probability of 50%.

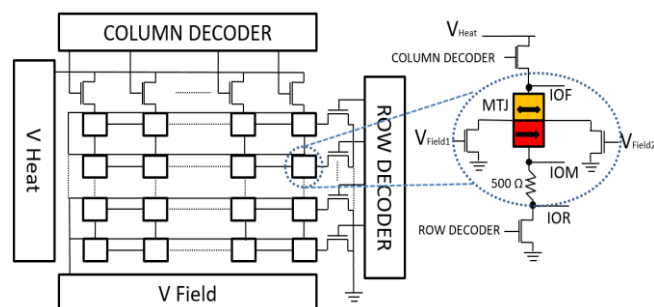


FIGURE 4. The architecture of the employed TAS-MRAM dies, and a memory cell including an MTJ and select transistor

It is noteworthy that if the CW_0 and CW_1 operation does not affect the resistance value of an MTJ, it can be broken or defective. Broken MTJs are the ones whose resistance and voltage value are sensibly lower than normal MTJs. The resistance and voltage of broken MTJs in our dies in either of the P or AP states are always less than 400 Ω and 100 mV, respectively. There are also some defective MTJs, not broken or shorted; they are always in either of the P or AP states. In fact, the magnetic direction of the FL is always fixed in the defective MTJs. Therefore, the resistance of an MTJ does not change even if the CW_0 / CW_1 operations repeated several times, that MTJ is defective.

IV. EXPERIMENTS AND RESULTS

In our experiments, we employ a Xilinx NEXYS2 FPGA board [23] and a National Instrument acquisition board [24] for writing in and reading from the dies, shown in Fig. 5. As seen in this figure, a PCB is designed and fabricated to route the address and read/write pins of the die to the FPGA board. It includes 3 digital potentiometers that can be controlled and set up through I²C protocol. The FPGA sets up addresses and commands these potentiometers to regulate the necessary voltages for writing in or reading from the cells of a die under test. The acquisition board is employed to read the voltage value of the addressed cells during the reading process. Three pads are shown in Fig. 4 and used to read the voltage of the MTJ at the three different points.

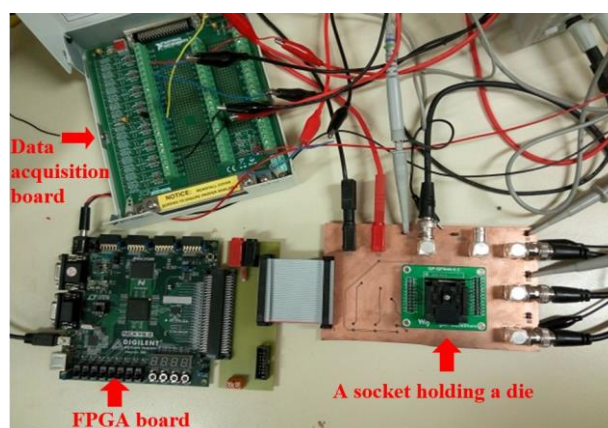


FIGURE 5. Equipment employed in the experiments

The flow of the experiments starts by distinguishing broken and defective cells from intact ones. Figs. 6a and 6b illustrate the cells of a die by applying CW_1 and CW_0 , respectively. As seen in Fig. 6a, there is one broken cell (blue) in the die under test. Six always-P-state cells (yellow) with a voltage value between 220 mV and 260 mV are also seen in Fig. 6a. Moreover, 21 always-AP-state cells (dark red) with a value of more than 300 mV are seen from Fig. 6b. In Fig. 6b there are some (around 45) light red cells with a voltage between 270 mV and 290 mV . Some of them are fixed, but some of them usually get a value in this range. One can consider wider voltage ranges to interpret the P and AP states and use such cells.

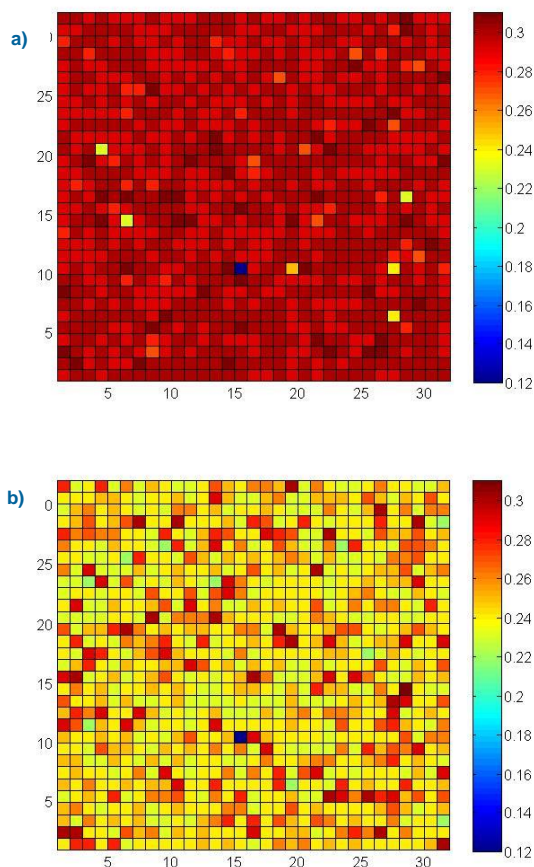


FIGURE 6. 32x32 cells in a die, a) applied CW_1 and b) applied CW_0

The probability of the voltage range corresponding to Fig. 6 is shown in Fig. 7. The common part between the two diagrams of this figure is correspondent to the light red cells in Fig. 6. These two diagrams are almost the same as all the other dies. The values in this diagram indicate that there is more variation in the P state. Therefore, in the next step, we run the sequence of CW_1 - UW_0 .

The second step in our experiments is to analyze different heating voltages to find out a proper one, by which 50% Hamming distance is obtained for inter-PUF variation. For this purpose, experiments begin from a heating voltage equal to 1.8 V. The sequence of CW_1 - UW_0 -R is applied to all the cells of

26 dies. Then, 0.1 V is deduced from the heating voltage and the sequence is repeated. This procedure is continued until the heating voltage reaches 1 V.

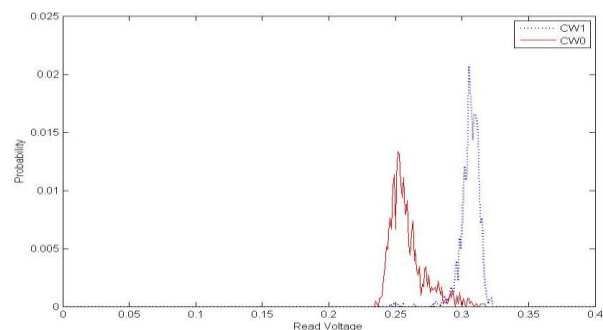


FIGURE 7. Probability of the voltage range corresponding to Fig. 6

Figure 8 illustrates the inter-PUF variation obtained by (1) for each of 26 dies in different heating voltages. The dashed lines in this figure show that for the heating voltages in the range of 1.4 V-1.6 V. For these three heating voltages the inter-PUF variations for all the dies are almost 40%-50%, which is an acceptable range. Fig. 9 shows the average of the inter-PUF variations among the all dies for the employed heating voltages. According to the green (dashed) bar in the figure, the best inter-PUF variation result is 49.8% for 1.4 V.

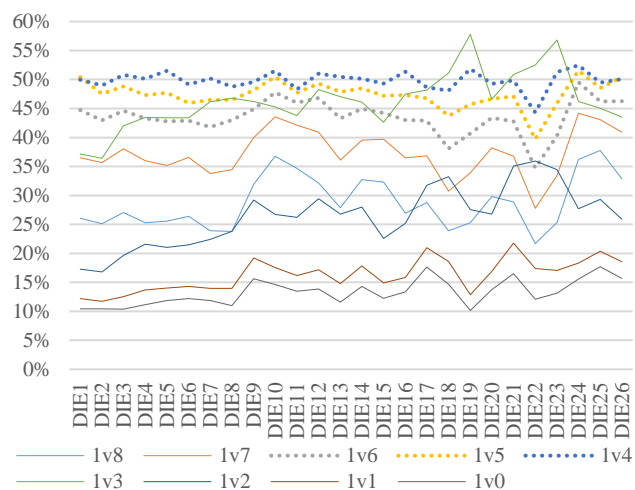


FIGURE 8. Inter-PUF variations corresponding to each die for different heating voltages

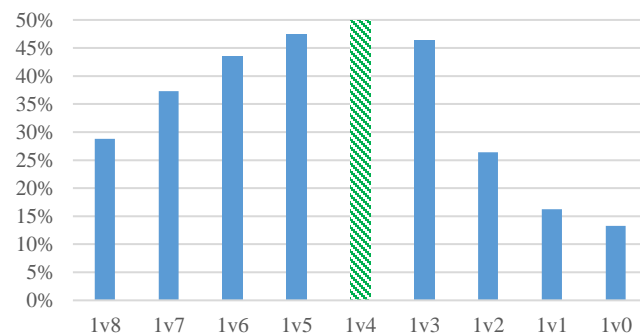


FIGURE 9. Average of the inter-PUF variations among all the dies for different heating voltages

The third step in our experiments is to measure intra-PUF variation. For this purpose, the sequence of CW_1-UW_0-R was run 1000 times on all the cells of the 26 dies. Like the previous step, nine heating voltages varying from 1.8 V to 1 V were analyzed. The results are reported in Fig. 10 and Fig. 11. Figure 10 shows the intra-PUF variations corresponding to each die for different heating voltages. Fig. 11 presents the average of the intra-PUF variations among all the dies for different heating voltages. In these figures, the lines and bars corresponding to the heating voltage 1.4 V are dashed (green). For this voltage, the intra-PUF variation obtained by (2) is 7.7%. This rate is acceptable and can be easily corrected using error correction methods like those proposed in [11, 16].

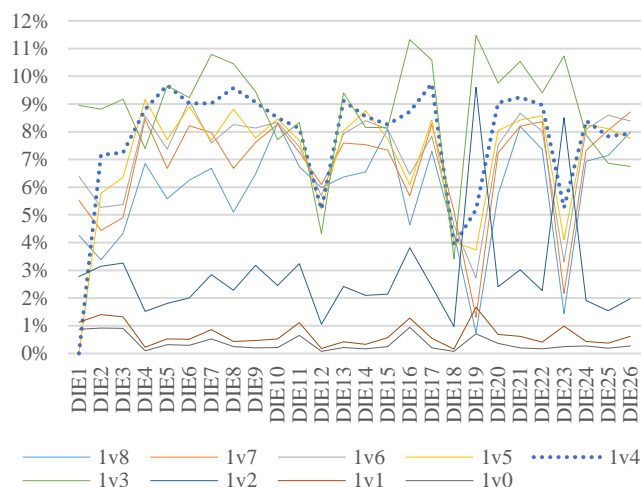


FIGURE 10. Intra-PUF variation corresponding to each die for different heating voltages

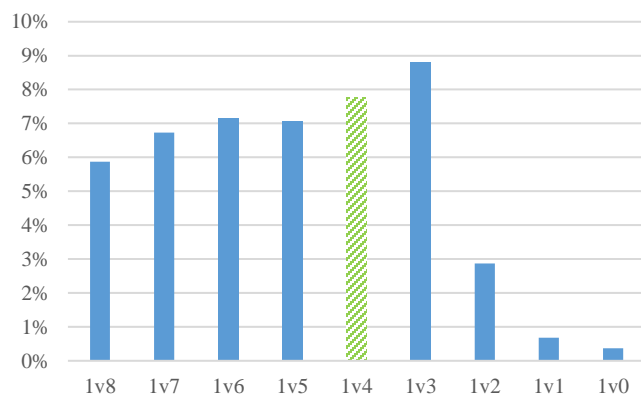


FIGURE 11. Average of the intra-PUF variations among all the dies for different heating voltages

As mentioned in Section II, uniformity is an important parameter for PUFs. Fig. 12 shows the average of the uniformity among all the dies for the different heating voltages. As shown in the figure, the best result (i.e., 0.93) is obtained using the heating voltage of 1.4 V.

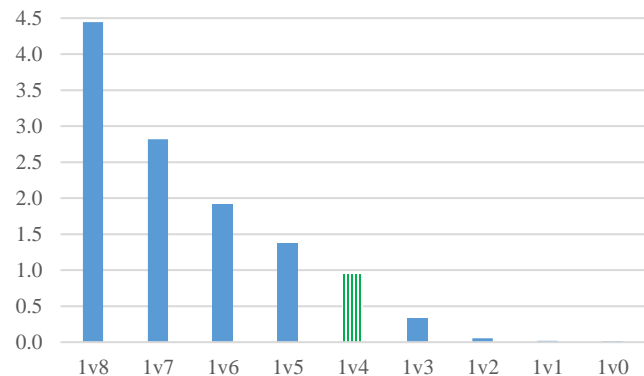


FIGURE 12. Average of the uniformity among all the dies for the different heating voltages

V. DISCUSSION

Although MRAM devices have many advantages over other types of NVMs, there are still existing issues regarding their reliability and security. One of the main concerns about MRAMs is their vulnerability to magnetic fields and temperature. Different methods have been used to solve these issues. For example, in T-MRAM, designing a magnetic shield or heat shield around the device can be a possible solution [31-32], but due to its cost it may not be practical or useful for applications such as IoTs. Moreover, MRAM devices can face an intentional magnetic field and temperature, which may not be protected by the packaging. For example, Everspin lists the maximum magnetic tolerance for their MRAM chips during write/read/standby to be only 1000e [25].

STT-MRAM offers high density, and consumes almost zero leakage power. This MRAM type is quite mature [26]. Unique characteristics of STT-MRAM make it a suitable candidate for several applications such as replacing conventional memories. However, its high and asymmetric read/write current introduces new challenges. For example, similar to the write current, the read current passes through the MTJ, applying a spin-transfer torque effect on the storage-(free-)layer magnetization. Therefore, one of the problems with STT-MRAM is the induced disturbance of the storage layer magnetic state during reading [9]. More importantly, different kinds of security attacks have been reported on STT-MRAM-PUFs. For example, in [10], a Correlation Power Analysis on the write operations of an STT-MRAM-PUF was performed successfully.

TAS-MRAM requires a single magnetic field and lower field values compared to T-MRAM, and thus it consumes less power than T-MRAM. Also, TAS-MRAM solves the limitation of conventional MRAM devices for PUF applications. First, addressing errors are reduced since the selection at write operation is driven by temperature. Second, TAS-MRAM has better reliability and less vulnerability than STT-MRAM to field disturbance [27]. Although external fields change the resistance state of a cell, the state after the field disturbance goes back to its first state [9, 27].

The results of some of the previous publications on MRAM-PUFs are provided in Table I. The studies were selected since the quality metrics of their designed MRAM-PUFs have also been reported without any post-processing method. However, they are not necessarily the most recent. In [28-30], promising MRAM-based PUF results are reported. However, the reported results are based on simulation studies, which do not fully capture the physical reality. In comparison, our results are based on experiments on the fabricated TAS-MRAM dies. In [10], the results are based on small-scale experiments on STT-MRAM while our results are based on larger-scale analyses. For using the MRAM-based PUFs in key generation protocols [28-30], it is necessary to employ some error reduction/correction methods for enhancing the PUF reliability. Our results show 49.8% of the inter-PUF between different TAS-MRAM-based PUF dies. Nevertheless, our findings show 7.7% of the intra-PUF, which is more satisfactory compared to one reported in the previous studies [28-30]. Therefore, a method for handling the error is necessary for key generation using the TAS-MRAM-based PUF. However, our results show that the proposed PUF can be used instead of SRAM PUFs for key-generation or authentication protocols, which are considered as future research. This is because the inter-PUF distances are near to the maximum and the intra-PUF error rate is in an acceptable range. This can reduce False Rejection Rate (FRR) and False Acceptance Rate (FAR) in authentication protocols. Numerical evaluation of FAR and FRR in key-generation protocols is considered as future work.

TABLE I
BASED PUF DESIGNS COMPARISON

MRAM Type	RE	UQ	Ref. #	Simulation Experimental
STT	5.0%	45.83%	[28]	Simulation
STT	2.25%	47 %	[29]	Simulation
STT	1%~10%	~50.1%	[30]	Simulation
STT	0%	~50%	[10]	Experimental
TAS	7.7%	49.8%	This work	Experimental

VI. CONCLUSION

In this work, practical experiments were performed on fabricated TAS-MRAM dies. The experiments included analyzing the efficiency of an implemented PUF on these dies. The major objective of this work was to investigate the stochastic switching behavior in the MTJ cells of the dies as a source of randomness. The experimental results of this work showed the possibility of obtaining a PUF on TAS-MRAM devices with almost 50% inter-PUF variation, 0.94 uniformity by accepting 7.7% intra-PUF variation, without using any extra hardware overhead. To realize this goal, we need to run on 1 kbits the sequence of the certain-writing-1 and uncertain-writing-0 operations by applying 2.2 V and 1.4 V for the heating voltage required in the writing operations in TAS-MRAM. The main conclusion is that our suggested circuits can be used to design resilient PUFs with TAS-MRAM.

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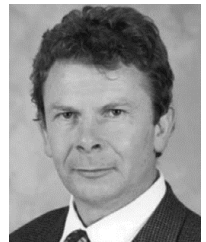
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