Parallel Relational Database Systems

I. Introduction

II. Optimization-Parallelization Strategies (Inter-operation)

III. Efficiency of Parallelism

IV. Optimization of Data Communication

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I. Introduction to Parallel Rel. DB

1. Motivation *[Dew 90, Val 93, Lu 94, ...]*

☛ *Relational Languages: Declaratives*

- *Regular Data Structures : Static Annotation*
- *Relational Language : Declarative*
  - ➡ *Automatic Parallelization*
- *Decision Support Queries : Complex, Huge DB, Join, Sort, Agregation*

2. Objectives:

☛ *Best Cost / Performance with respect to Mainframe DPS8 / GCOS, IBM 30390 / VMS, ...)*

☛ *High Performance:*

- *Minimizes the Response Time*
- *Maximizes the Parallel System Throughput*

☛ *" Scalability " :*

- *Adding New Resources (CPU, Disk, Memory)*
- *Adding New Users*

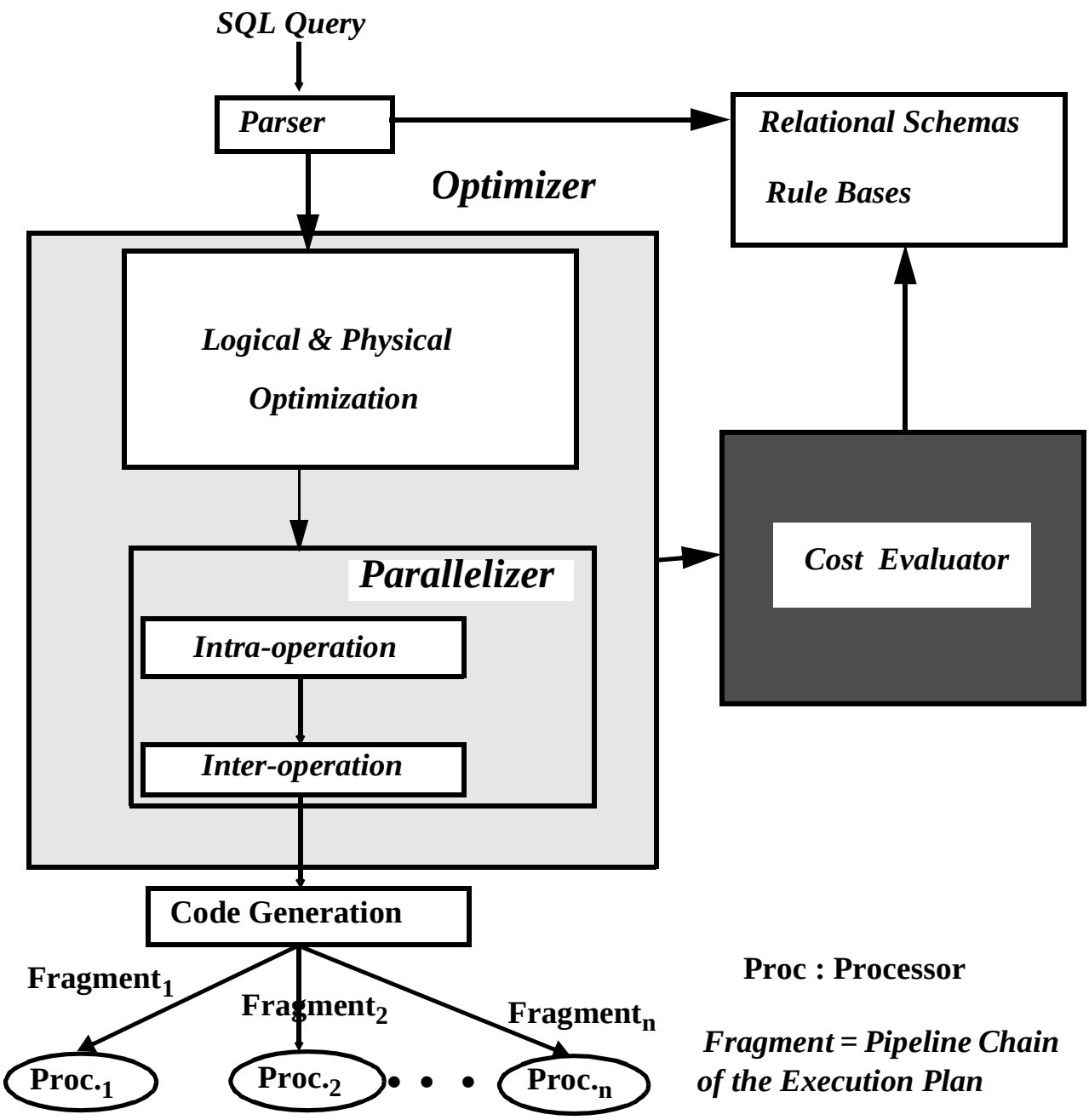
⇨ *Holding the Same Performance*

☛ *Availability*

II. Optimization-Parallelization Strategies (Inter-operation)

1. Introduction

☛ *Query Compiler Architecture for Parallel Database Systems*



SPJ Query Parallelisation:

Parallelism Extraction & Resource Allocation

A. Parallelism Extraction

1. Data Partitioning: Approaches & Methods [Liv 87, Cop 88, Dew 92]

- ***Partitioning Degree of each base relation?***

2. Parallelism degrees of Joins?

3. Parallelization Strategies (Inter-operator)

- ***Approaches***

⇨ ***Two- Phase Approach :*** Φ_1 ; Φ_2

◆ ***XPRS [Hon92, Sto 88], Papyrus [Gan 92, Has 94, Chek 95], Gamma Proj. [Kab 98], ...***

⇨ ***One-Phase Approach: packs*** Φ_1 & Φ_2 : ***into one process***

◆ ***[Sch90, Che92, Zia93, Lan93,...]***

Φ_1 : ***Physical Optimization (without considering the resources)***

Φ_2 : ***Parallelization: Parallelism Extraction & Resource Alloc.***

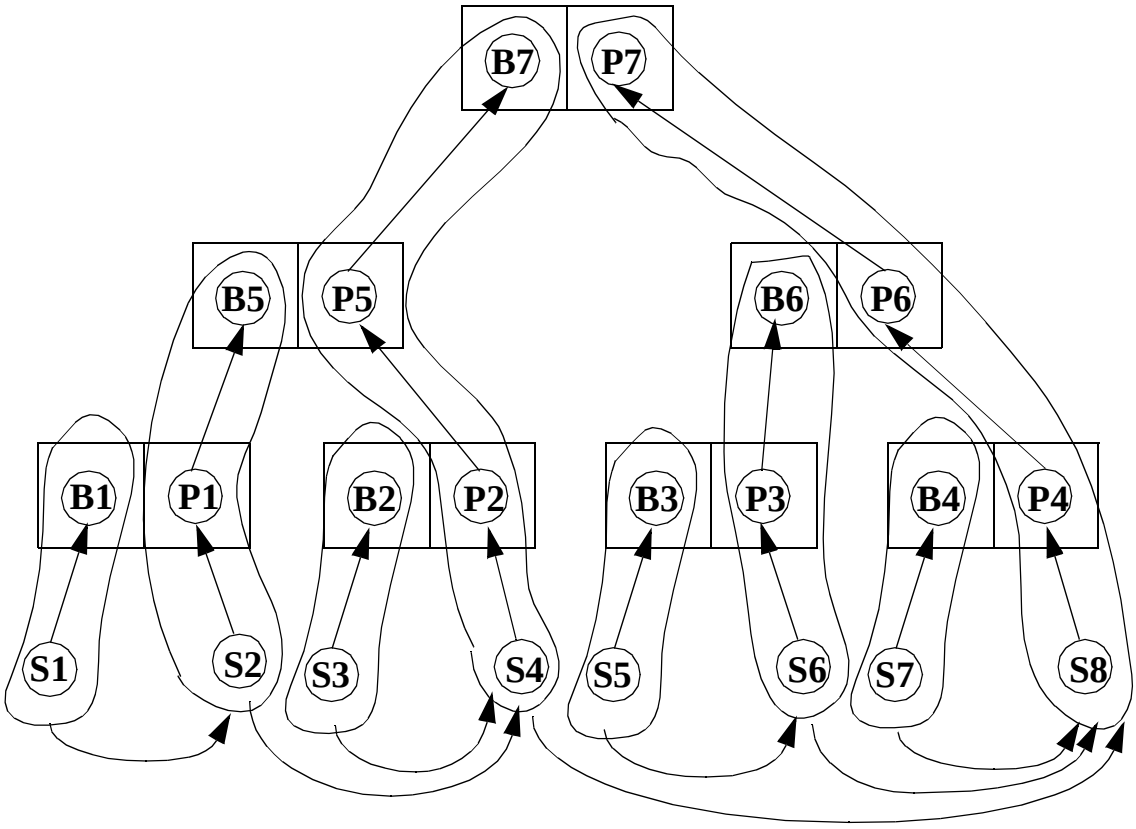
4. Generation of Parallel Programms :

Query = R1xR2xR3xR4xR5xR6xR7xR8

B. Resource Allocation (Mapping)

1. Data (relations) Placement : Alloc_R

2. Tasks (Operator) Placement : Alloc_T

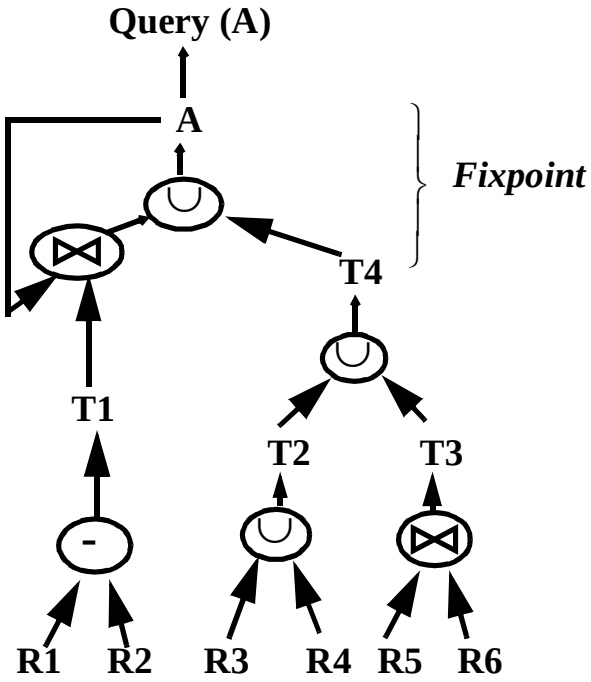



```
SEQ
    PAR
        PIPE Scan S1 - Build J1 ENDPIPE
        PIPE Scan S3 - Build J2 ENDPIPE
        PIPE Scan S5 - Build J3 ENDPIPE
        PIPE Scan S7 - Build J4 ENDPIPE
    ENPAR;
    PAR
        PIPE Scan S2 - Probe J1- Build J5 ENDPIPE
        PIPE Scan S6 - Probe J3 - Build J6 ENDPIPE
    ENDPAR;
    PIPE Scan S4 - Probe J2 - Probe J5 -Build J7 ENDPIPE
    PIPE Scan S8 - Probe J4 - Probe J6 - Probe J7 ENDPIPE
ENSEQ
```

Bushy Tree

III. Efficiency of Parallelism

- ***Shared-Nothing Architecture***
- ***Rel. Size [Bit 83], & Parameters [Sch 90], [Val 88]***



- **Simple Hash-Join Algorithm (Build + Probe)**
  Build (R) holds in memory

LRT $(T \leftarrow R \bowtie S) = T_{ef} + T_d + T_{com}$ where

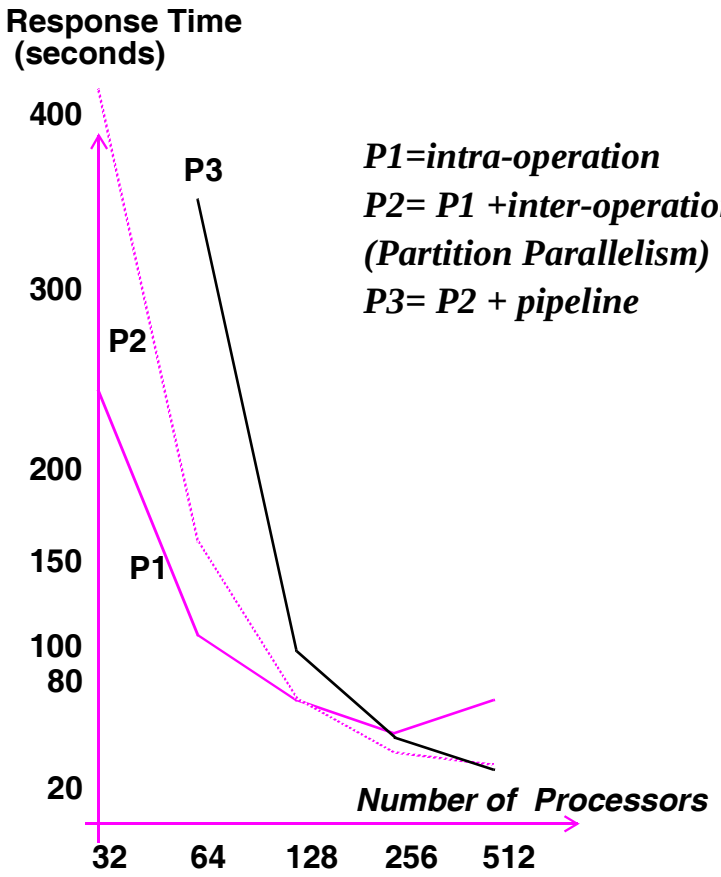
T_{ef} **= (|R|/d).th + ((||R||/d) + (||S||/d)).CR + (||R||/d/q).(||S||/d).CJO + |T|.I + ||T||/d.CW**

Build Time + Read Time + ComparisonTime +
Time for Moving a Tuple + WriteTime

T_d **= (|T|/d).th and** T_{com} **= ((|T|/d).trf + p.msg).** $\lceil d/p \rceil$

|R| : Number of Tuples in R=10^6
||R|| : Numbre of Pages in R
th : Time to hash a tuple (200 B)
CR : Time to read 1 page (18KB)= 8 ms
CW : Time to write 1 page= 16 ms
CJO : Time for joining 2 unsorted pages
d : Number of proc. of source operation
p : Number of proc. of destination operation

Trf : Time to transfer a tuple
msg : Time to process a message
CPU = 4 MIPS



◆ *Efficiency of Parallelism :*

- *Intra-Operation with Lower NB of Processors*
- *Pipeline with Large NB of Processors*

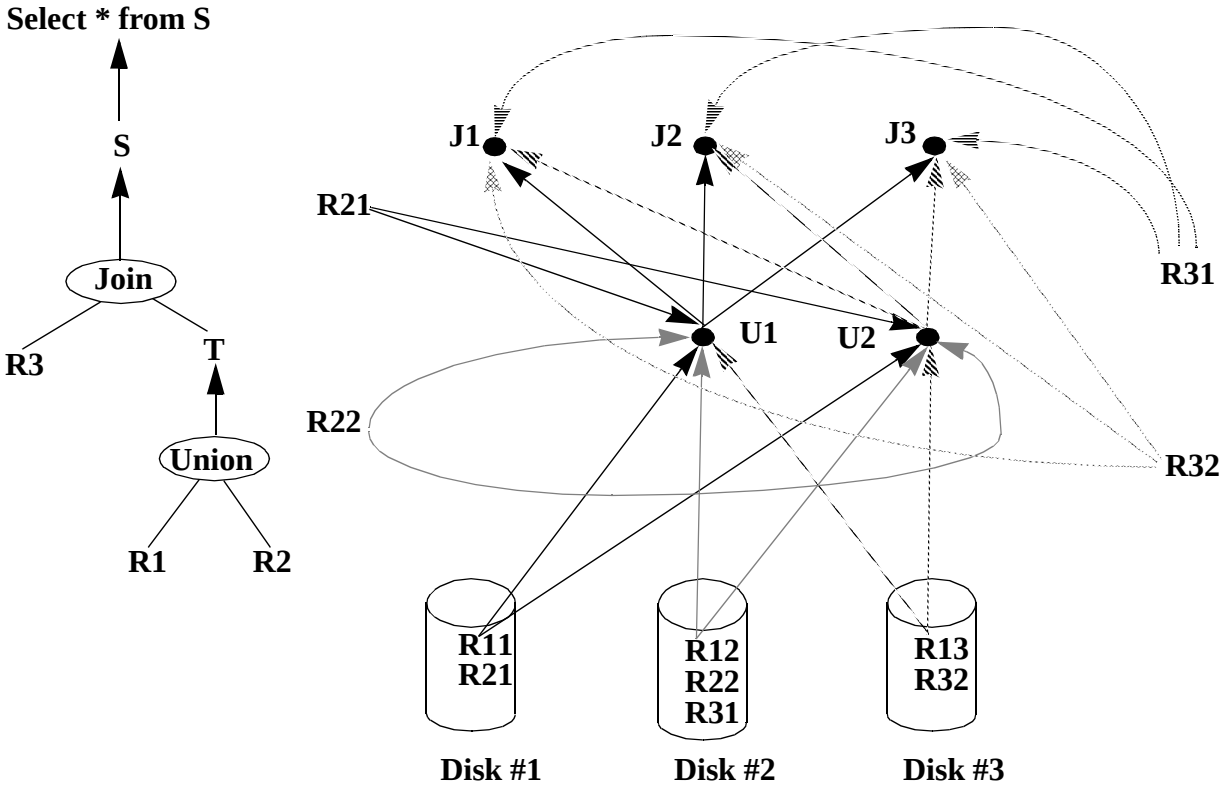
➥ *The Plague of Parallelism : Cost of Data Communication*

IV. Optimization of Data Communication

1. Logical Optimization : JSP -->PSJ (Reducing the Vol. of Data)

2. Physical Optimization : the order in which the joins are executed

3. Parallelization Phase :
Cost of Tuple Redistributing



A Simple SQL Query and Associated Data Flow Graph

Methods : Tree Coloring [HAS 95] Propagation Method[Ham 93]

- ***Partitioning Attributes & /Same Partitioning Function***
- ***Number of Processors***

Propagation Method : Partition Attribute & Number of Processors