Robust Planning with (L)RTDP

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Abstract

Stochastic Shortest Path problems (SSPs), a subclass of Markov Decision Problems (MDPs), can be efficiently dealt with using *Real-Time Dynamic Programming* (RTDP). Yet, MDP models are often uncertain (obtained through statistics or guessing). The usual approach is robust planning: searching for the best policy under the worst model. This paper shows how RTDP can be made robust in the common case where transition probabilities are known to lie in a given interval.

1 Introduction

In decision-theoretic planning, Markov Decision Problems [Bertsekas and Tsitsiklis, 1996] are of major interest when a probabilistic model of the domain is available. A number of algorithms make it possible to find plans (policies) optimizing the expected long-term utility. Yet, optimal policy convergence results all depend on the assumption that the probabilistic model of the domain is accurate.

Unfortunately, a large number of MDP models are based on uncertain probabilities (and rewards). Many rely on statistical models of physical or natural systems, such as plant control or animal behavior analysis. These statistical models are sometimes based on simulations (themselves being mathematical models), observations of a real system or human expertise.

Working with uncertain models first requires answering two closely related questions: 1– how to model this uncertainty, and 2– how to use the resulting model. Existing work shows that uncertainty is sometimes represented as a set of possible models, each assigned a model probability [Munos, 2001]. The simplest example is a set of possible models that are assumed equally probable [Bagnell *et al.*, 2001; Nilim and Ghaoui, 2004]. Rather than construct a possibly infinite set of models we represent model uncertainty by allowing each probability in a single model to lie in an interval [Givan *et al.*, 2000; Hosaka *et al.*, 2001].

Uncertain probabilities have been investigated in resource allocation problems [Munos, 2001] — investigating efficient exploration [Strehl and Littman, 2004] and state aggregation [Givan *et al.*, 2000] — and policy robustness [Bagnell *et al.*, 2001; Hosaka *et al.*, 2001; Nilim and Ghaoui, 2004]. We focus on the later, considering a two-player game where the opponent chooses from the possible models to reduce the long-term utility.

Our principal aim is to develop an efficient planner for a common sub-class of MDPs for which some policies are guaranteed to eventually terminate in a goal state: Stochastic Shortest Path (SSP) problems. The greedy *Real-Time Dynamic Programming* algorithm (RTDP) [Barto *et al.*, 1995] is particularly suitable for SSPs, as it finds good policies quickly and does not require complete exploration of the state space.

This paper shows that RTDP can be made robust. In Section 2, we present SSPs, RTDP and robustness. Then Sec. 3 explains how RTDP can be turned into a robust algorithm. Finally, experiments are presented to analyse the behavior of the algorithm, before a discussion and conclusion.[1]

2 Background

2.1 Stochastic Shortest Path Problems

A Stochastic Shortest Path problem [Bertsekas and Tsitsiklis, 1996] is defined here as a tuple $\langle S, s_0, G, A, T, c\rangle$. It describes a control problem where S is the finite set of **states** of the system, $s_0 \in S$ is a starting state, and $G \subseteq S$ is a set of goal states. A is the finite set of possible **actions** a. Actions control transitions from one state s to another state s' according to the system's probabilistic dynamics, described by the **transition function** T defined as $T(s, a, s') = Pr(s_{t+1} = s'|s_t = s, a_t = a)$. The aim is to optimize a performance measure based on the **cost function** $c : S \times A \times S \to \mathbb{R}^+$.[2]

SSPs assume a goal state is reachable from any state in S, at least for the optimal policy, so that one cannot get stuck in a looping subset of states. An algorithm solving an SSP has to find a **policy** that maps states to probability distributions over actions $\pi : S \to \Pi(A)$ which optimizes the **long-term cost** J defined as the expected sum of *costs* to a goal state.

In this paper, we only consider SSPs for planning purposes, with full knowledge of tuple defining the problem: $\langle S, s_0, G, A, T, c\rangle$. In this framework, well-known stochastic dynamic programming algorithms such as *Value Iteration* (VI) make it possible to find a deterministic policy that corresponds to the minimal expected long-term cost J^*. Value Iteration works by computing the function $J^*(s)$ that gives

[1] Work presented in more details in [Buffet and Aberdeen, 2004].

[2] As the model is not certain, we do not make the usual assumption $c(s, a) = \mathbb{E}_{s'}[c(s, a, s')]$.

the expected sum of costs of the optimal policies. It is the unique solution of the fixed point Bellman equation:

$$J(s) \quad = \quad \min_{a\in A} \sum_{s'\in S} T(s,a,s')\left[c(s,a,s') + J(s')\right]. \quad (1)$$

Updating J with this formula leads to asymptotic convergence to J^*. For convenience, we also introduce the Q-value: $Q(s,a) = \sum_{s'\in S} T(s,a,s')[c(s,a,s') + V(s')]$.

SSPs can easily be viewed as shortest path problems where choosing a path only probabilistically leads you to the expected destination. They can represent a useful sub set of MDPs, as they are essentially finite-horizon MDPs with no discount factor.

2.2 RTDP

Trial based[3] *Real-Time Dynamic Programming* (RTDP), introduced in [Barto *et al.*, 1995], uses the fact that the SSP costs are positive and the additional assumption that every trial will reach a goal state with probability 1. Thus, with a zero initialization of the long-term cost function J, both J and Q-values monotonically increase during their iterative computation.

The idea behind RTDP (Algorithm 1) is to follow paths from the start state s_0, always greedily choosing actions with the lowest long-term cost and updating $Q(s,a)$ as states s are encountered. In other words, the action chosen is the one expected to lead to the lowest future costs, until the iterative computations show that another action may do better.

Algorithm 1 RTDP algorithm for SSPs

```
RTDP($s$:state) // $s = s_0$
repeat
   RTDPTRIAL($s$)
until // no termination condition

RTDPTRIAL($s$:state)
while ¬GOAL($s$) do
   $a$ =GREEDYACTION($s$)
   J($s$)=QVALUE($s, a$)
   $s$ =PICKNEXTSTATE($s, a$)
end while
```

RTDP has the advantage of quickly avoiding plans that lead to high costs. Thus, the exploration looks mainly at a promising subset of the state space. Yet, because it follows paths by simulating the system's dynamics, rare transitions are only rarely taken into account. The use of a simulation makes it possible to get good policies early, but at the expense of a slow convergence, because of the bad update frequency of rare transitions.

2.3 Robust Value Iteration

Pessimism and Optimism — We now turn to the problem of taking the model's uncertainty into account when looking for a "best" policy. The (possibly infinite) set of alternative models is denoted $\mathcal{M}$.

A simplistic approach computes the average model over $\mathcal{M}$, or the most probable model in $\mathcal{M}$, then uses standard SSP optimization methods. Such approaches guarantee nothing about the long-term cost of the policy if the true model differs from the one chosen for optimization.

We follow the approach described in [Bagnell *et al.*, 2001], that consists of finding a policy that behaves well under the worst possible model. This amounts to considering a two-player zero-sum game where a player's gain is its opponent's loss. The player chooses a policy over actions while its "disturber" opponent simultaneously chooses a policy over models (as this is a simultaneous game, optimal policies can be stochastic). This results in a max-min-like algorithm:

$$\max_{\pi_{\mathcal{M}}\in\Pi_{\mathcal{M}}} \min_{\pi_A\in\Pi_A} J_{\pi_{\mathcal{M}},\pi_A}(s_0).$$

In this SSP game, Value iteration does converge to a fixed solution [Patek and Bertsekas, 1999].

It is also possible to be optimistic, considering that both players collaborate (as they endure the same costs), which turns the $\max$ into a $\min$ in previous formula. This second case is equivalent to a classical SSP where a decision consists of choosing an action and a local model.

Locality — Such a max-min algorithm would be particularly expensive to implement. Even restricting the search to a deterministic policy over models, it requires computing the optimal long-term cost function for each model before picking the worst model found and the optimal policy associated with it. However, a simpler process may be used to compute J while looking simultaneously for the worst model. It requires the hypothesis that next-state distributions $T(s,a,\cdot)$ are independent from one state-action pair (s,a) to another. This assumption does not always hold. However, this makes things easier for the opponent because it now has larger set of models from which to choose. This consequence is conservative for producing robust policies.

Because we assume independence at the state-action level (not only at the state level), it is equivalent to a situation where the second player makes a decision depending on the current state and the first player's action. This situation amounts to a *sequential* game where the previous players move is known to the next player: both players can act in a deterministic way without loss of efficiency.

The result of this assumption is that the worst model can be chosen locally when Q is updated for a given state-action pair. As can be seen from Algorithm 2, the worst local model m_s^a may change while Q-values evolve. Previous updates of reachable states' long-term costs may change their relative ordering, changing which outcomes are considered to be worst.

The key contribution of this paper is to show that RTDP can be made *robust*, allowing for planning in very large and uncertain domains, retaining worst (or best) case behaviour guarantees.

3 Robust RTDP

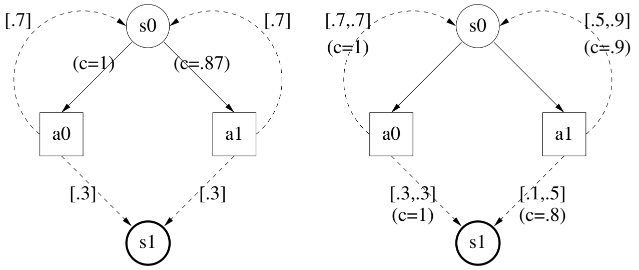
From now on, we consider **interval-based** uncertain SSPs, where $T(s,a,s')$ is known to be in an interval $[Pr^{\min}(s'|s,a), Pr^{\max}(s'|s,a)]$. Figure 1 is an example. We

[3] We always assume a trial based RTDP implementation.

Algorithm 2 Robust Value Iteration (for an SSP)

```
Initialize $J$ to 0.
repeat
    for all $s$: state do
        for all $a$: action do
            $Q_{\max}(s,a) \leftarrow \max_{m_s^a \in \mathcal{M}_s^a} \sum_{s' \in S} T_{m_s^a}(s,a,s') \left[J(s') + c_{m_s^a}(s,a,s')\right]$
        end for
        $J(i) \leftarrow \min_{a \in A} Q_{\max}(s,a)$
    end for
until $J$ converges
```

discuss the pessimistic approach, the optimistic one leading to similar results.



a) certain SSP b) uncertain SSP

Figure 1: Two views of one SSP, depending on whether model uncertainty is taken into account (costs in parenthesis). In the uncertain SSP, action a_0 will be prefered as it quickly reaches the goal s_1.

For a given state-action pair (s,a), there is a list $R = (s'_1, \cdots, s'_k)$ of reachable states. For each reachable state $T(s,a,s'_i) \in I_i = [p_i^{\min}, p_i^{\max}]$. Thus, possible models are the ones that comply with these interval constraints while ensuring $\sum_i T(s,a,s'_i) = 1$. Fig. 2 illustrates this with three reachable states.

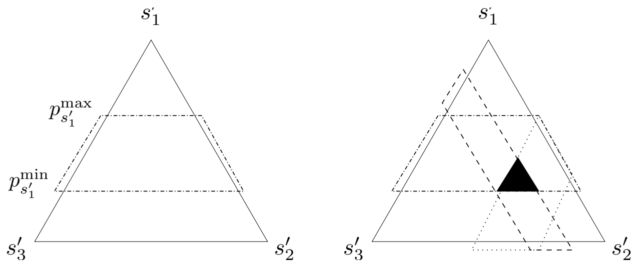


Figure 2: A triangle is a probability simplex representing all possible probability distributions with three different outcomes ($Pr(s'_i) = 1$ at the s'_i vertex). On the left triangle is the trapezium showing the interval constraint for s'_1. The right triangle shows possible models at the intersection of the three interval constraints.

Worst Local Models — The maximisation step to compute $Q(s,a)$ in Alg. 2 is done by giving the highest probability to the worst outcome. This requires firstly sorting reachable states in decreasing order of the values: $c(s,a,s'_1) + J(s'_1) \geq c(s,a,s'_2) + J(s'_2) \geq \cdots c(s,a,s'_k) + J(s'_k)$. Then, the worst distribution is the one giving the highest probability to the first state s'_1, then to s'_2, and so on up to s'_k. As pointed out in [Givan *et al.*, 2000], this is equivalent to finding the highest index $r \in [1..k]$ such that

$$\sum_{i=1}^{r-1} p_i^{\max} + \sum_{i=r}^{k} p_i^{\min} \leq 1.$$

The resulting transition probabilities are

$$Pr(s'_i) = \begin{cases} p_i^{\max} \text{ if } i < r \\ p_i^{\min} \text{ if } i > r \end{cases} \qquad (2)$$

$$Pr(s'_r) = 1 - \sum_{i=1, i \neq r}^{k} Pr(s'_i). \qquad (3)$$

Using the pre-computed bound $B_{\min} = \sum_{i=1}^{k} p_i^{\min}$, Alg. 3 gives a complete implementation. The *insertion sort* algorithm[4] is chosen as the list will usually be ordered from previous J updates.

Algorithm 3 Worst Model for State-Action Pair (s,a)

```
WORSTMODEL($s$: state, $a$: action)
$R = (s'_1, \cdots, s'_k) =$ REACHABLESTATES($s$,$a$)
SORT($R$)
$i = 1$, $bound = B_{\min}$
while $(bound - p_i^{\min} + p_i^{\max} < 1)$ do
    $bound \leftarrow bound - p_i^{\min} + p_i^{\max}$
    $Pr(s'_i) \leftarrow p_i^{\max}$
    $i \leftarrow i + 1$
end while
$r = i$
$Pr(s'_r) \leftarrow 1 - (bound - p_r^{\min})$
for all $i \in \{r+1, \ldots, k\}$ do
    $Pr(s'_i) \leftarrow p_i^{\min}$
end for
return $(R, Pr(\cdot))$
```

To summarise, Robust VI on an interval-based SSP consists of applying normal Value Iteration with the transition probabilities updated through Alg. 3.

We need only a single worst model to compute the corresponding Q-value. Yet, because several reachable states s'_i may have the same value $c(s,a,s'_i) + J(s'_i)$ as s'_r (we call this set of states S'_r), there may be an infinite number of equivalent worst local models. Any model differing only in how the probability mass is distributed among the equally bad states of S'_r is also a worst local model.

Worst Global Models — Contrary to VI, RTDP does not necessarily visit the complete state-space. This is why [Barto *et al.*, 1995] introduces the notion of *relevant* states, which we extend to the uncertain case: a state s is *relevant* for $\mathcal{M}$ if

[4] http://en.wikipedia.org/wiki/Insertion_sort

there exists a start state s_0, a model $m \in \mathcal{M}$ and an optimal policy π under that model such that s can be reached from state s_0 when the controller uses π under m.

This notion is important because two equally worst local models on a given state-action pair may forbid different states, so that for two models m_1 and m_2, a state may be relevant in m_1 but not in m_2. Yet, RTDP should not find an optimal policy just for the relevant states of only one worst global model. Neither does the policy have to apply to all possible states. It should apply to all *reachable* states under *any* model (for optimal policies) i.e., for relevant states. But covering all relevant states in the worst model used for updating Q-values does not necessarily cover all relevant states in $\mathcal{M}$: it depends on the model used to choose the next state, i.e., to simulate the system's dynamics.

To avoid missing relevant states, each local model used for the simulation should ensure that all reachable states can be visited. As can be seen in Fig. 2, the set of possible local models for a state-action pair is an n-dimensional convex polytope. Any model *inside* this polytope, excluding the boundary, is therefore adequate because, for all s'_i, it ensures that $P(s'_i|s,a) > 0$.

So there exists a global model m_d that can be used to effectively simulate the system's dynamics without missing any potentially reachable state.

3.1 Robust (Trial-Based) RTDP

Robust RTDP differs from the original RTDP in that:

- Each time the algorithm is updating a state's evaluation, an opponent is looking for a worst local model which serves to compute the Q-values.
- For exploration purposes, the algorithm follows dynamics of the system that consider all possible transitions (using model m_d).
- "Relevant" states are now the states reachable by following any optimal policy under any possible model.

From this, we adapt to our context convergence Theorem 3 from [Barto *et al.*, 1995] and the corresponding proof, discussing mainly its changes.

Proposition 1. *In* **uncertain** *undiscounted stochastic shortest path problems,* **robust** *Trial-Based RTDP with the initial state of each trial restricted to a set of start states, converges (with probability one) to J^* on the set of relevant states, and the controller's policy converges to an optimal policy (possibly nonstationary) on the set of relevant states, under the same conditions as Theorem 3 in [Barto* et al.*, 1995].*

Proof. The proof outline is:

1. As shown in Sec. 2.3, robustness is achieved by considering a Stochastic Shortest Path *Game* (SSPG). In particular, this one is a *sequential* SSPG, for which Bertsekas and Tsitsiklis [1996] establish value iteration and Q-learning algorithms. Provided all states are visited infinitely often, *asynchronous* value iteration can also be used to solve sequential SSPGs (AVI-SSPG).
2. We adapt the proof of Barto *et al.* [1995] to transform AVI-SSPG to an RTDP algorithm (RTDP-SSPG). The transformation ensures: 1- the additional $\max_{m \in \mathcal{M}}$ in the update formula (Alg. 2, Line 5) does not break the requirement that J_t is increasing and non-overestimating; and 2- *relevant* states are visited infinitely often, guaranteeing a complete and optimal policy.
3. RTDP-SSPG and robust RTDP differ because RTDP-SSPG assumes an optimal opponent. In the robust case we wish to plan for opponents that may choose the wrong model. Thus robust RTDP is RTDP-SSPG following model m_d — for simulation only — that allows states to be visited that RTDP-SSPG optimal opponents would forbid. Under m_d (described in Sec. 3), no state reachable under any model in $\mathcal{M}$ is excluded from the search. Thus we have redefined an expanded set of relevant states for added robustness. All relevant states are still visited infinitely often.

□

Our use of the model m_d for simulation ensures that the policy will cover all states *any* opponent could lead us to, but assumes the worst model will apply afterwards.

4 Experiments

Labelled RTDP [Bonet and Geffner, 2003] is a modified version of RTDP which can be made robust in a similar way. The experiments conducted illustrate the behavior of *robust LRTDP*. To that end, it is compared to Bagnell *et al.*'s *Robust Value Iteration*, as well as *LRTDP*. In all cases, the convergence criteria is $\epsilon = 10^{-3}$ for LRTDP, and for VI we stop when the maximum change in a state long-term cost over one iteration is less than 10^{-3}.

4.1 Heart

In this first experiment, we compare a non-robust optimal policy with a robust one on the small example of Fig. 1-b. Table 1 shows the theoretical expected long-term costs of each policy on the normal (most probable) model, as well as the pessimistic and optimistic ones. The robust policy is largely better in the pessimistic case.

Table 1: Theoretical evaluation of robust and non-robust policies on various models, which match empirical evaluation.

	Normal	Pessimistic	Optimistic
Non-robust	2.90	8.90	1.70
Robust	3.33	3.33	3.33

4.2 Mountain-Car

We use here the mountain-car problem as defined in [Sutton and Barto, 1998]: starting from the bottom of a valley, a car has to get enough momentum to reach the top of a mountain (see Fig. 3). The same dynamics as described in the mountain car software[5] have been employed. The objective is to minimize the number of time steps to reach goal.

[5] http://www.cs.ualberta.ca/~sutton/··· MountainCar/MountainCar.html

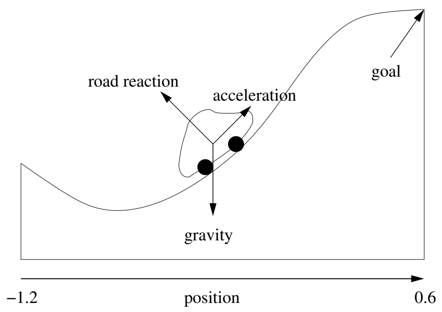


Figure 3: The mountain-car problem.

The continuous state-space is discretized (32×32 grid) and the corresponding uncertain model of transitions is obtained by sampling 1000 transitions from each state-action pair (s, a). For each transition, we computed intervals in which the true model lies with 95% confidence.

Results — Preliminary remark: simulating a path generally shows a car oscillating several times before leaving the valley. This has two main explanations: 1- the speed gathering is just sufficient to reach the summit, but not over it; and 2- the discretized model is not accurate enough and applying the policy obtained on the true mathematical model (instead of the discretized one) should be much better.

Fig. 4 shows the long-term cost function obtained by using *value iteration*, LRTDP, and their robust counterparts on the mountain-car problem. The x and y axes are the car's position and speed. The z axis is the expected cost to the goal. On the surface is an example path from the start state to a goal state: it follows the greedy policy under the average model.

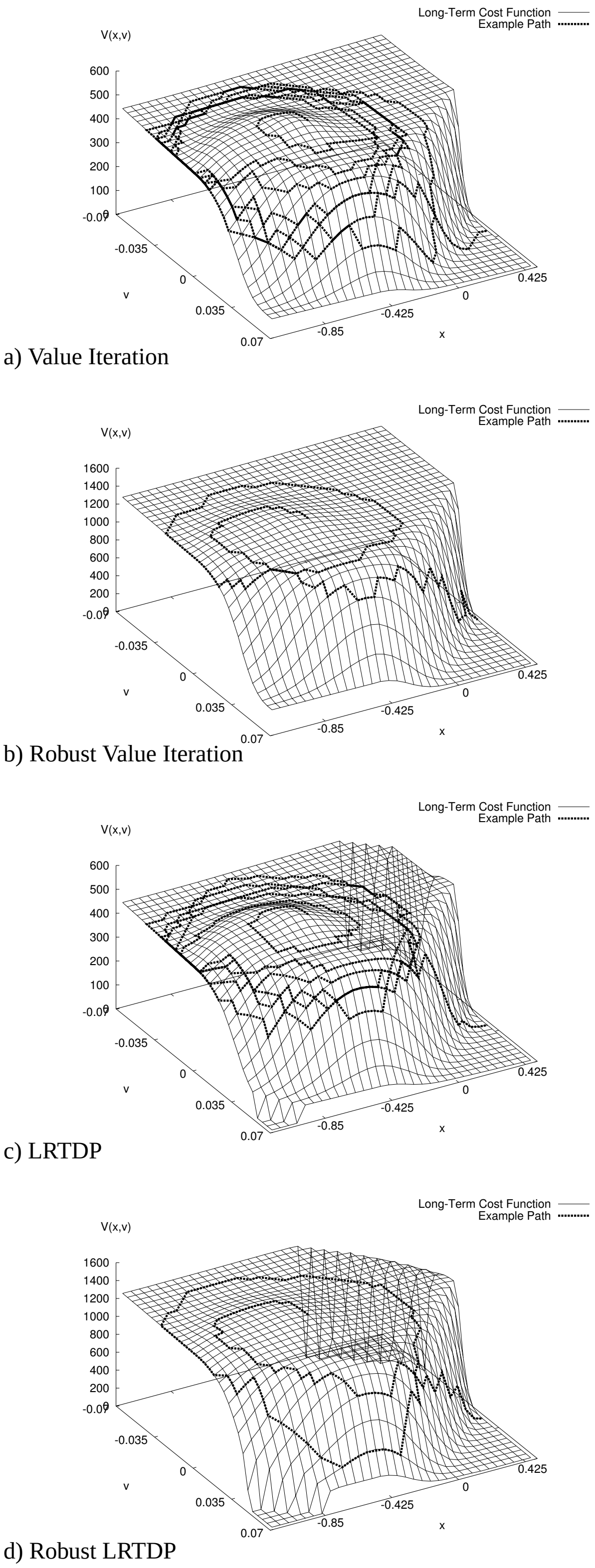
The general shape of the surface obtained is always the same, with some unexplored parts of the state-space in LRTDP and Robust LRTDP (as expected). The vertical scales are much larger in robust cases. This reflects the fact that reaching the goal is much more time-consuming under a pessimistic model. Because J can here be interpreted as the average time to the goal, these graphs show how a small uncertainty can lead to longer policies. Here the times are multiplied by more than 2.5.

While executing the four different algorithms, an evaluation of the current greedy policy was made every $10 * nStates = 10\,240$ Q-value updates. The result appears in Fig. 5, the y axis being the expected cost to the goal from the start state. On both sub-figures, LRTDP-based algorithms obtain good policies quickly, but have slow convergence times of VI=2.83×10^6 updates, LRTDP=6.76×10^6, rVI=8.31×10^6, rLRTDP=11.06×10^6.

Other experiments [Buffet and Aberdeen, 2004] also confirm these results, with one example showing that LRTDP can have a faster convergence than VI, and another one illustrating this approach on a temporal planning problem.

5 Discussion and Conclusion

A simple extension to this work, suggested by Hosaka *et al.* [2001], is to find the set of worst models as we do in this



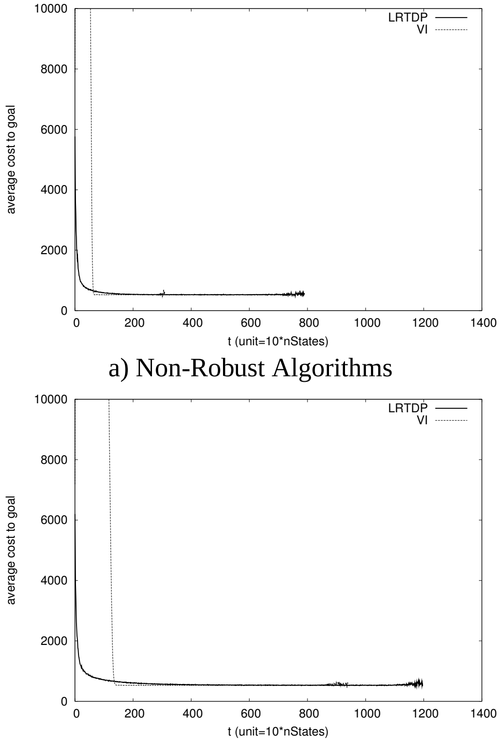
a) Value Iteration

b) Robust Value Iteration

c) LRTDP

d) Robust LRTDP

Figure 4: Long-term cost functions for the mountain-car problem. In each case, an example path is generated based on the most likely model.



a) Non-Robust Algorithms

b) Robust Algorithms

Figure 5: Average cost to goal for the mountain-car problem, measured every $10 * nStates$ updates of Q-values.

paper, but then to choose from these the model that has the lowest cost under an *optimistic model*. This improvement is presented in [Buffet and Aberdeen, 2004].

The approach to robustness adopted in this paper considers the knowledge of a set of possible models. An open question is whether it is possible to use more information from our uncertain model, by taking the probability distribution over possible models into account.

Model uncertainty has similarly been considered in model learning while planning [Strehl and Littman, 2004]. The algorithm proposed is optimistic, but does not seem to adapt well to our framework as an evolving model can break the “no-overestimation” assumption: $\forall s \in S,\ t \geq 0,\ J_t(s) \leq J^*(s)$. It is still important to notice that *robust* RTDP would not suffer on-line, as the real dynamics can be employed to pick a next state (the worst model appearing only in the long-term cost update formula).

Finally, a crucial assumption in RTDP is that a goal state must be reachable from any state. We present an algorithm tackling this problem in [Buffet, 2004].

Conclusion — Recent works show that model uncertainty is a major issue in decision-theoretic planning. We have proposed a modification of the RTDP algorithm enabling it to compute robust policies efficiently in large and uncertain domains. Model uncertainty is represented through confidence intervals on transition probabilities. The convergence of the resulting algorithm is sketched. We demonstrate robust LRTDP on a domain where statistics are used to estimate appropriate intervals.

Acknowledgments

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