A Classification of Weak Asynchronous Models of Distributed Computing

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- All nodes run the same deterministic finite-state machine 80
- Task: decide some graph property by consensus.



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Is there a red node?



Easy if nodes can change their answer.

 ${\tt Esparza \& Reiter - Weak \ Asynchronous \ Models \ of \ Distributed \ Computing}$

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- Task: decide some graph property *by consensus*.
- Consistency condition: all legal runs must yield the same answer.

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Is the graph a star?



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Is the graph a star?



Easy if nodes can count their neighbors.

| Four parameters | | | |
|-----------------|------------|-----------|----------|
| Detection | Acceptance | Selection | Fairness |

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|--|------------|-----------|----------|
| d : non-counting | | | |
| $A \qquad B \\ \{Sees: \\ \{A, B\} \}$ | | | |

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|--|------------|-----------|----------|
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| | | | |
| D : counting | | | |
| $\begin{array}{c} \textbf{A} \textbf{B} \\ \textbf{Sees:} \\ \{\{A, A, B\}\} \end{array}$ | | | |

| Detection | Acceptance | Selection | Fairness |
|--|---|-----------|----------|
| d : non-counting | a: halting | | |
| $ \begin{array}{c} A \\ \hline \\ \\ \hline \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\$ | ? \rightarrow ? \rightarrow Yes Answers are final. | | |
| D : counting | | | |



| Detection | Acceptance | Selection | Fairness |
|---|---|-----------|----------|
| d : non-counting | a: halting | | |
| $ \begin{array}{c} \textbf{A} \textbf{A} \textbf{B} \\ \hline \textbf{Sees:} \\ \{A, B\} \end{array} $ | ? \rightarrow ? \rightarrow Yes Answers are final. | | |
| D : counting | A: stabilizing | | |
| A A B { <i>Sees:</i> {{(A, A, B}} | No \rightarrow Yes \rightarrow Yes Answers can change. | | |
| | | | |






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Synchronous selection ***\$* < Exclusive selection **Sf














































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