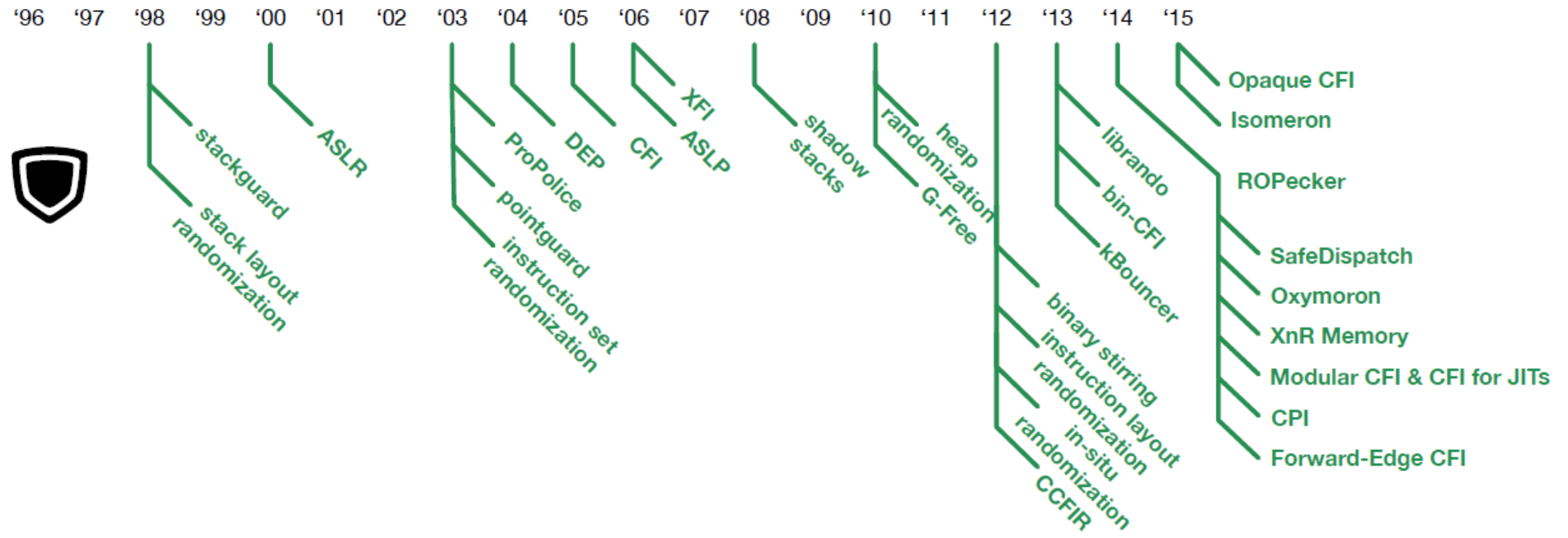
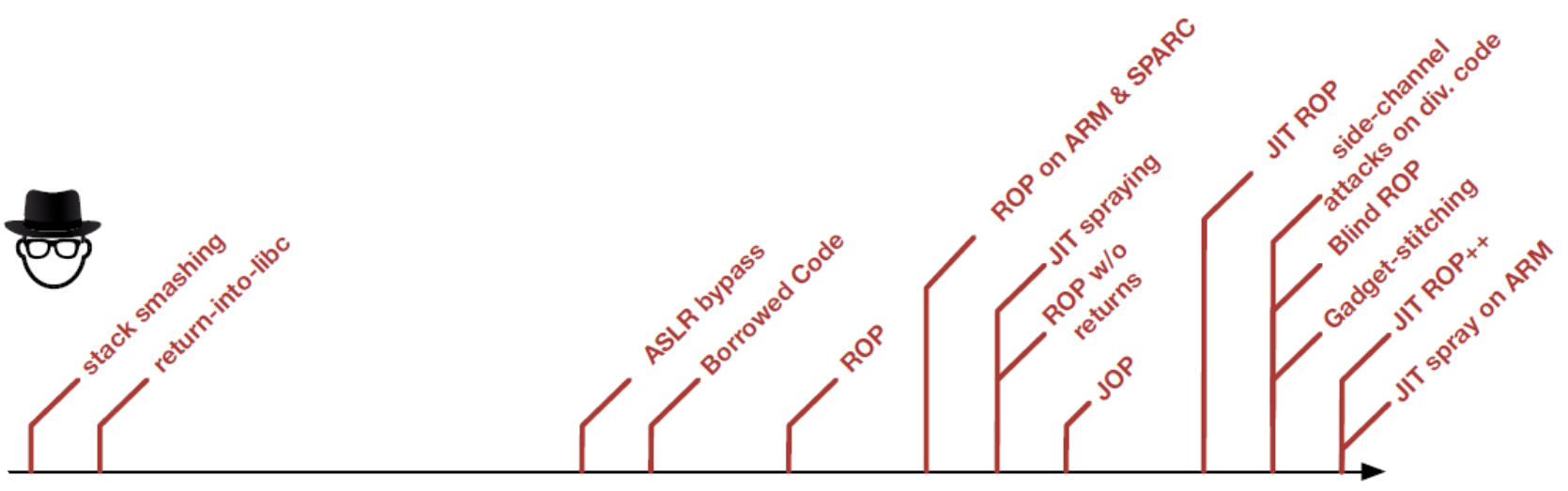


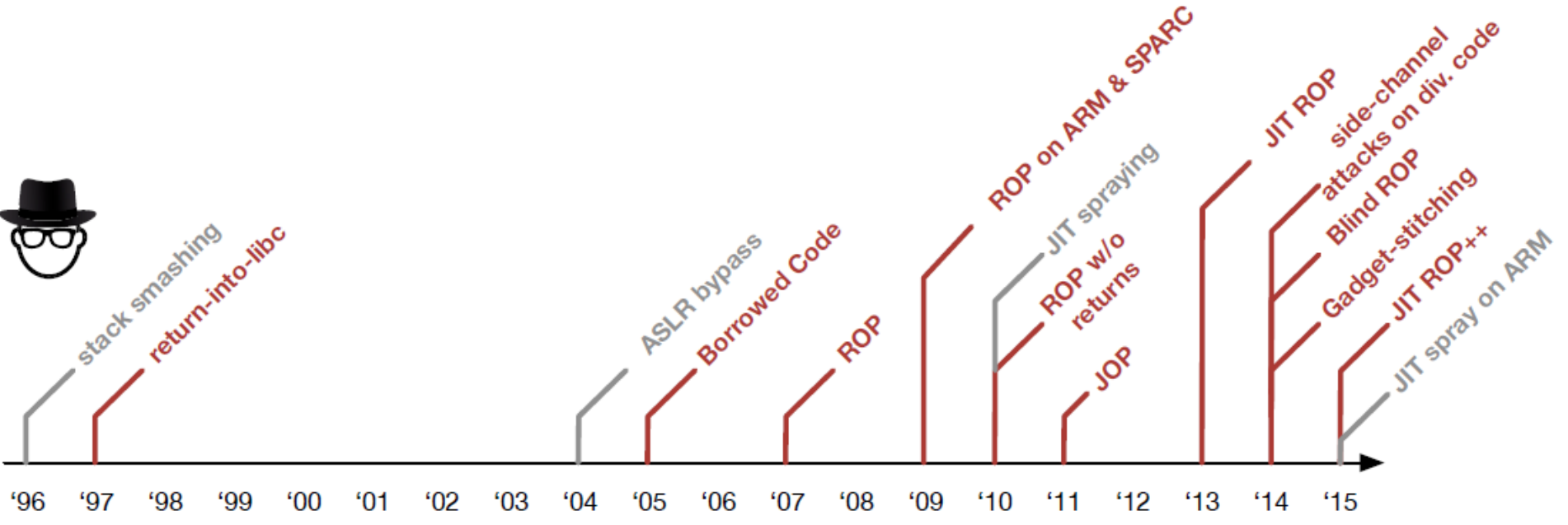
Opaque Control-Flow Integrity

Vishwath Mohan[‡], Per Larsen[§], Stefan Brunthaler[§], Kevin W. Hamlen[‡], Michael Franz[§]

The University of Texas at Dallas[‡]

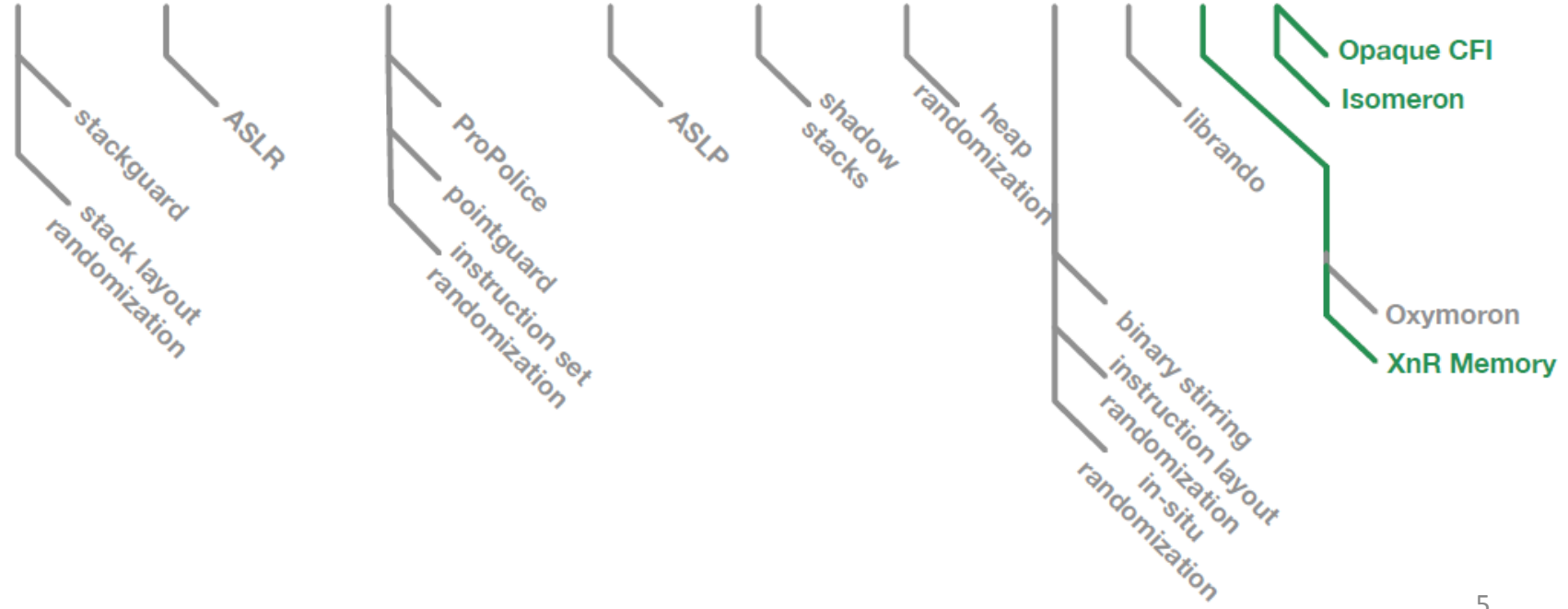
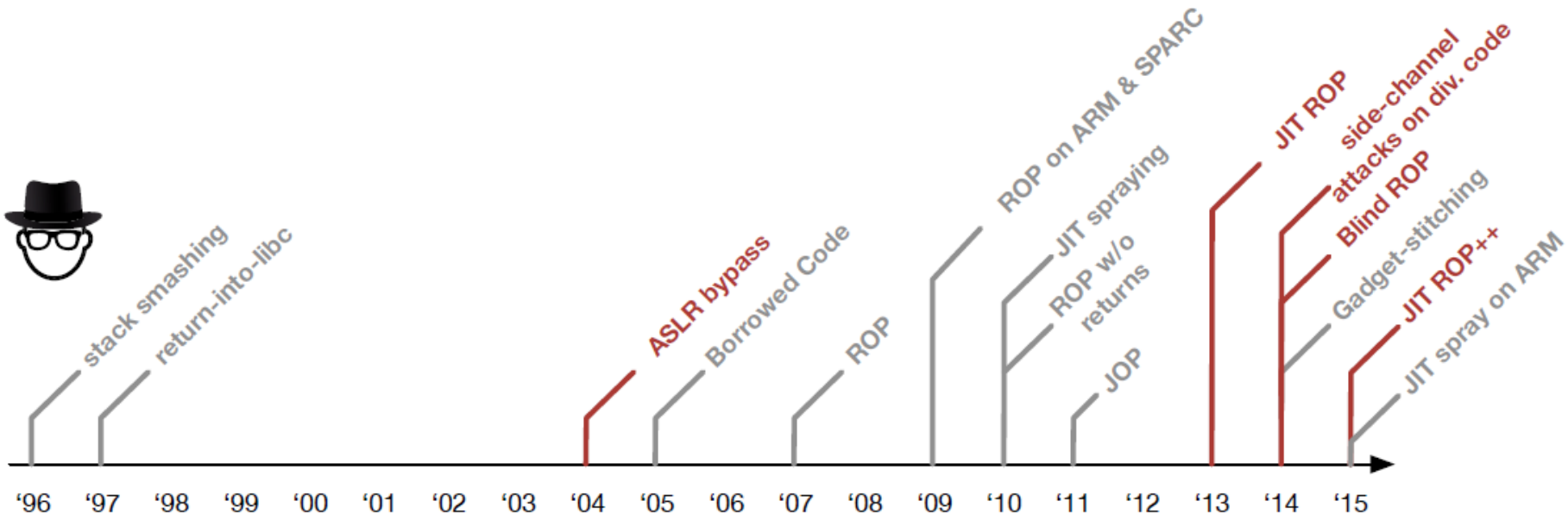
University of California, Irvine[§]

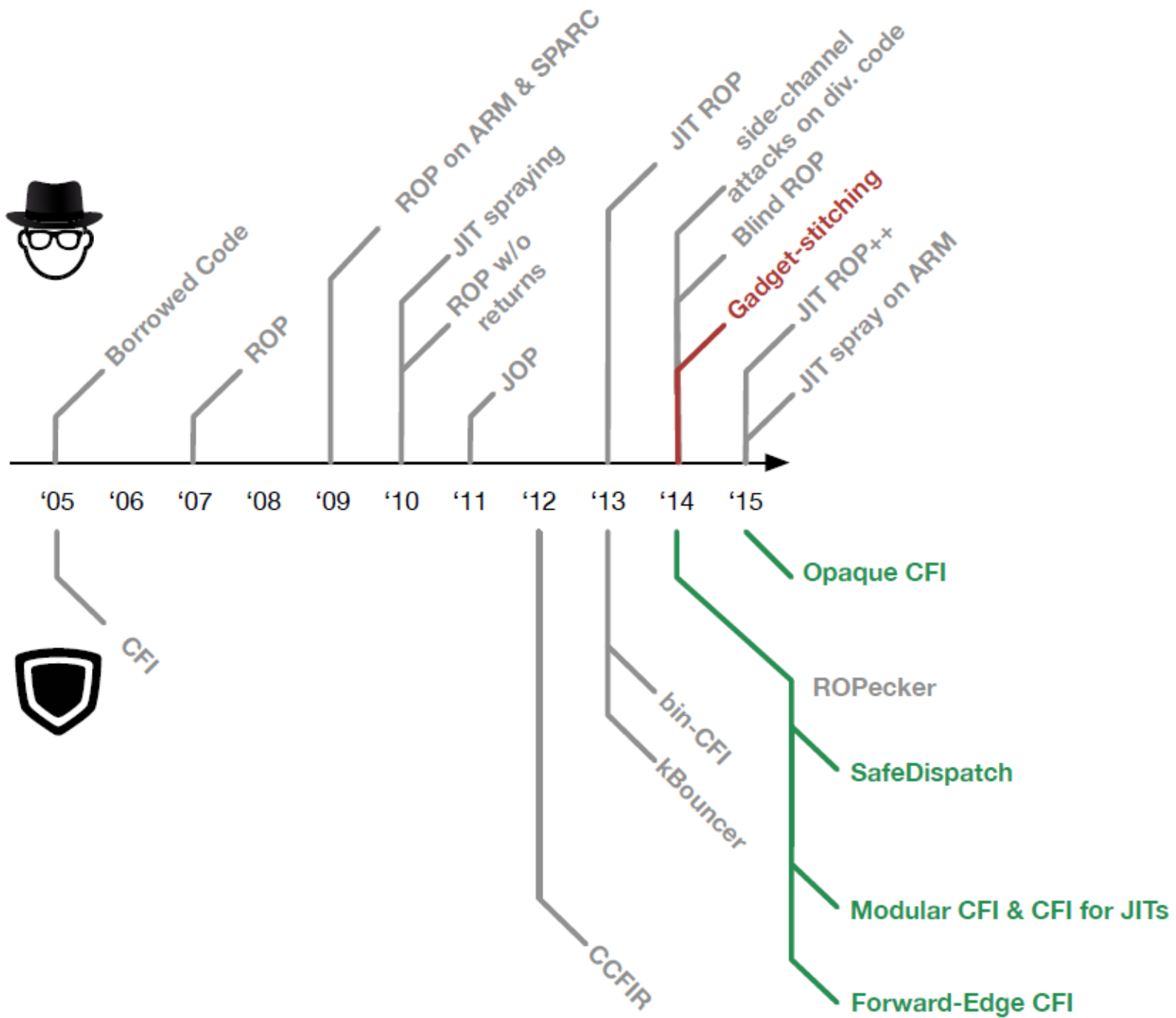




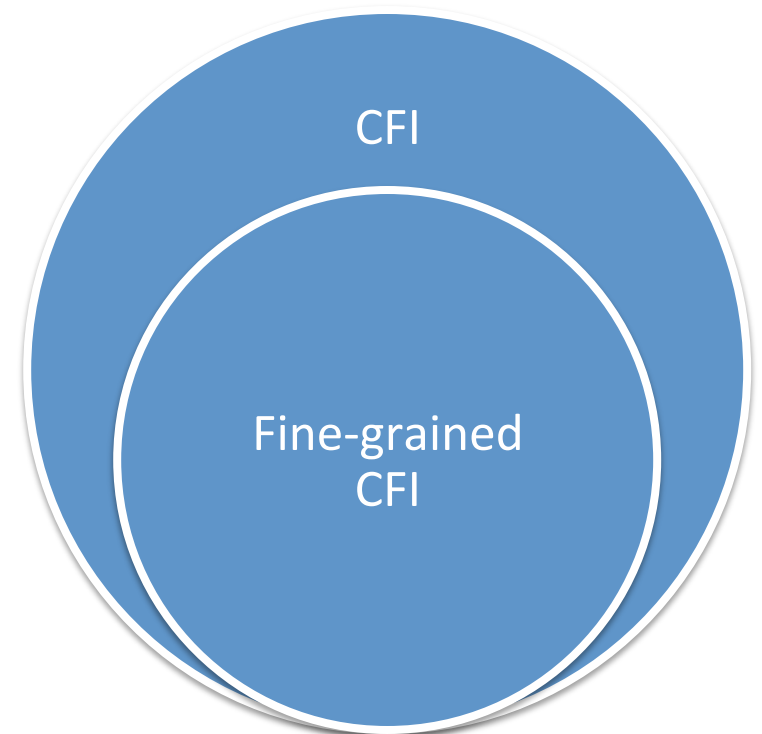
Code Reuse Attacks

- Needs
 - Location of code
 - Hijack control-flow
- Defensive options
 - Randomization
 - Control Flow Integrity

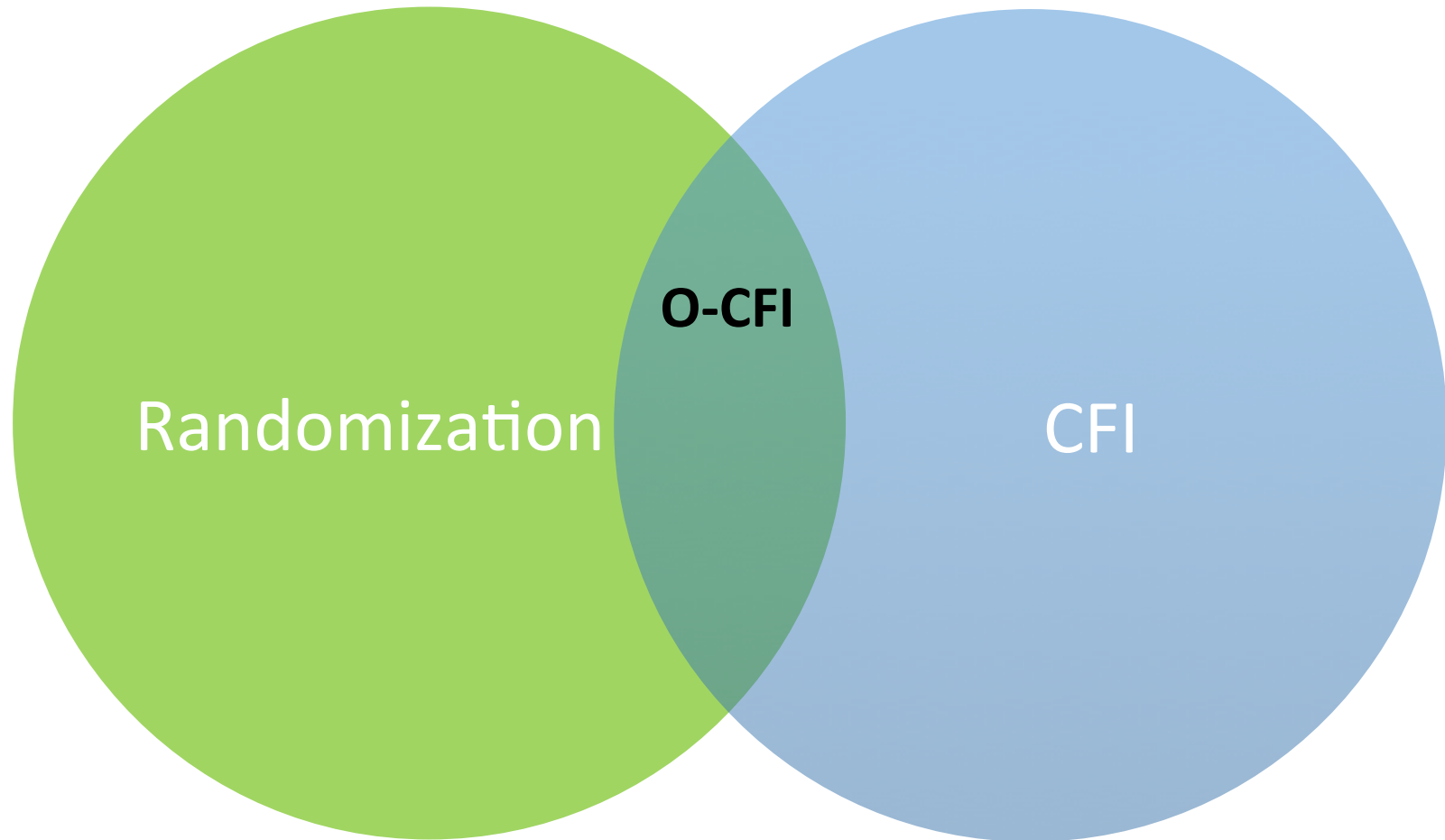


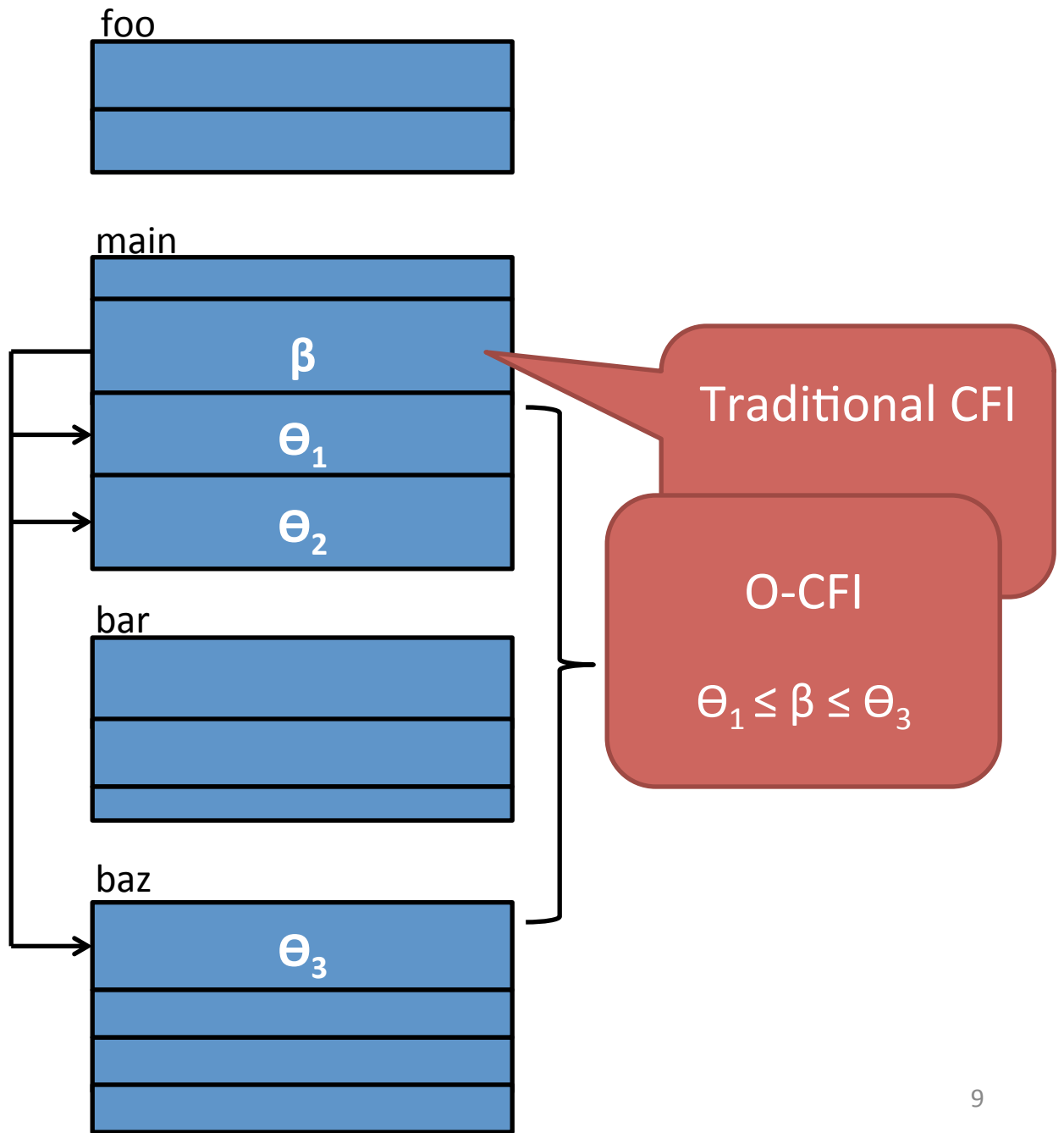


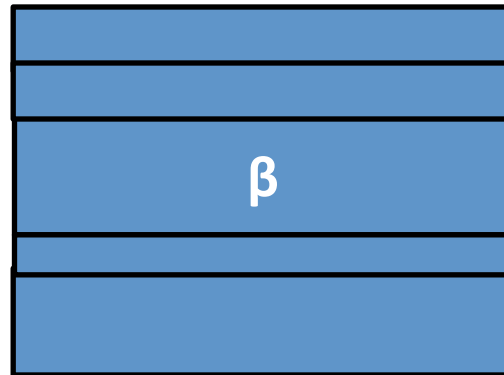
Where does that leave us?



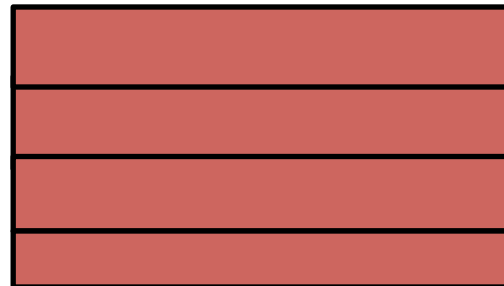
Opaque Control-Flow Integrity



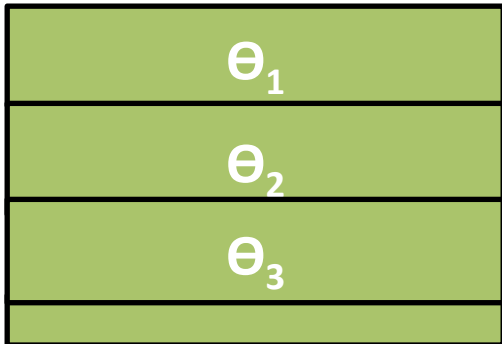
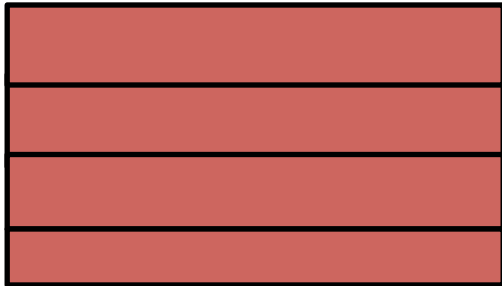




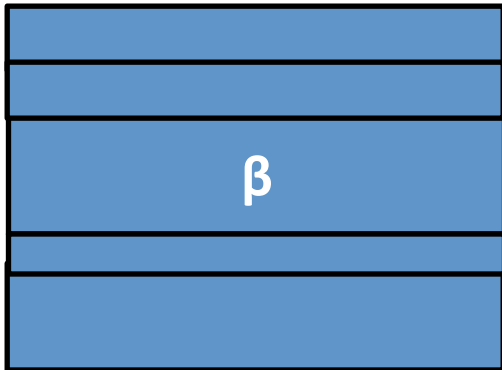
Randomization I:
Shuffle clusters



BLT



Randomization II:
Shuffle basic
blocks



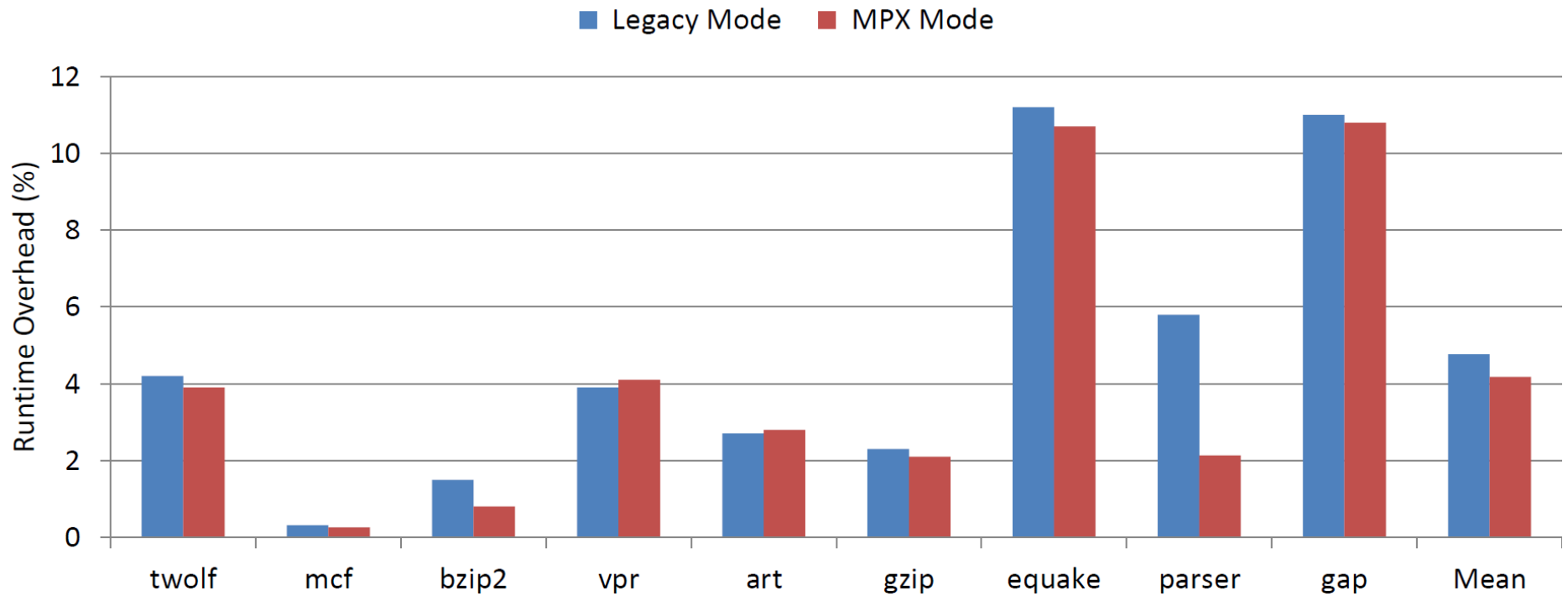
Accelerated Bounds Checks

- MPX mode on supported chipsets

Syntax	Description
<code>bndmov bnd, m64</code>	Move upper and lower bound from m64 to bound register bnd.
<code>bndcl bnd, r/m32</code>	Generate a #BR if r/m32 is less than the lower bound in bnd.
<code>bndcu bnd, r/m32</code>	Generate a #BR if r/m32 is higher than the upper bound in bnd.

- Legacy mode as fallback

Performance



Legacy mode overhead 4.7%

MPX mode overhead \approx 4.2%

Security

- When the CFI policy is opaque

Gadget Chain Size	Chance(%)
2	2.0
3	0.8
4	0.01
5	-

Security

- When the CFI policy is not opaque
 - Expressed as CSP
 - Attempted constructing `VirtualAlloc` payload
 - Across Mona/custom tool, no payload found

Conclusion

- Coarse-grained CFI with randomization
 - Advantages of both
- Effective against state-of-the-art exploits
 - JIT-ROP, BROP, Gadget stitching
- Efficient
 - 4.7% overhead in legacy mode

Thank you

Extra Resources

Optimizing Guards

- Actual guard implementation
 - PittSField inspired guards
 - Want minimal chunk size
 - Comparison instructions rather large (~ 7 bytes)
- How efficient can we be?

Optimizing Guards

Description	Original Code	Rewritten Code (MPX-mode)	Rewritten Code (Legacy-mode)
Indirect Branches	<code>call/jmp r/[m]</code>	<pre> 1: mov [esp-4], eax 2: mov eax, r/[m] 3: cmp byte ptr [eax], 0xF4 4: cmovz eax, [eax+1] — chunk boundary — 5: bndmov bndl, gs:[branch_id] 6: bndcu bndl, eax 7: jmp 9 — chunk boundary — 8: xor eax, eax 9: and al, align_mask 10: bndcl bndl, eax 11: xchg eax, [esp-4] 12: call/jmp [esp-4] </pre>	<pre> 1: push ecx 2: push eax 3: mov eax, r/[m] 4: cmp byte ptr [eax], 0xF4 5: cmovz eax, [eax+1] — chunk boundary — 6: mov ecx, branch_id 7: cmp eax, gs:[ecx] 8: jb 10 9: cmp gs:[ecx+4], eax 10: jbe abort — chunk boundary — 11: and al, align_mask 12: xchg eax, [esp] 13: pop ecx 14: pop ecx 15: call/jmp [esp-8] </pre>
Returns	<code>ret {n}</code>	<pre> — chunk boundary — 1: xchg eax, [esp] 2: and al, align_mask 3: bndmov bndl, gs:[branch_id] 4: jmp 8 — chunk boundary — 5: xor eax, eax 6: bndcu bndl, eax 7: bndcl bndl, eax 8: xchg eax, [esp] 9: ret {n} </pre>	<pre> — chunk boundary — 1: xchg eax, [esp] 2: cmp eax, gs:[branch_id] 3: jb 9 4: and al, align_mask — chunk boundary — 5: cmp eax, gs:[branch_id + 4] 6: jae 9 7: xchg eax, [esp] 8: ret {n} — chunk boundary — 9: jmp abort </pre>

Coarse Grained Insecurity

		CFI [1]	bin-CFI [50]	CCFIR [49]	kBouncer [33]	ROPecker [7]	ROPGuard [16]	EMET [30]
DeMott [12]	Feb 2014							☹
Göktaş et al. [18]	May 2014	☹	☹	☹				☹
Davi et al. [11]	Aug 2014		☹		☹	☹	☹	☹
Göktaş et al. [19]	Aug 2014				☹	☹		
Carlini and Wagner [6]	Aug 2014				☹	☹		